Java StampedLock: Usage Considerations



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Learning Objectives in this Part of the Lesson

- Understand the structure, functionality of the Java StampedLock class
- Know the key methods in Java StampedLock
- Recognize how to apply Java StampedLock in practice
- Appreciate Java StampedLock usage considerations



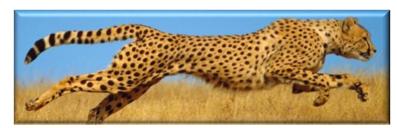
Java StampedLock Usage Considerations

StampedLock often much faster than ReentrantReadWriteLock

SYNCHRONIZED	OPTIMISTIC	RWLOCK	STAMPED
1996.6	1174	116393	64077
2312.7	1174	116617	47897
2100.9	1122	117746	65921
2285.1	1182.9	115605	73500
2173.6	1184.9	118346	32857
2173.78	1167.56	116941.4	56850.4

19 readers & 1 writer

Optimistic read mode works very well with little/no contention



See www.takipiblog.com/java-8-stampedlocks-vs-readwritelocks-and-synchronized

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ReentrantReadWriteLock is very slow...

StampedLock often much faster than ReentrantReadWriteLock

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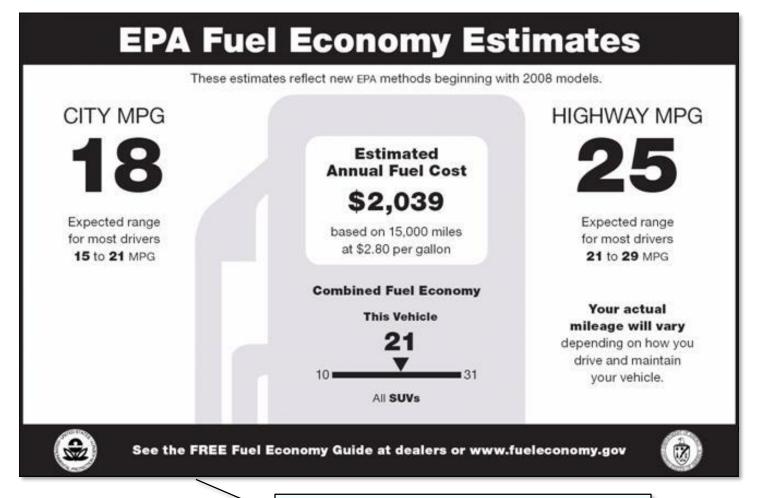
StampedLock with "reading mode" works better than ReentrantReadWriteLock

StampedLock often much faster than ReentrantReadWriteLock

SYNCHRONIZED	OPTIMISTIC	RWLOCK	STAMPED
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2312.7	1174	116617	47897
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Synchronized statements perform quite well

StampedLock often much faster than ReentrantReadWriteLock



However, your mileage may vary!

See en.wiktionary.org/wiki/your mileage may vary

StampedLock often much faster than ReentrantReadWriteLock

RWLOCK	STAMPED	SYNCHRONIZED	OPTIMISTIC
1960.8	165.1	177.4	387.9
1473.6	111.3	192.1	382.8
2119.7	216.8	173.3	403.6
2772.2	221.9	205.4	403.9
2721.4	189.3	181.2	394.2
2209.54	180.88	185.88	394.48

10 readers & 10 writers

Optimistic read mode works less well with more contention

See www.takipiblog.com/java-8-stampedlocks-vs-readwritelocks-and-synchronized

StampedLock often much faster than ReentrantReadWriteLock

RWLOCK	STAMPED	SYNCHRONIZED	OPTIMISTIC
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10 readers & 10 writers

However, ReentrantReadWriteLock is still much slower..



StampedLock often much faster than ReentrantReadWriteLock

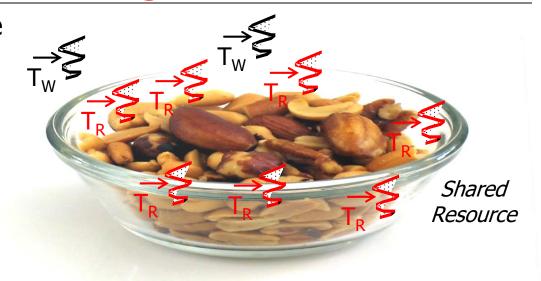
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221.9	205.4	403.9		
189.3	181.2	394.2		
180.88	185.88	394.48		
10 readers & 10 writers				
	216.8 221.9 189.3 180.88	216.8 173.3 221.9 205.4 189.3 181.2 180.88 185.88		

StampedLock & synchronized statements both do quite well

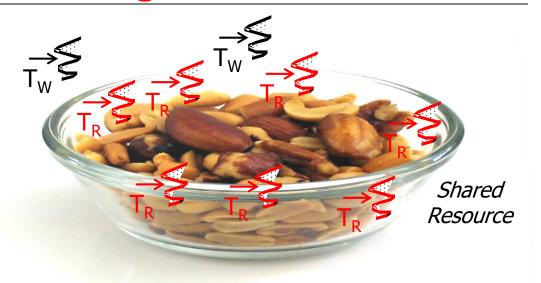
 Java StampedLock speedups are only fully realized under certain conditions



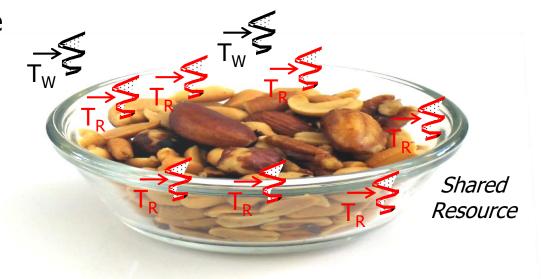
- Java StampedLock speedups are only fully realized under certain conditions, e.g.
 - Frequency of reads to writes
 - Ideally, many more reads than writes



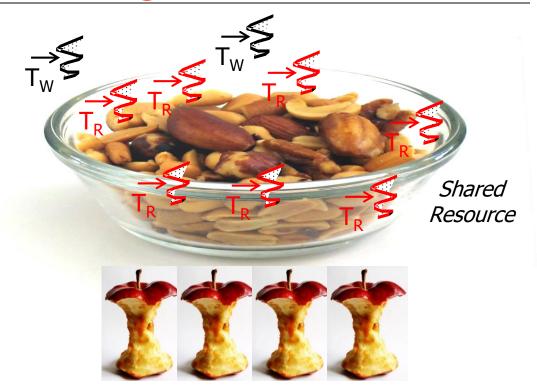
- Java StampedLock speedups are only fully realized under certain conditions, e.g.
 - Frequency of reads to writes
 - Duration of read & write operations
 - Ideally, read operations should be non-trivial or else locking costs may dominate



- Java StampedLock speedups are only fully realized under certain conditions, e.g.
 - Frequency of reads to writes
 - Duration of read & write operations
 - "Contention" for the data
 - Ideally, *many* concurrent readers



- Java StampedLock speedups are only fully realized under certain conditions, e.g.
 - Frequency of reads to writes
 - Duration of read & write operations
 - "Contention" for the data
 - Number of processor cores
 - Ideally, *many* cores



 StampedLock can be harder to use than ReentrantReadWriteLock



- StampedLock can be harder to use than ReentrantReadWriteLock
 - Many more methods

© ReentrantReadWriteLock() © ReentrantReadWriteLock() © ReentrantReadWriteLock(boolean) ● writeLock():WriteLock ● readLock():ReadLock © isFair():boolean ● getReadLockCount():int ● isWriteLocked():boolean ● isWriteLockedByCurrentThread():boolean ● getWriteHoldCount():int ● getReadHoldCount():int © fhasQueuedThreads():boolean © fhasQueuedThread(Thread):boolean © getQueueLength():int ● hasWaiters(Condition):boolean

<<Java Class>>

<<Java Class>> StampedLock StampedLock() writeLock():long tryWriteLock():long tryWriteLock(long,TimeUnit):long writeLockInterruptibly():long readLock():long tryReadLock():long tryReadLock(long,TimeUnit):long readLockInterruptibly():long tryOptimisticRead():long validate(long):boolean unlockWrite(long):void unlockRead(long):void unlock(long):void tryConvertToWriteLock(long):long tryConvertToReadLock(long):long tryConvertToOptimisticRead(long):long tryUnlockWrite():boolean tryUnlockRead():boolean isWriteLocked():boolean isReadLocked():boolean getReadLockCount():int toString() asReadLock():Lock asWriteLock():Lock

asReadWriteLock():ReadWriteLock

getWaitQueueLength(Condition):int

toString()

- StampedLock can be harder to use than ReentrantReadWriteLock
 - Many more methods
 - More intricate semantics& usage patterns

Conditional writes & lock upgrades are tricky to program

```
void moveIfAtOrigin(double newX,
                     double newY) {
  long stamp = sl.readLock();
    try {
      while (x == 0.0 \&\& y == 0.0) {
        long ws =
          sl.tryConvertToWriteLock
                (stamp);
        if (ws != 0L) {
          stamp = ws;
          x = newX; y = newY;
          break:
        } else {
           sl.unlockRead(stamp);
           stamp = sl.writeLock();
       finally
       sl.unlock(stamp); }
```

- StampedLock can be harder to use than ReentrantReadWriteLock
 - Many more methods
 - More intricate semantics
 & usage patterns
 - Invariants are tricky with optimistic read locks

```
class Boooom {
  StampedLock mS =
    new StampedLock();
  int mX = 0;
  int mY = 1;
     Thread T1
  while (true) {
    mS.writeLock();
    mX++; mY++;
    mS.writeUnlock();
  }
     Thread T2
  do {
    stamp = mS.tryOptimisticRead();
    z = 1 / (mX - mY);
  } while (mS.validate(stamp));
```

- StampedLock can be harder to use than ReentrantReadWriteLock
 - Many more methods
 - More intricate semantics& usage patterns
 - Invariants are tricky with optimistic read locks

```
class Boooom {
  StampedLock mS =
    new StampedLock();
  int mX = 0;
  int mY = 1;
                  Create a StampedLock
                   to protect two fields
     Thread T1
  while (true) {
    mS.writeLock();
    mX++; mY++;
    mS.writeUnlock();
  }
     Thread T2
  do {
    stamp = mS.tryOptimisticRead();
    z = 1 / (mX - mY);
  } while (mS.validate(stamp));
```

- StampedLock can be harder to use than ReentrantReadWriteLock
 - Many more methods
 - More intricate semantics
 & usage patterns
 - Invariants are tricky with optimistic read locks
 - Fields read in optimistic mode may be inconsistent since their values can change unpredictably

```
class Boooom {
  StampedLock mS =
    new StampedLock();
  int mX = 0;
  int mY = 1;
                  Want to establish the
                 invariant mX == mY - 1
     Thread T1
  while (true)
    mS.writeLock();
    mX++; mY++;
    mS.writeUnlock();
     Thread T2
  do {
    stamp = mS.tryOptimisticRead();
    z = 1 / (mX - mY);
  } while (mS.validate(stamp));
```

- StampedLock can be harder to use than ReentrantReadWriteLock
 - Many more methods
 - More intricate semantics
 & usage patterns
 - Invariants are tricky with optimistic read locks
 - Fields read in optimistic mode may be inconsistent since their values can change unpredictably

Since no read lock is held, mX & mY may be reordered, such that invariant mX == mY - 1 may not hold

```
class Boooom {
  StampedLock mS =
    new StampedLock();
  int mX = 0;
  int mY = 1;
     Thread T1
  while (true) {
    mS.writeLock();
    mX++; mY++;
    mS.writeUnlock();
     Thread T2
  do {
    stamp = mS.tryOptimisticRead();
    z = 1 / (mX - mY);
    while (mS.validate(stamp));
```

- StampedLock can be harder to use than ReentrantReadWriteLock
 - Many more methods
 - More intricate semantics
 & usage patterns
 - Invariants are tricky with optimistic read locks
 - Non-reentrant

```
class SomeComponent {
 private StampedLock sl =
   new StampedLock();
 public oid someMethod1() {
    long stam = sl.readLock();
    someMethod2
 private void someMethod2()
    long stamp = sl.readLock()
```

- StampedLock is usually the best choice for readers-writer locks in Java 8+!
 - Despite its complexity & lack of reentrant semantics

Class StampedLock

java.lang.Object java.util.concurrent.locks.StampedLock

All Implemented Interfaces:

Serializable

public class StampedLock
extends Object
implements Serializable

A capability-based lock with three modes for controlling read/write access. The state of a StampedLock consists of a version and mode. Lock acquisition methods return a stamp that represents and controls access with respect to a lock state; "try" versions of these methods may instead return the special value zero to represent failure to acquire access. Lock release and conversion methods require stamps as arguments, and fail if they do not match the state of the lock. The three modes are:

- Writing. Method writeLock() possibly blocks waiting for exclusive access, returning a stamp that can be used in method unlockWrite(long) to release the lock. Untimed and timed versions of tryWriteLock are also provided. When the lock is held in write mode, no read locks may be obtained, and all optimistic read validations will fail.
- Reading. Method readLock() possibly blocks waiting for non-exclusive access, returning a stamp that can be used in method unlockRead(long) to release the lock. Untimed and timed versions of tryReadLock are also provided.
- Optimistic Reading. Method tryOptimisticRead() returns a non-zero stamp only if the lock is not currently held in write mode. Method validate(long) returns true if the lock has not been acquired in write mode since obtaining a given stamp. This mode can be thought of as an extremely weak version of a read-lock, that can be broken by a writer at any time. The use of optimistic mode for short read-only code segments

See www.javaspecialists.eu/archive/Issue215.html

End of Java StampedLock: Usage Considerations