Java ConditionObject (Part 1)



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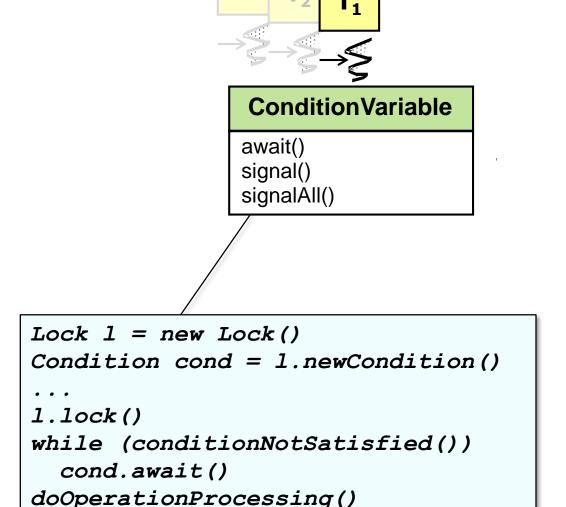
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Learning Objectives in this Part of the Lesson

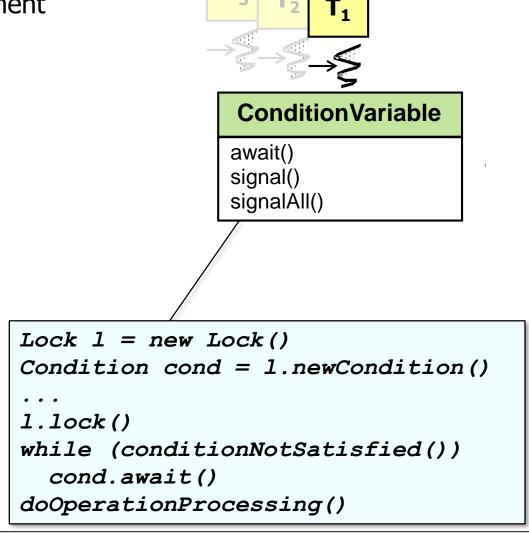
Understand what condition variables are



Learning Objectives in this Part of the Lesson

- Understand what condition variables are
- Know what pattern they implement





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```
ConditionVariable

await()
signal()
signalAll()
```

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Lock 1 = new Lock()
Condition cond = 1.newCondition()
...
1.lock()
while (conditionNotSatisfied())
cond.await()
doOperationProcessing()
```

Condition variables can be tricky, so I recommend you rewatch this lesson & read the links carefully

 A CV is a synchronizer that allows a thread to (repeatedly) suspend its execution until a condition is satisfied



Wheel of Pain - Conan the Barbarian

See <u>blog.dcoles.net/2012/02/understanding</u>
-how-to-use-condition.html

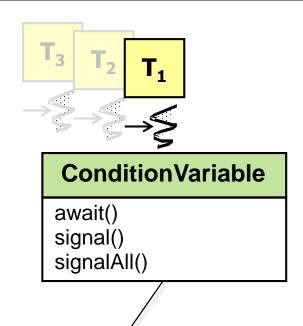
- A CV is a synchronizer that allows a thread to (repeatedly) suspend its execution until a condition is satisfied
 - A thread whose execution is suspended on a CV is said to be "blocked" on the CV



Tree of Woe – Conan the Barbarian

 A CV is implemented as a queue of threads that wait for certain condition(s) to be satisfied

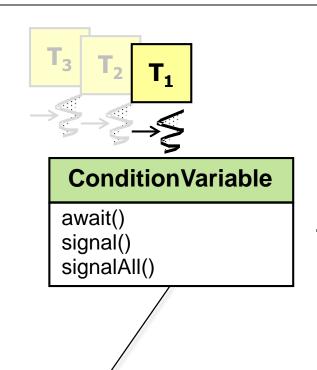




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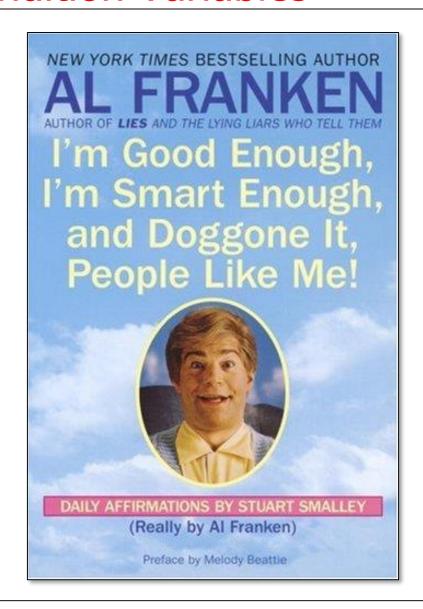
- A CV is implemented as a queue of threads that wait for certain condition(s) to be satisfied
 - This queue of threads is known as the "wait set"





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 - Inefficient use of resources
 - e.g., due to excessive "busy waiting"



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 - Inefficient use of resources
 - Insufficient to ensure *coordination*
 - e.g., what to do when a thread encounters shared state that it can't do any work upon (yet)

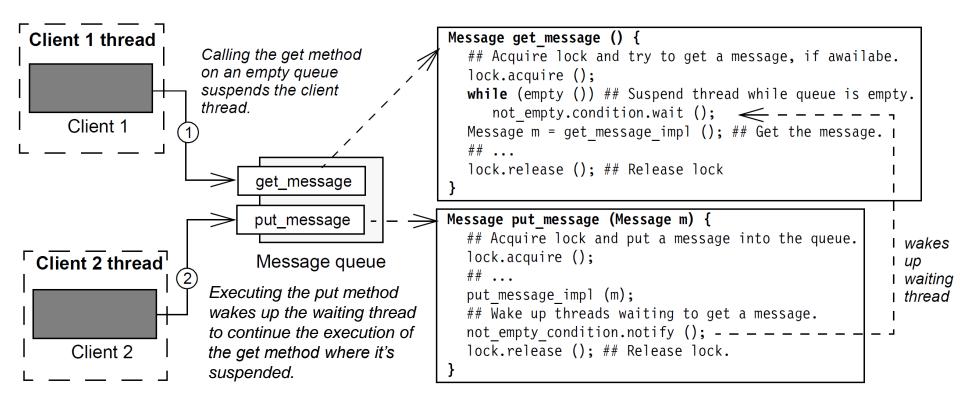


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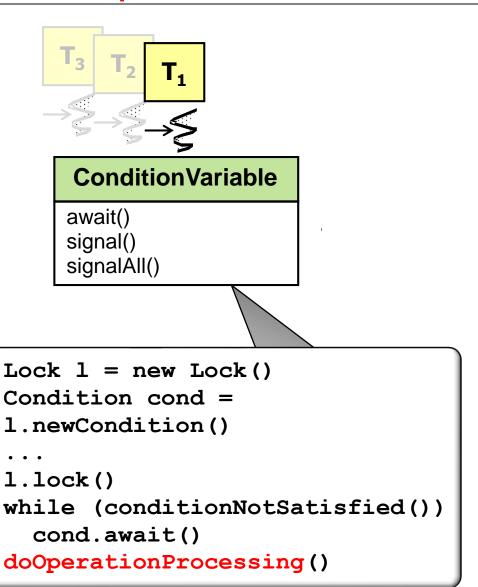


Waiting on an empty list

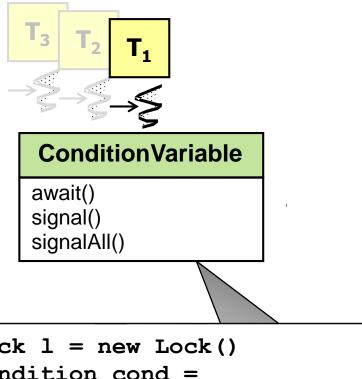
CVs are most often used to implement the Guarded Suspension pattern



 This pattern is applied to operations that can run only when a condition is satisfied

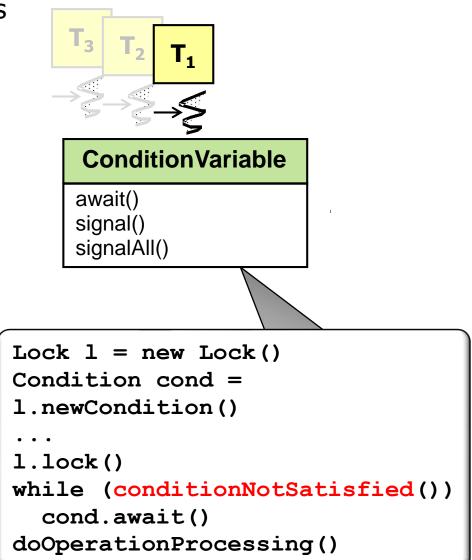


- This pattern is applied to operations that can run only when a condition is satisfied, e.g.,
 - a lock is acquired

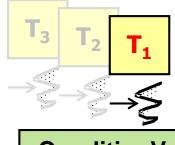


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- This pattern is applied to operations that can run only when a condition is satisfied, e.g.,
 - a lock is acquired
 - a precondition holds



 In this example thread T₁ uses a CV to suspend its execution until thread T_n notifies it that shared state it's waiting on *may* now be satisfied

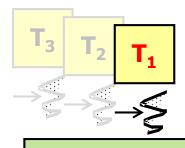




await() signal() signalAll()

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ConditionVariable

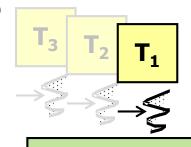
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Note the tentative nature of "may"...

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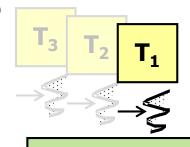
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First, a lock must be acquired..

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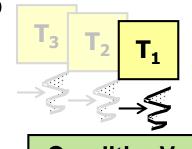
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Second, a condition is checked (in a loop) with the lock held..

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- In this example thread T₁ uses a CV to suspend its execution until thread T_n notifies it that shared state it's waiting on *may* now be satisfied
 - A condition can be arbitrarily complex



ConditionVariable

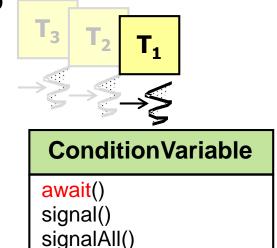
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e.g., a method call, an expression involving shared state, etc.

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Any state shared between threads must be protected by a lock associated with the CV

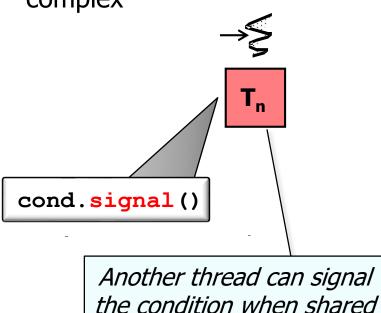
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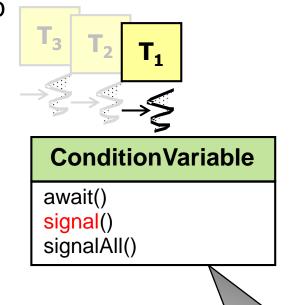
The calling thread will block (possibly repeatedly) while the condition is not satisfied (await() atomically releases the lock)

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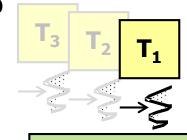


state may now be true



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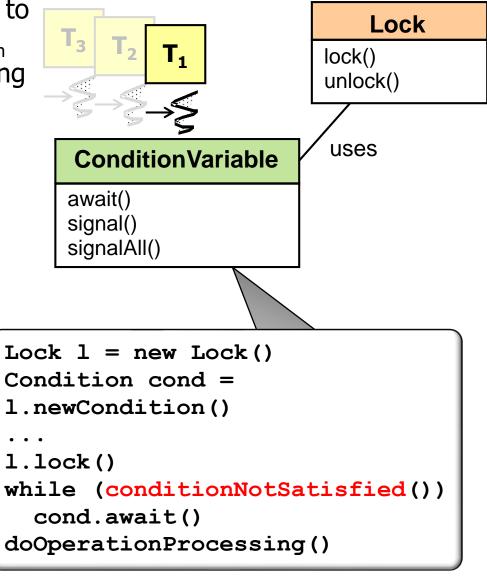
ConditionVariable

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await() reacquires the lock & the condition is rechecked in the loop

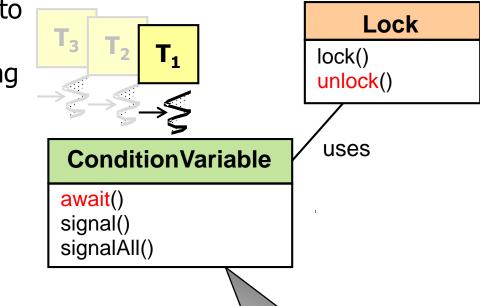
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 - A condition can be arbitrarily complex
 - Waiting on a CV releases the lock
 & suspends the thread atomically



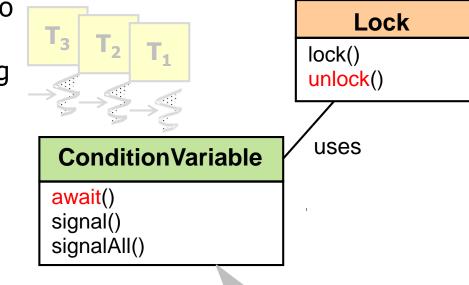


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The lock is released when the thread is suspended on the CV

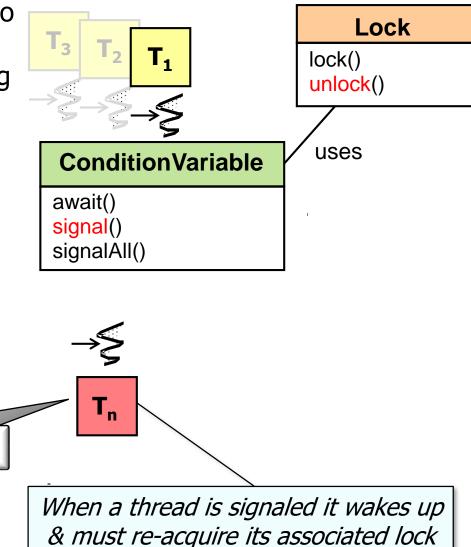
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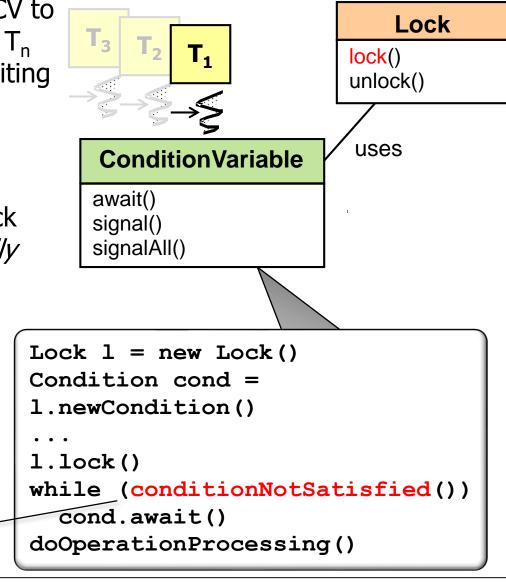
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cond.signal()

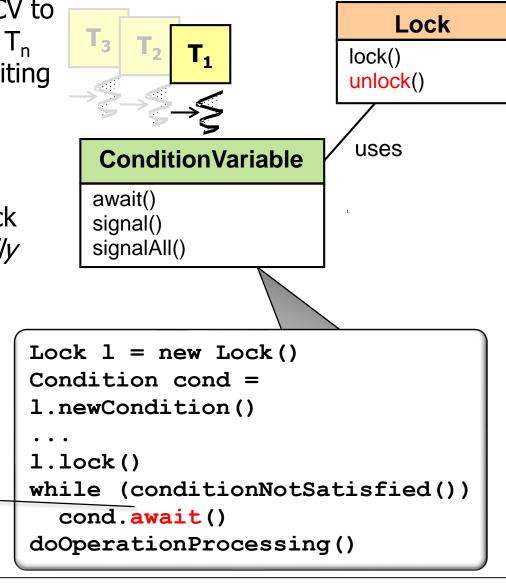
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After lock is re-acquired the thread can reevaluate its condition to see if it's satisfied



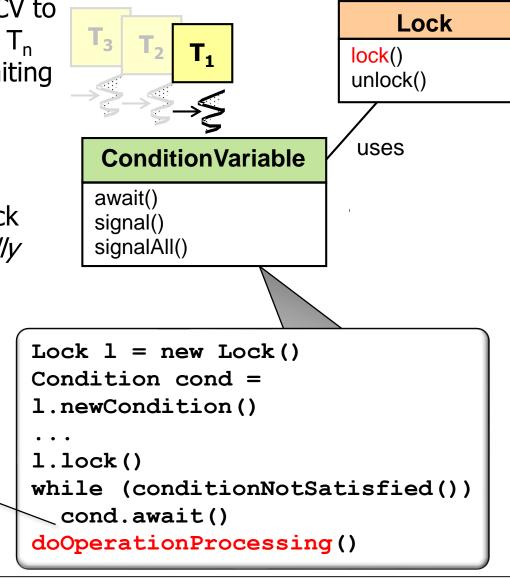
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If condition is not satisfied the thread must wait (which releases the lock atomically)



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 - A condition can be arbitrarily complex
 - Waiting on a CV releases the lock
 & suspends the thread atomically
 - Thread T₁ is suspended until thread T_n signals the CV

After the lock is re-acquired & the condition is satisfied the operation can proceed (with lock held)



End of Java ConditionObject (Part 1)