Introduction to the Specific Notification Pattern



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Learning Objectives in this Part of the Lesson

Understand the *Specific Notification* pattern

Specific Notification for Java Thread Synchronization

Tom Cargill

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Abstract

Java supports thread synchronization by means of monitorlike primitives. The weak semantics of Java's signaling mechanism provides little control over the order in which threads acquire resources, which encourages the use of the Haphazard Notification pattern, in which an *arbitrary* thread is selected from a set of threads competing for a resource. For synchronization problems in which such arbitrary selection of threads is unacceptable, the Specific Notification pattern may be used to designate exactly which thread should proceed. Specific Notification provides an explicit mechanism for thread selection and scheduling.

0. Introduction

To study Java's threads, I first tackled some of the classic exercises, like the "Dining Philosophers" and the "Readers and Writers." The solutions that I obtained were reasonable, but I felt uncomfortable with the degree to which I had to depend on serendipitous treatment with respect to contention for locks and notifications. The solutions were free of deadlock, but were not fair in all circumstances. I thought I might have to resign myself to tolerating some unfairness in Java. Next, I built a multithreaded NNTP¹ client, in which several

¹ B. Kantor, P. Lapsley, Network News Transfer Protocol, Internic RFC 977, 1986. threads could have active requests outstanding with an NNTP server. The fundamental correctness of this class depended on waiting threads being reactivated in *exactly* the right order to receive their responses from the server. In coding this class I applied the Specific Notification mechanism described below. With new insight, I returned to the earlier exercises and found that Specific Notification provided complete solutions to those problems. I therefore propose the Specific Notification pattern.

Section 1 summarizes the semantics of Java's thread synchronization mechanisms, contrasting them with classical monitors; this section may be omitted by readers who have a detailed

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Context

- A family of threads in Java that need to cooperate by synchronizing access to shared resources in a specific way
 - e.g., FIFO, LIFO, priority, etc.



Problem

 Java built-in monitor objects provide apps no control over the order in which threads acquire a resource

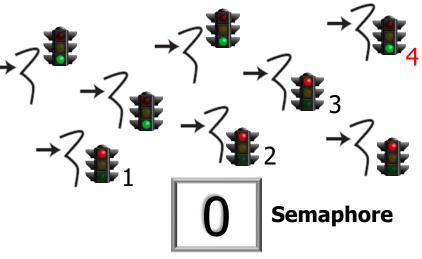




a.k.a., "Haphazard Notification"

• Problem

- Java built-in monitor objects provide apps no control over the order in which threads acquire a resource, e.g.
 - Selection of thread(s) after notify() or notifyAll()



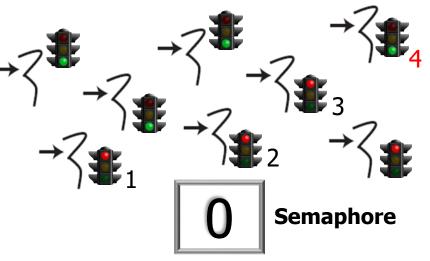




• Problem

- Java built-in monitor objects provide apps no control over the order in which threads acquire a resource, e.g.
 - Selection of thread(s) after notify() or notifyAll()
 - Scheduling of thread(s) after notify() or notifyAll()







Solution – Apply the Specific Notification pattern

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See <u>www.dre.vanderbilt.edu/~schmidt/PDF/</u> <u>specific-notification.pdf</u> (especially Listing 3)

Solution – Apply the Specific Notification pattern

 Provide a non-haphazard mechanism for selecting/scheduling threads Specific Notification for Java Thread Synchronization

Tom Cargill Consultant Box 69, Louisville, CO 80027 www.sni.net/~cargill

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Solution – Apply the Specific Notification pattern

- Provide a non-haphazard mechanism for selecting/scheduling threads
- Designate *exactly* which thread in a family of threads should proceed next after notify() or notifyAll()

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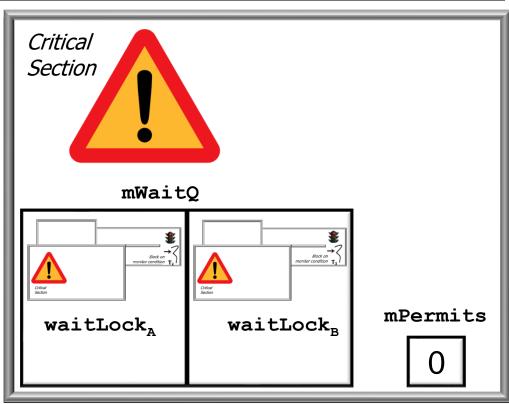
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Solution Outline

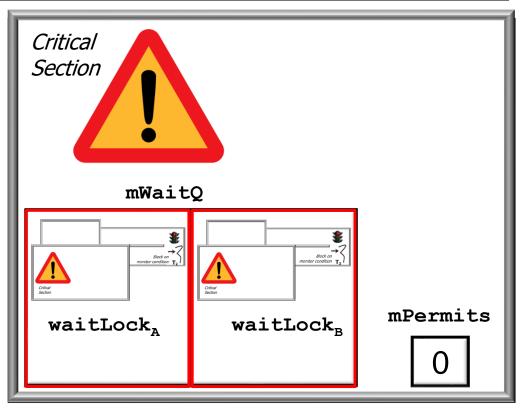
 Put threads to sleep via wait() calls in monitor objects



Can also use ReentrantLock & ConditionObject

Solution Outline

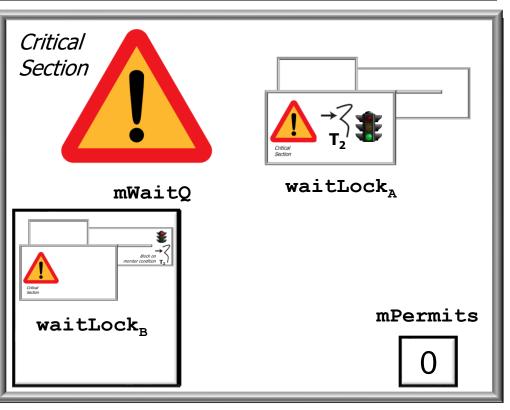
- Put threads to sleep via wait() calls in monitor objects
 - One monitor object is used for each thread that must be individually notified



Solution Outline

- Put threads to sleep via wait() calls in monitor objects
- One waiting thread is notified in a specific order so it can run next
 - e.g., FIFO, LIFO, priority, etc.





End of Introduction to the Specific Notification Pattern