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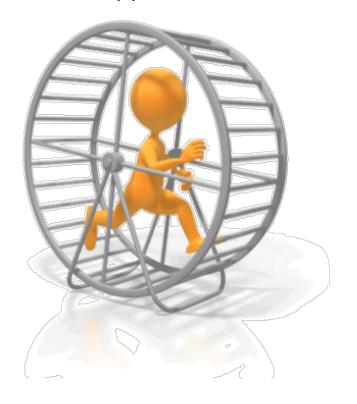


Learning Objectives in this Part of the Lesson

- Understand how Java volatile variables provide concurrent programs with thread-safe mechanisms to read from & write to single variables
- Know how to use a Java volatile variable in practice
- Appreciate usage considerations for Java volatile variables



Concurrent apps should use volatile variables carefully to avoid "busy waiting"



```
class LoopMayNeverEnd {
  volatile boolean mDone = false;
  void work() {
    // Thread T2 read
    while (!mDone) {
      // do work \
         If "do work" isn't time consuming
           this loop will spin excessively..
  void stopWork() {
    // Thread T1 write
    mDone = true;
```

- Concurrent apps should use volatile variables carefully to avoid "busy waiting"
 - Busy waiting is most effective when encapsulated in higherlevel concurrency libraries



"Engineering Concurrent Library Components"

Doug Lea

Day 2 - April 3, 2013 - 1:30 PM - Salon C

phillyemergingtech.com

```
public class AtomicLong
  private volatile long value;
  private static final Unsafe unsafe
    = Unsafe.getUnsafe();
  private static final long
    valueOffset;
  static {
    valueOffset = unsafe.
      objectFieldOffset
        (AtomicLong
        .class
        .getDeclaredField("value"));
```

 Complex operations that perform multiple instructions can't use volatile by itself

```
volatile int counter = 0;
// In Thread t1
counter++;
// load counter into register r1
// increment register r1
// store register r1 into counter
// In Thread t2
counter--;
// load counter into register r1
// decrement register r1
// store register r1 into counter
```

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 - Incrementing an integer

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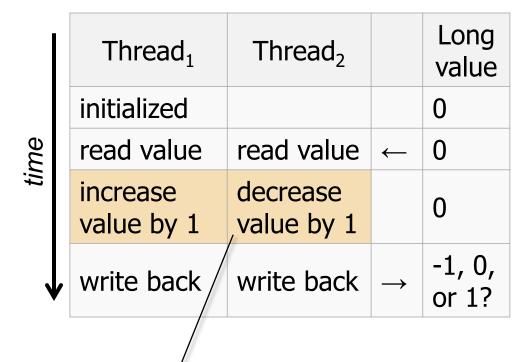
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If these steps interleave in multiple threads the results may be inconsistent

See en.wikipedia.org/wiki/Write-write_conflict

- Complex operations that perform multiple instructions can't use volatile by itself, e.g.
 - Incrementing an integer
 - Use an atomic variable instead of a volatile variable

```
AtomicLong mCounter =
  new AtomicLong(0);
// In Thread t1
mCounter.getAndIncrement();
// load counter into register r1
// increment register r1
// store register r1 into counter
// In Thread t2
mCounter.getAndDecrement();
// load counter into register r1
// decrement register r1
// store register r1 into counter
```

 Declaring an array or an object as volatile only makes the reference volatile

```
public class Vector<E> ... {
  /**
   * The number of elements or
   * the size of the vector.
   */
  protected int elementCount;
  /**
   * The elements of the vector.
   */
  protected Object[] elementData;
volatile Vector v = new Vector();
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Volatile variable

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 - However, the contents pointed to by the reference are not volatile

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Non-volatile fields
```

- Declaring an array or an object as volatile only makes the reference volatile
 - However, the contents pointed to by the reference are not volatile
 - Therefore, more powerful types of synchronization are needed

```
public class Vector<E> ... {
  public synchronized E set
       (int location, E object) {
    if (location < elementCount) {</pre>
      E result = (E)
        elementData[location];
       elementData[location] =
         object;
      return result:
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```

 Java semantics of volatile aren't the same as in C or C++

```
In C and C++ redit
In C, and consequently C++, the volatile keyword was intended to [1]

    allow access to memory mapped devices

    allow uses of variables between setjmp and longjmp

    allow uses of sig atomic t variables in signal handlers.

Operations on volatile variables are not atomic, nor do they establish a
proper happens-before relationship for threading. This is according to the
relevant standards (C, C++, POSIX, WIN32),[2] and this is the matter of fact for
the vast majority of current implementations. Thus, the usage of volatile
keyword as a portable synchronization mechanism is discouraged by many
C/C++ groups.[3][4][5]
Example of memory-mapped I/O in C [edit]
In this example, the code sets the value stored in foo to 0. It then starts to
poll that value repeatedly until it changes to 255:
  static int foo;
  void bar(void) {
       foo = 0;
       while (foo != 255)
```

An optimizing compiler will notice that no other code can possibly change the value stored in foo, and will assume that it will remain equal to of at all times. The compiler will therefore replace the function body with an infinite loop similar

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- Java semantics of volatile aren't the same as in C or C++
 - Volatiles in C/C++ aren't atomic & don't create a happens-before relationship
 - They largely just disable compiler optimizations



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End of Usage Considerations of Java Volatile Variables