## Java Atomic Classes & Operations (Part 1)



Douglas C. Schmidt

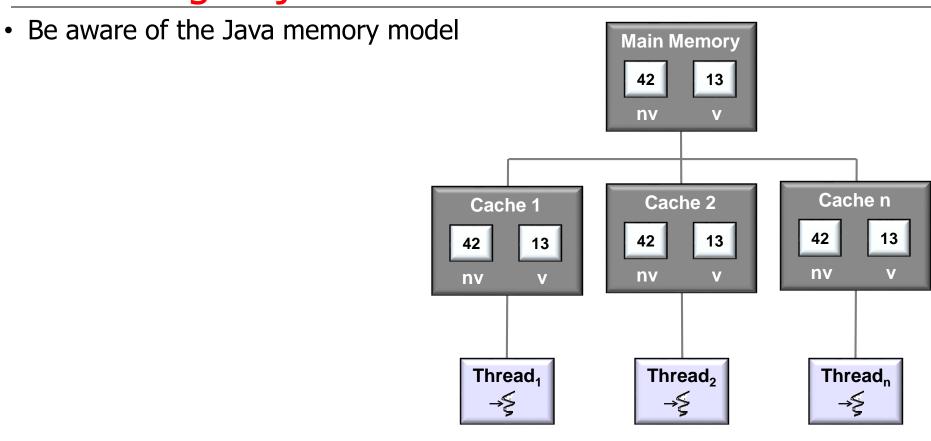
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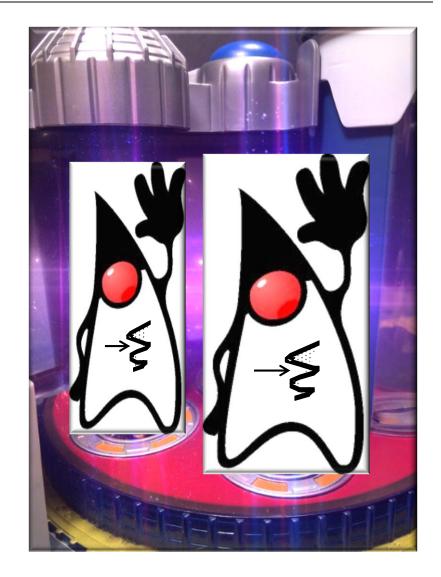


#### Learning Objectives in this Part of the Lesson



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- Be aware of the Java memory model
- Understand how Java atomic operations provide concurrent programs with lockfree, thread-safe mechanisms to read from & write to single variables



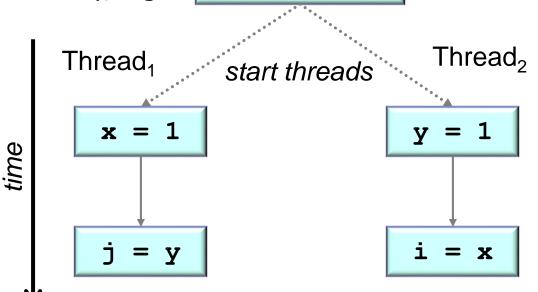
 Java's memory model defines semantics of **Main Memory** multi-threaded access to shared memory 13 42 nv Cache n Cache 2 Cache 1 13 42 13 13 nv V nv V nv Thread<sub>1</sub> Thread<sub>2</sub> Thread<sub>n</sub>

• Java's memory model defines semantics of multi-threaded access to shared memory, e.g.

x = y = 0

 Which instruction reorderings are allowed in memory

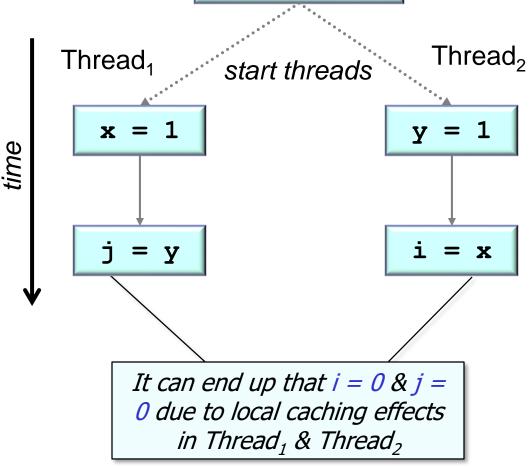
> OUT OF ORDER



There are a number of potential sources of reordering, e.g., the Java compiler, the JIT, & processor caches, etc.

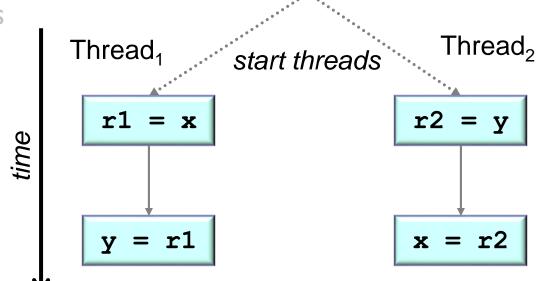
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- Which instruction reorderings are allowed in memory
  - Should not be overly restrictive, to enable hardware optimizations



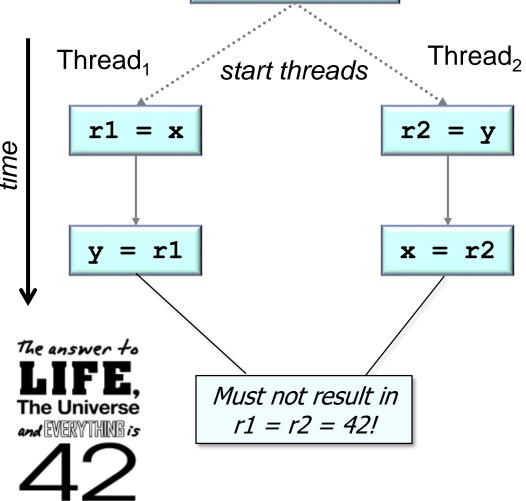
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- Which instruction reorderings are allowed in memory
- Which program outputs may occur in a correct JVM implementation

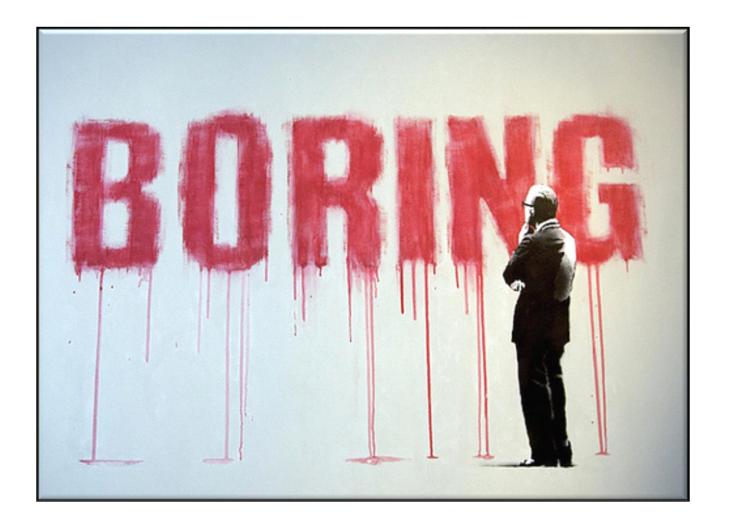


- Java's memory model defines semantics of multi-threaded access to shared memory, e.g.
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- Which instruction reorderings are allowed in memory
- Which program outputs may occur in a correct JVM implementation
  - Should not be too generous such that values appear randomly!



Reading about Java's memory model is as much fun as watching paint dry...



See <a href="https://www.cs.umd.edu/users/pugh/java/memoryModel/jsr-133-faq.html">www.cs.umd.edu/users/pugh/java/memoryModel/jsr-133-faq.html</a>

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Fortunately, you needn't understand all these memory model details – you just need to know how to use Java synchronizers properly!!

 The java.util.concurrent.atomic package several types of atomic actions on objects

#### Package java.util.concurrent.atomic

A small toolkit of classes that support lock-free thread-safe programming on single variables.

See: Description

Class	Description
AtomicBoolean	A boolean value that may be updated atomically.
AtomicInteger	An int value that may be updated atomically.
AtomicIntegerArray	An int array in which elements may be updated atomically.
AtomicIntegerFieldUpdater <t></t>	A reflection-based utility that enables atomic updates to designated volatile int fields of designated classes.
AtomicLong	A long value that may be updated atomically.
AtomicLongArray	A long array in which elements may be updated atomically.
AtomicLongFieldUpdater <t></t>	A reflection-based utility that enables atomic updates to designated volatile long fields of

See <u>docs.oracle.com/javase/8/docs/api/java/util/concurrent/atomic/package-summary.html</u>

- The java.util.concurrent.atomic package several types of atomic actions on objects
  - Atomic variables
    - Provide lock-free thread-safe operations on single variables



See <a href="mailto:docs.oracle.com/javase/tutorial/essential/concurrency/atomicvars.html">docs.oracle.com/javase/tutorial/essential/concurrency/atomicvars.html</a>

- The java.util.concurrent.atomic package several types of atomic actions on objects
  - Atomic variables
    - Provide lock-free thread-safe operations on single variables
      - e.g., AtomicLong supports atomic "compare-and-swap" operations

```
<<Java Class>>
               AtomicLong
AtomicLong(long)
AtomicLong()
get():long
set(long):void
flazySet(long):void
getAndSet(long):long
compareAndSet(long,long):boolean
weakCompareAndSet(long,long):boolean
getAndIncrement():long
getAndDecrement():long
getAndAdd(long):long
fincrementAndGet():long
decrementAndGet():long

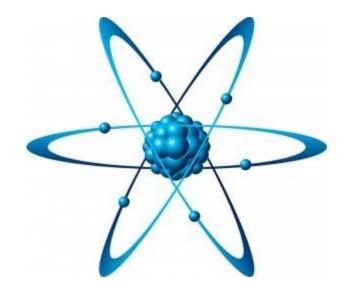
√ addAndGet(long):long

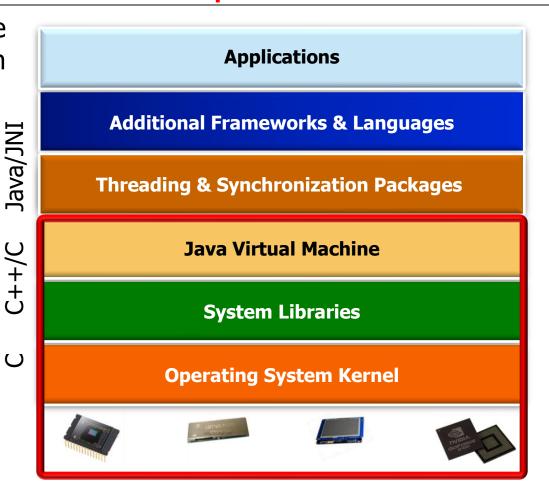
getAndUpdate(LongUnaryOperator):long
updateAndGet(LongUnaryOperator):long
getAndAccumulate(long,LongBinaryOperator):long
accumulateAndGet(long,LongBinaryOperator):long
toString()
intValue():int
longValue():long
floatValue():float
doubleValue():double
```

- The java.util.concurrent.atomic package several types of atomic actions on objects
  - Atomic variables
  - LongAdder
    - Allows multiple threads to update a common sum efficiently under high contention

- <<Java Class>>
- LongAdder
- √LongAdder()
- add(long):void
- increment():void
- decrement():void
- sum():long
- reset():void
- sumThenReset():long
- toString()
- longValue():long
- intValue():int
- floatValue():float
- doubleValue():double

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void lock(int \*mutex) {

continue;

```
int compareAndSwap(int *loc,
                                             int expected,
                                             int updated) {
                           START ATOMIC();
                           int oldValue = *loc;
                           if (oldValue == expected)
                              *loc = updated;
                           END ATOMIC();
                           return oldValue;
while (compareAndSwap(mutex, 0, 1) == 1)
```

The **lock**() method uses compareAndSwap() to implement mutual exclusion (mutex) via a "spin-lock"

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compareAndSwap() checks if the location pointed to by **mutex** is 0 & iff that's true it sets the value to 1

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                           if (oldValue == expected)
                              *loc = updated;
                           END ATOMIC();
                           return oldValue:
while (compareAndSwap(mutex, 0, 1) == 1)
```

If compareAndSwap() returns 1 that means the mutex is "acquired" so the loop keeps spinning

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int compareAndSwap(int *loc,
                                                 int expected,
                                                 int updated) {
                              START ATOMIC();
                              int oldValue = *loc;
                              if (oldValue == expected)
                                  *loc = updated;
                              END ATOMIC();
                              return oldValue:
void lock(int *mutex) {
  while (compareAndSwap(mutex, 0, 1) == 1)
    continue;
void unlock(int *mutex) {
  START ATOMIC();
                                 The unlock() method atomically
  *mutex = 0;
                                  resets the mutex value to 0
  END ATOMIC();
```

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int compareAndSwap(int *loc,
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void lock(int *mutex) {
  while (compareAndSwap(mutex, 0, 1) == 1)
void unlock(int *mutex) {
                                 The unlock() method atomically
                                   resets the mutex value to 0
```

```
START ATOMIC();
*mutex = 0;
END ATOMIC();
```

continue;

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  - CAS "compare-and-swap"

void lock(int \*mutex) {

void unlock(int \*mutex) {

```
int compareAndSwap(int *loc,
                                               int expected,
                                               int updated) {
                            START ATOMIC();
                            int oldValue = *loc;
                            if (oldValue == expected)
                               *loc = updated;
                            END ATOMIC();
                            return oldValue:
while (compareAndSwap(mutex, 0, 1) == 1)
                               The unlock() method atomically
                               resets the mutex value to 0
```

```
*mutex = 0;
END ATOMIC();
```

START ATOMIC();

continue;

- Atomic operations can be implemented other ways
  - e.g., "test-and-set"

```
int testAndSet(int *loc) {
   int oldValue;
   START_ATOMIC();
   oldValue = *loc;
   *loc = 1; // 1 == locked
   END_ATOMIC();
   return oldValue;
}
```

Test-and-set atomically modifies the contents of a memory location & returns its old value

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Test-and-set **atomically** modifies the contents of a memory location & returns its old value

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Test-and-set atomically modifies the contents of a memory location & returns its old value

- Atomic operations can be implemented other ways
  - e.g., "test-and-set"

Test-and-set can also be used to implement a spin-lock mutex

```
int testAndSet(int *loc) {
  int oldValue;
  START ATOMIC();
  oldValue = *loc;
  *loc = 1; // 1 == locked
  END ATOMIC();
  return oldValue;
void lock(int *loc) {
  while (testAndSet(loc) == 1);
void unlock(int *loc) {
  START ATOMIC();
  *loc = 0;
  END ATOMIC();
```

 compareAndSwap() provides a more general solution than the testAndSet()

```
int testAndSet(int *loc) {
  int oldValue;
  START ATOMIC();
  oldValue = *loc;
  *loc = 1; // 1 == locked
  END ATOMIC();
  return oldValue;
int compareAndSwap (int *loc,
                    int expected,
                    int updated) {
  START ATOMIC();
  int oldValue = *loc;
  if (oldValue == expected)
     *loc = updated;
  END ATOMIC();
  return oldValue:
```

- compareAndSwap() provides a more general solution than the testAndSet()
  - e.g., it can set the value to something other than 1 or 0

```
int testAndSet(int *loc) {
  int oldValue;
  START ATOMIC();
  oldValue = *loc;
  *loc = 1; // 1 == locked
  END ATOMIC();
  return oldValue;
int compareAndSwap(int *loc,
                    int expected,
                    int updated) {
  START ATOMIC();
  int oldValue = *loc;
  if (oldValue == expected)
     *loc = updated;
  END ATOMIC();
  return oldValue:
```

 One "human" known use of atomic operations is a Star Trek transporter



See en.wikipedia.org/wiki/Transporter\_(Star\_Trek)

- One "human" known use of atomic operations is a Star Trek transporter
  - Converts a person/object into an energy pattern & "beams" them to a destination where they're converted back into matter



- One "human" known use of atomic operations is a Star Trek transporter
  - Converts a person/object into an energy pattern & "beams" them to a destination where they're converted back into matter
  - This process must occur atomically or a horrible accident will occur!



 Another "human" known use of atomic operations is "apparition" in Harry Potter



See <a href="https://harrypotter.fandom.com/wiki/Apparition">harrypotter.fandom.com/wiki/Apparition</a>

- Another "human" known use of atomic operations is "apparition" in Harry Potter
  - If the user focuses properly they disappear from their current location & instantly reappear at the desired location



- Another "human" known use of atomic operations is "apparition" in Harry Potter
  - If the user focuses properly they disappear from their current location & instantly reappear at the desired location
  - However, "spinching" occurs if a wizard or witch fails to apparate atomically!



See harrypotter.fandom.com/wiki/Splinching

# End of Atomic Classes & Operations (Part 1)