

Java ReentrantLock

(Part 1)



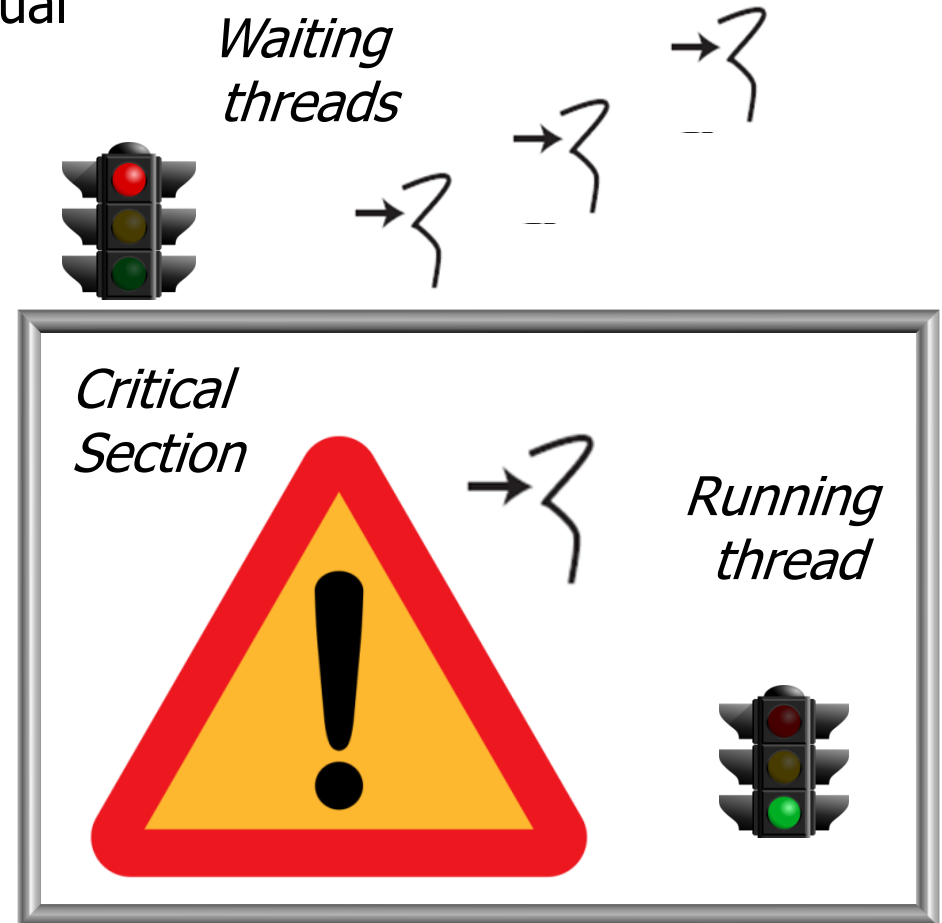
Douglas C. Schmidt
d.schmidt@vanderbilt.edu
www.dre.vanderbilt.edu/~schmidt

**Institute for Software
Integrated Systems
Vanderbilt University
Nashville, Tennessee, USA**



Learning Objectives in this Part of the Lesson

- Understand how the concept of mutual exclusion in concurrent programs



See en.wikipedia.org/wiki/Mutual_exclusion

Overview of Mutual Exclusion Locks

Overview of Mutual Exclusion Locks

- A mutual exclusion lock (mutex) defines a “critical section”

A critical section is group of instructions or region of code that must be executed atomically

*Critical
Section*



See en.wikipedia.org/wiki/Critical_section

Overview of Mutual Exclusion Locks

- A mutual exclusion lock (mutex) defines a “critical section”
 - Ensures only one thread can run in a block of code at a time

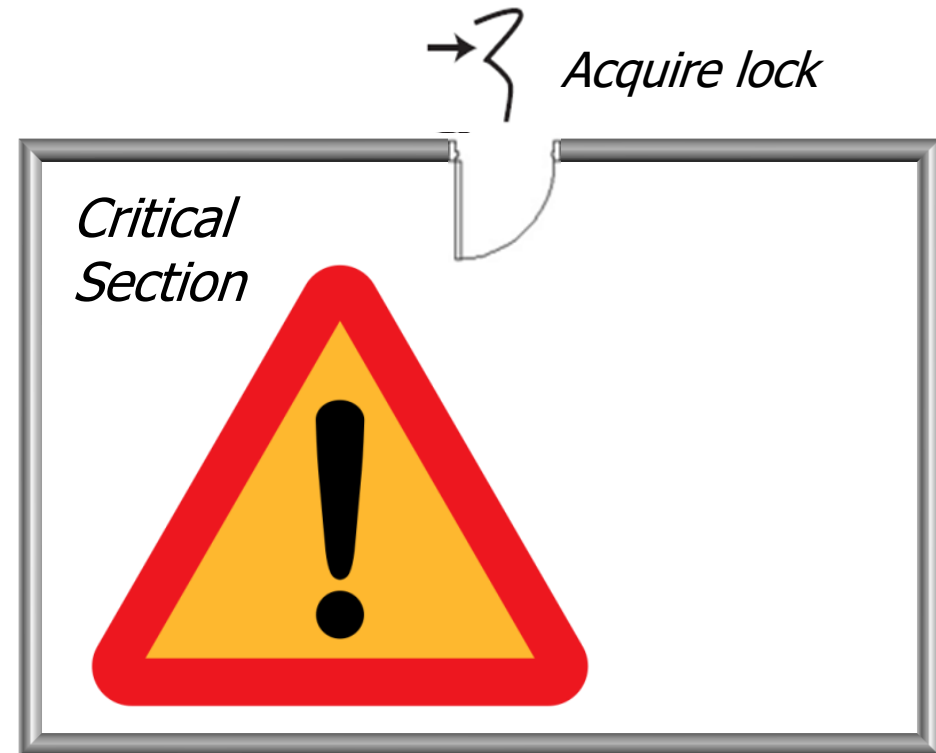


*Critical
Section*



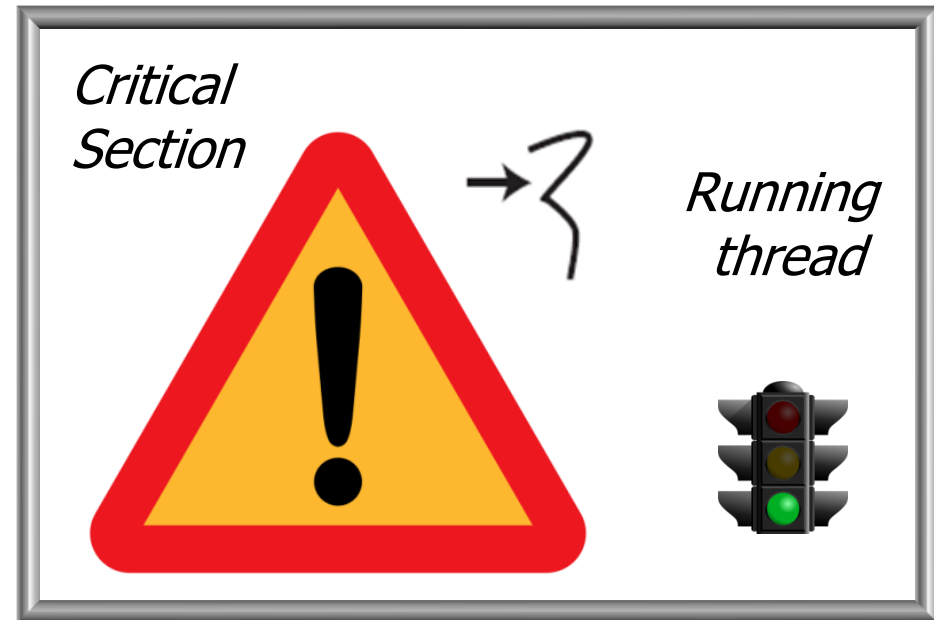
Overview of Mutual Exclusion Locks

- A mutual exclusion lock (mutex) defines a “critical section”
 - Ensures only one thread can run in a block of code at a time



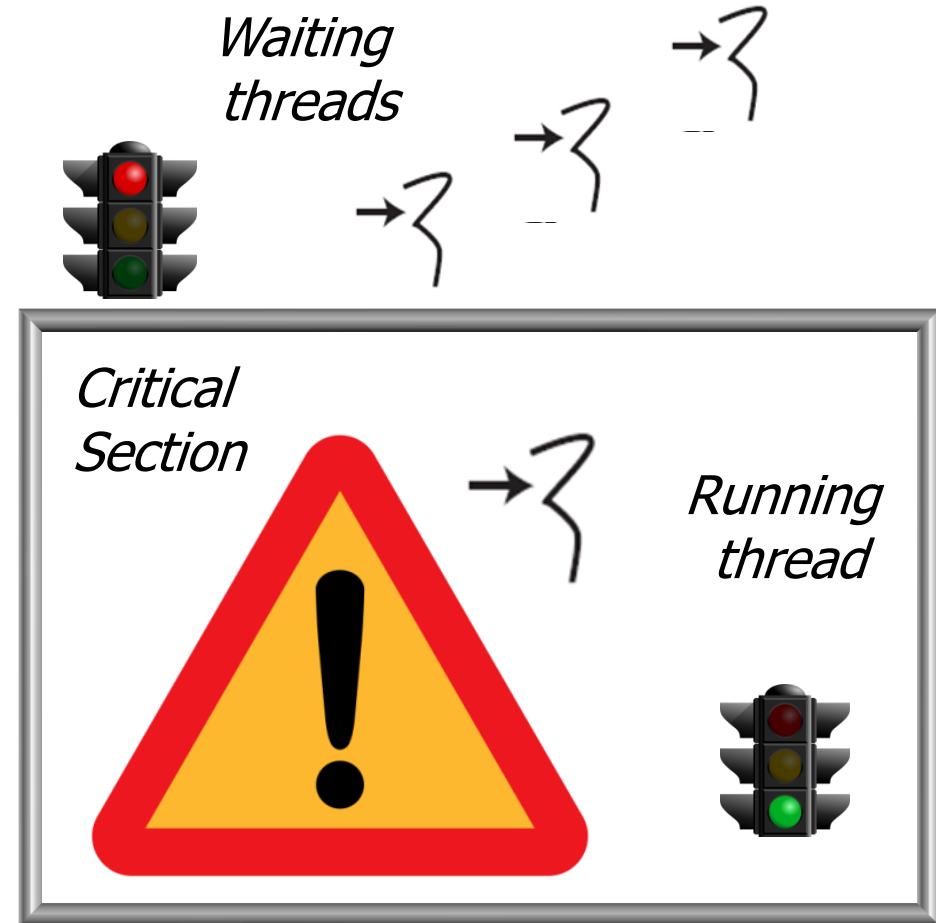
Overview of Mutual Exclusion Locks

- A mutual exclusion lock (mutex) defines a “critical section”
 - Ensures only one thread can run in a block of code at a time



Overview of Mutual Exclusion Locks

- A mutual exclusion lock (mutex) defines a “critical section”
 - Ensures only one thread can run in a block of code at a time
- Other threads are kept “at bay”
 - Prevent corruption of shared (mutable) data that can be set/get by concurrent operations



Other threads must obey the locking protocol or chaos will ensue!!

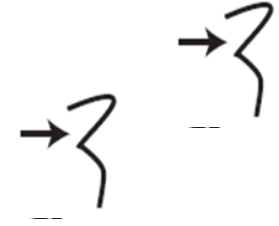
Overview of Mutual Exclusion Locks

- A mutual exclusion lock (mutex) defines a “critical section”
 - Ensures only one thread can run in a block of code at a time
- Other threads are kept “at bay”
 - Prevent corruption of shared (mutable) data that can be set/get by concurrent operations
- Race conditions could occur if multiple threads run within a critical section

Race conditions can arise when a program depends on the sequence or timing of threads for it to operate properly



Waiting threads



Critical Section



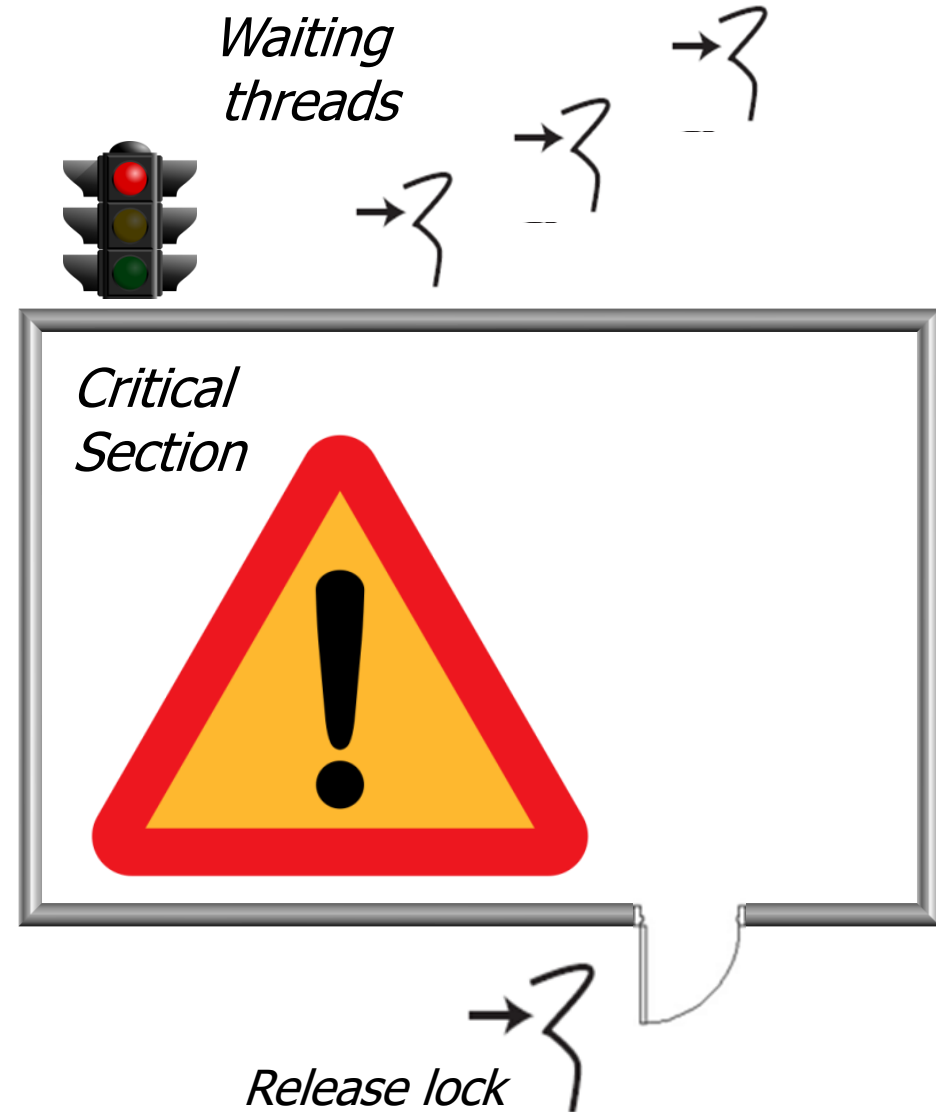
Running threads



See en.wikipedia.org/wiki/Race_condition

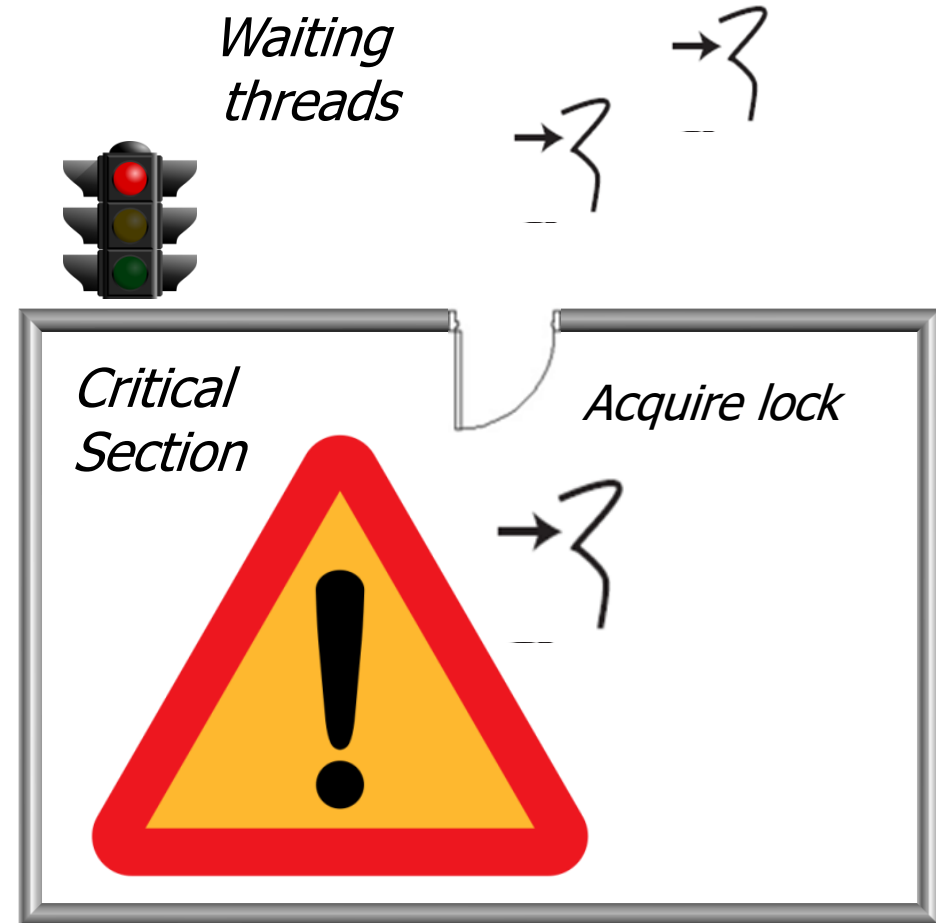
Overview of Mutual Exclusion Locks

- A mutual exclusion lock (mutex) defines a “critical section”
 - Ensures only one thread can run in a block of code at a time
 - Other threads are kept “at bay”
 - After a thread leaves a critical section another thread can enter & start running



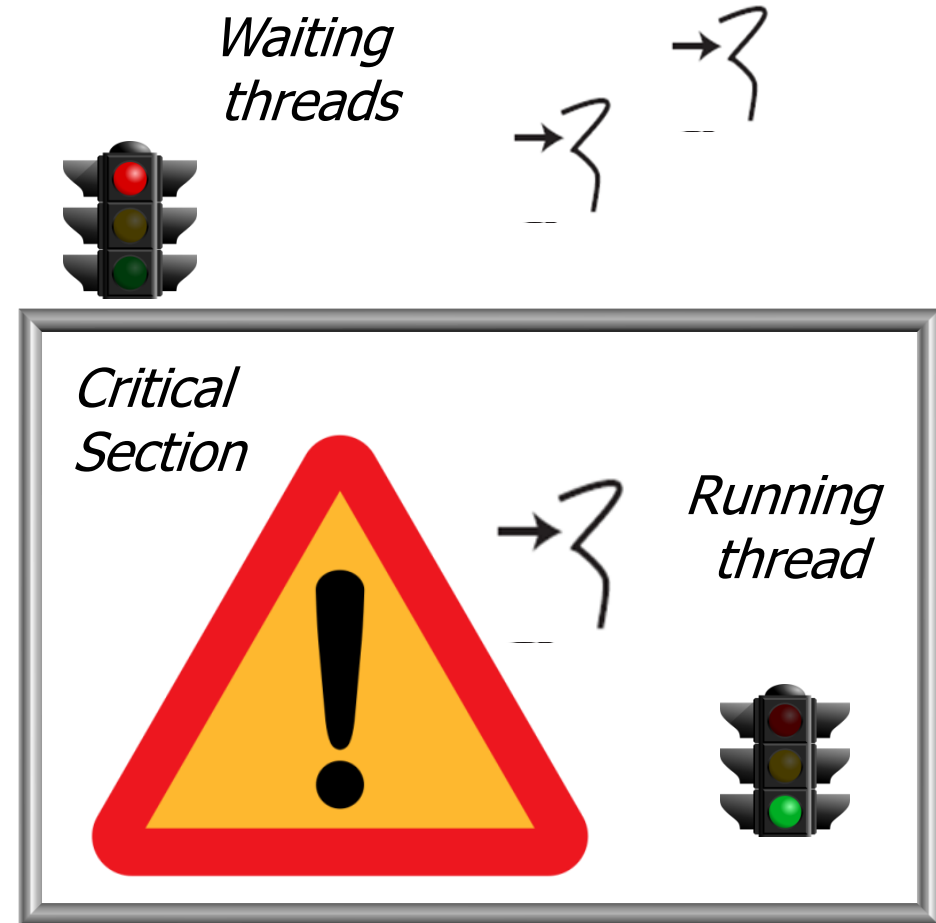
Overview of Mutual Exclusion Locks

- A mutual exclusion lock (mutex) defines a “critical section”
 - Ensures only one thread can run in a block of code at a time
 - Other threads are kept “at bay”
 - After a thread leaves a critical section another thread can enter & start running



Overview of Mutual Exclusion Locks

- A mutual exclusion lock (mutex) defines a “critical section”
 - Ensures only one thread can run in a block of code at a time
 - Other threads are kept “at bay”
 - After a thread leaves a critical section another thread can enter & start running



Overview of Mutual Exclusion Locks

- A mutex is typically implemented in hardware via atomic operations

Atomic operations appear to occur instantaneously & either change the state of the system successful or have no effect



See en.wikipedia.org/wiki/Linearizability

Overview of Mutual Exclusion Locks

- A mutex is typically implemented in hardware via atomic operations
- Implemented in Java via the `compareAndSwap*()` methods in the `Unsafe` class

Concurrency

And few words about concurrency with `Unsafe`. `compareAndSwap` methods are atomic and can be used to implement high-performance lock-free data structures.

For example, consider the problem to increment value in the shared object using lot of threads.

First we define simple interface `Counter`:

```
interface Counter {  
    void increment();  
    long getCounter();  
}
```

Then we define worker thread `CounterClient`, that uses `Counter`:

```
class CounterClient implements Runnable {  
    private Counter c;  
    private int num;  
  
    public CounterClient(Counter c, int num) {  
        this.c = c;  
        this.num = num;  
    }  
  
    @Override  
    public void run() {  
        for (int i = 0; i < num; i++) {  
            c.increment();  
        }  
    }  
}
```

See earlier discussion of "*Java Atomic Classes & Operations*"

Human Known Use of Mutual Exclusion Locks

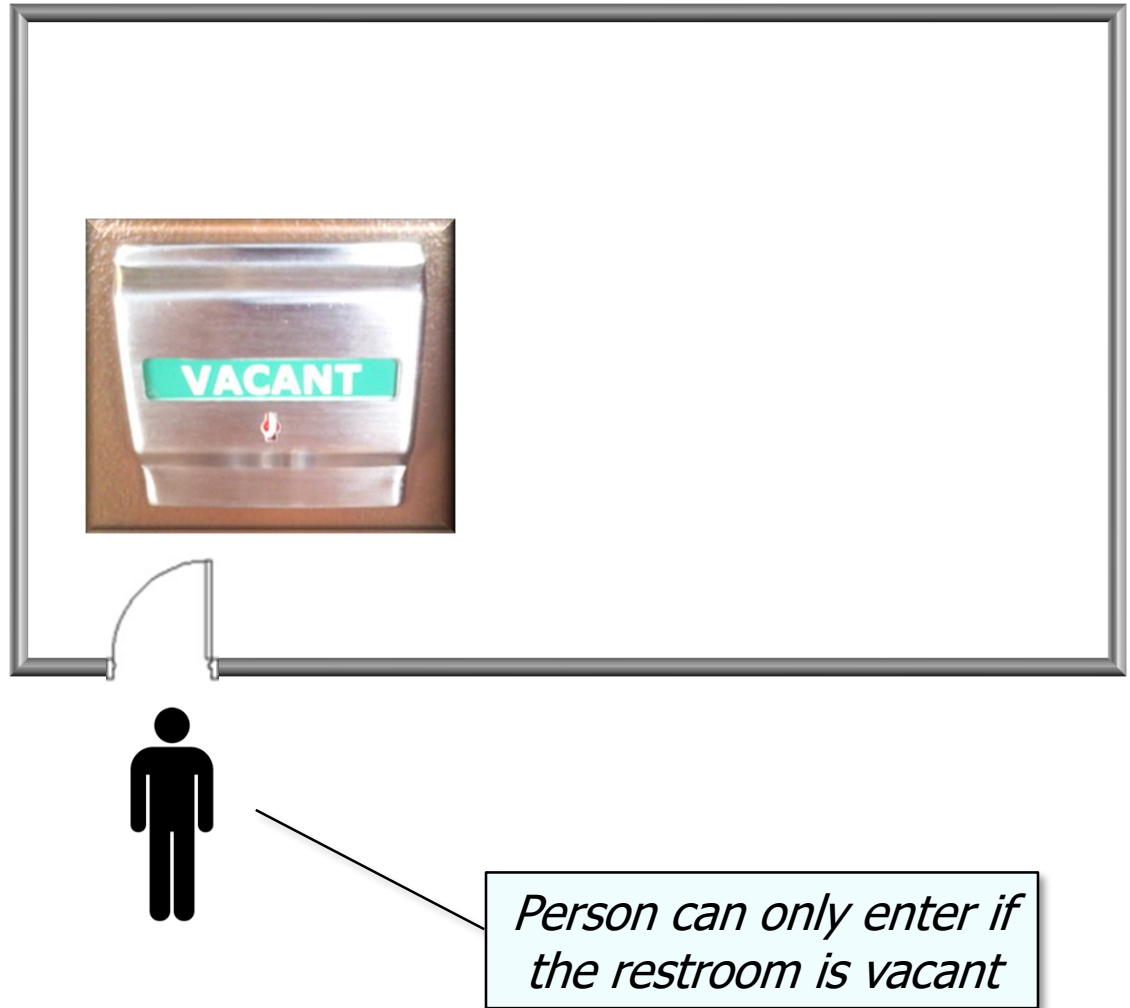
Human Known Use of Mutual Exclusion Locks

- A human known use of mutual exclusion locks is an airplane restroom protocol



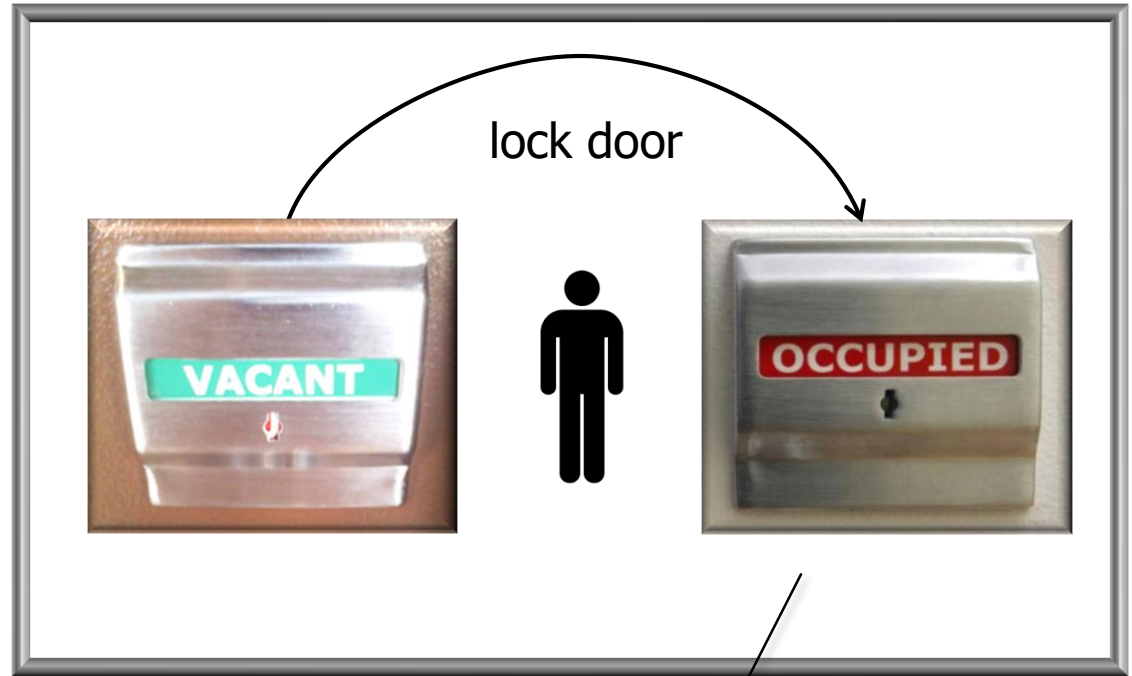
Human Known Use of Mutual Exclusion Locks

- A human known use of mutual exclusion locks is an airplane restroom protocol



Human Known Use of Mutual Exclusion Locks

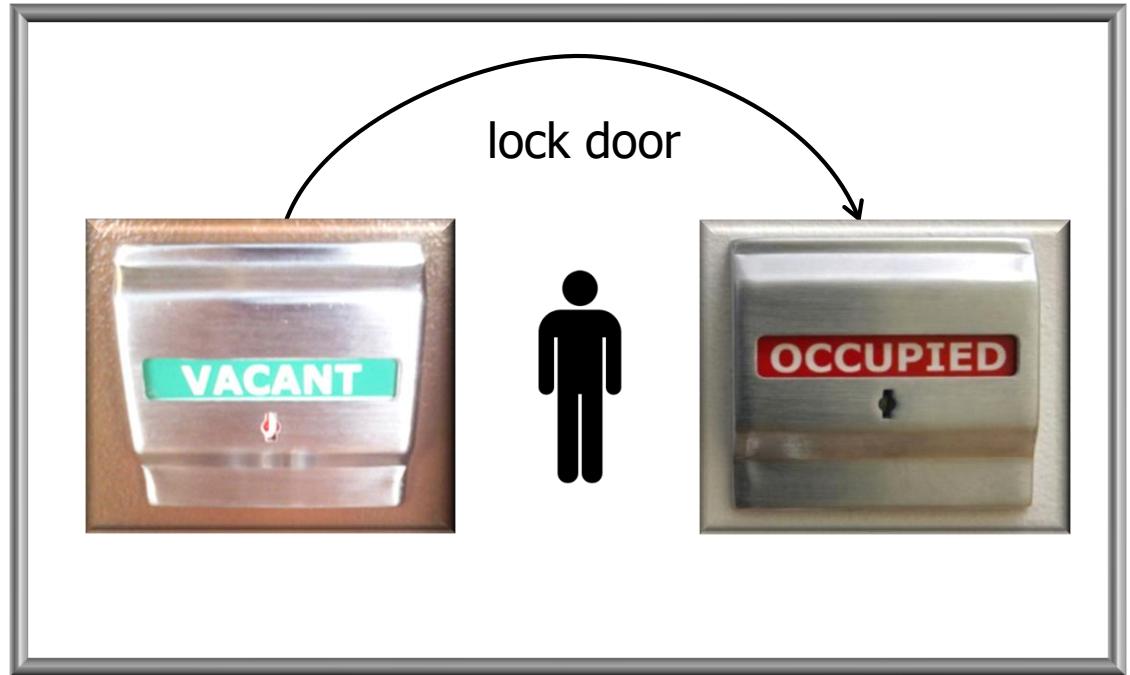
- A human known use of mutual exclusion locks is an airplane restroom protocol



*Person atomically
enters & locks the door*

Human Known Use of Mutual Exclusion Locks

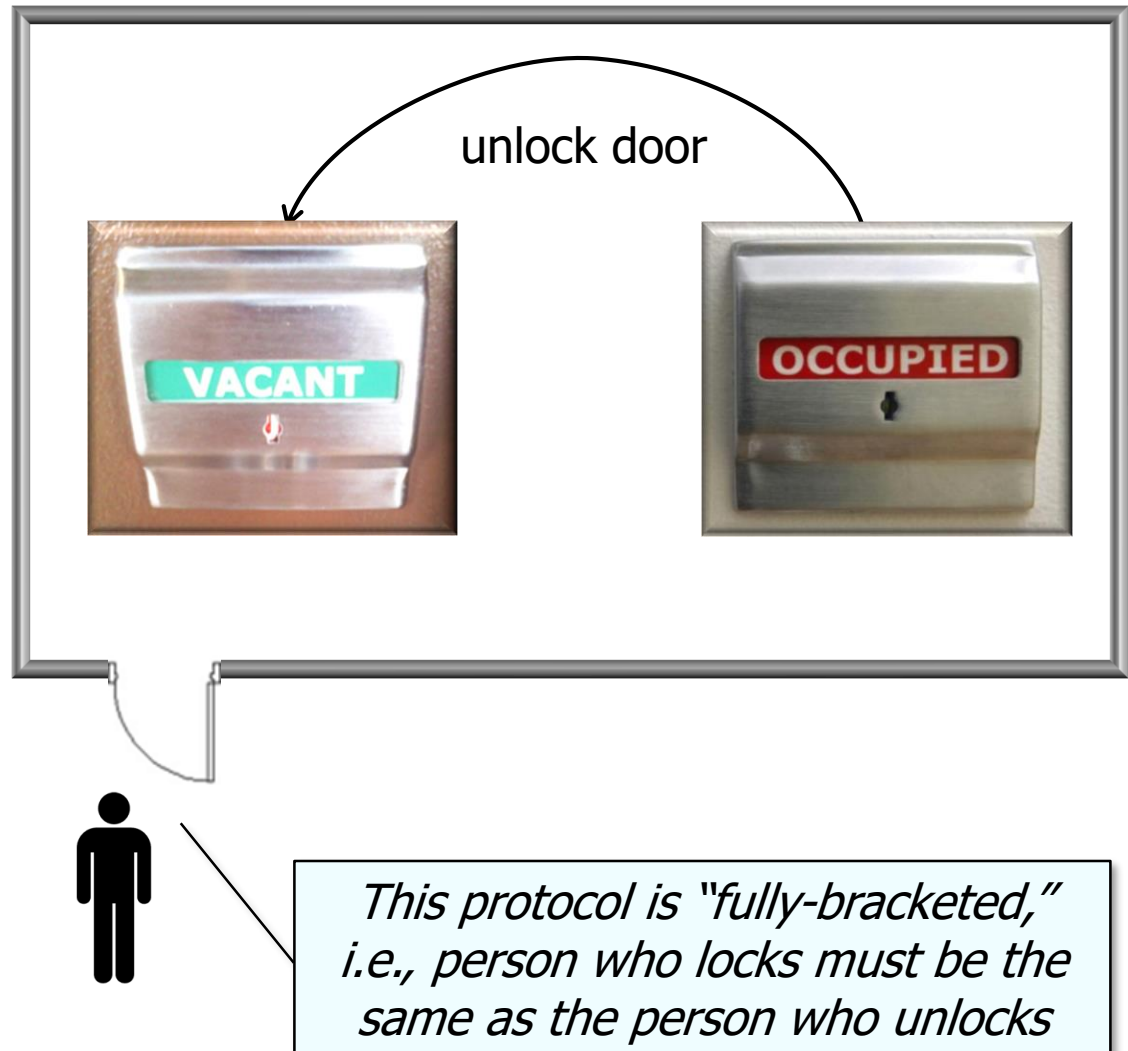
- A human known use of mutual exclusion locks is an airplane restroom protocol



Other people who want to use the restroom must wait while it's in use

Human Known Use of Mutual Exclusion Locks

- A human known use of mutual exclusion locks is an airplane restroom protocol



Human Known Use of Mutual Exclusion Locks

- A human known use of mutual exclusion locks is an airplane restroom protocol



*Once the restroom is vacant
another person can enter*

End of Java ReentrantLock (Part 1)