Java Monitor Objects: Usage Considerations



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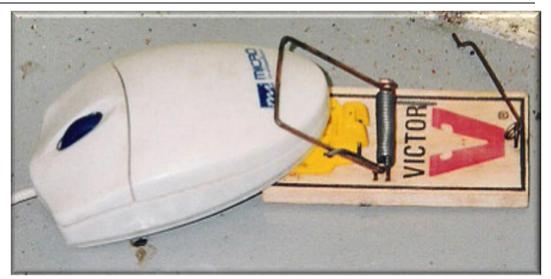
Learning Objectives in this Lesson

 Appreciate Java monitor object usage considerations



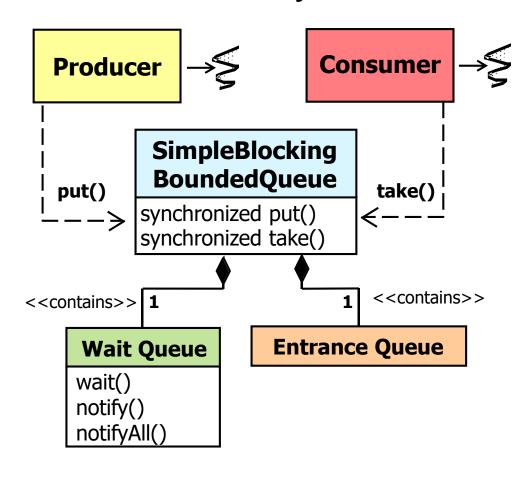
Learning Objectives in this Lesson

- Appreciate Java monitor object usage considerations
 - In particular, know common traps & pitfalls of Java's builtin monitor objects



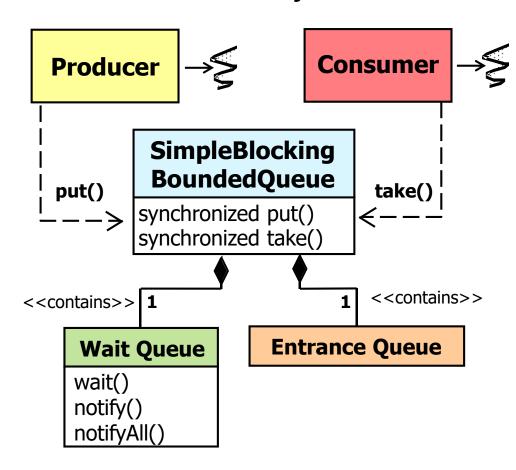
Programmers must be aware of issues with Java monitor objects



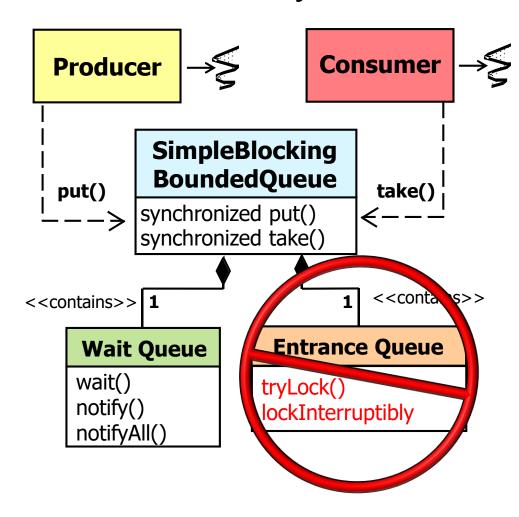


- Programmers must be aware of issues with Java monitor objects
 - Monitor objects are limited





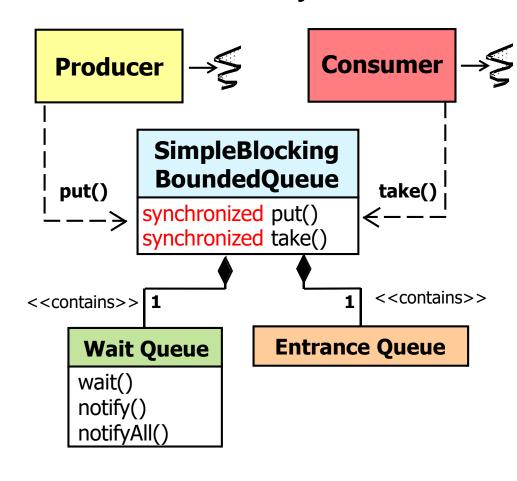
- Programmers must be aware of issues with Java monitor objects
 - Monitor objects are limited, e.g.
 - No non-blocking, timed, or interruptible synchronizers



See lessons on "Java ReentrantLocks" for examples of these capabilities

- Programmers must be aware of issues with Java monitor objects
 - Monitor objects are limited, e.g.
 - No non-blocking, timed, or interruptible synchronizers
 - Only one wait queue & one entrance queue





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 - No non-blocking, timed, or interruptible synchronizers
 - Only one wait queue & one entrance queue
 - May yield "nested monitor lockout"

```
public class BuggyLock {
  Object mMonObj = new Object();
  boolean mLocked;
  synchronized void lock() {
    while (mLocked)
      synchronized(mMonObj)
      { mMonObj.wait(); }
    mLocked = true;
lock() is a synchronized method
  synchronized void unlock() {
    mLocked = false:
    synchronized(mMonObj)
    { mMonObj.notify(); }
```

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  synchronized void lock() {
    while (mLocked)
      synchronized(mMonObj)
      { mMonObj.wait();/}
    mLocked = true;
       BuggyLock monitor lock is still
     held here, so unlock() never runs!
  synchronized void unlock() {
    mLocked = false;
    synchronized(mMonObj)
    { mMonObj.notify(); }
```

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 - Monitor objects are limited, e.g.
 - No non-blocking, timed, or interruptible synchronizers
 - Only one wait queue & one entrance queue
 - May yield "nested monitor lockout"
 - Doesn't support "two lock queue" optimizations

```
class LinkedBlockingQueue<E>
      extends AbstractQueue<E>
     implements BlockingQueue<E>,
  /** Lock held by take, poll,
      etc */
  private final ReentrantLock
    takeLock =
      new ReentrantLock();
  /** Lock held by put, offer,
      etc */
  private final ReentrantLock
    putLock =
       new ReentrantLock();
```

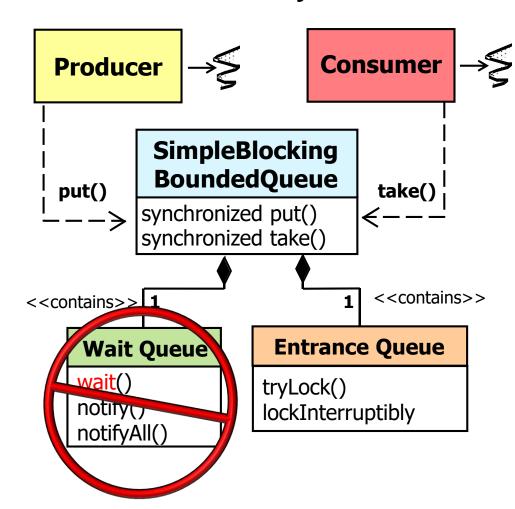
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// end of this block

```
Producer
                           Consumer
            SimpleBlocking
            BoundedQueue
                                 take()
  put()
          synchronized put()
          synchronized take()
                                 <<contains>>
<<contains>> 1
                       Entrance Queue
     Wait Queue
    wait()
    notify()
    notifyAll()
```

Scoped locking is inefficient for certain concurrent algorithms, e.g., it may require redundant checks for internal state(s)

- Programmers must be aware of issues with Java monitor objects
 - Monitor objects are limited, e.g.
 - No non-blocking, timed, or interruptible synchronizers
 - Only one wait queue & one entrance queue
 - Synchronized statements only support scoped locking
 - No support for sensible timed waits...



See <u>stackoverflow.com/questions/3397722/how-to-differentiate-when-waitlong-timeout-exit-for-notify-or-timeout</u>

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 - Monitor objects are limited
 - Choosing between notify()
 & notifyAll() is tricky

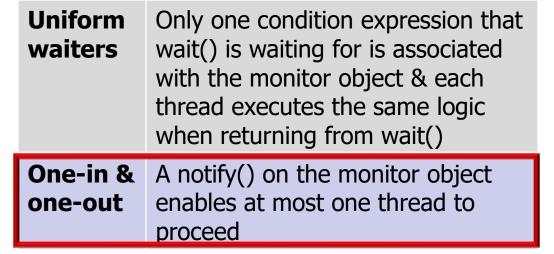


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Uniform waiters	Only one condition expression that wait() is waiting for is associated with the monitor object & each thread executes the same logic when returning from wait()
One-in & one-out	A notify() on the monitor object enables at most one thread to proceed

Conditions under which notify() can be used

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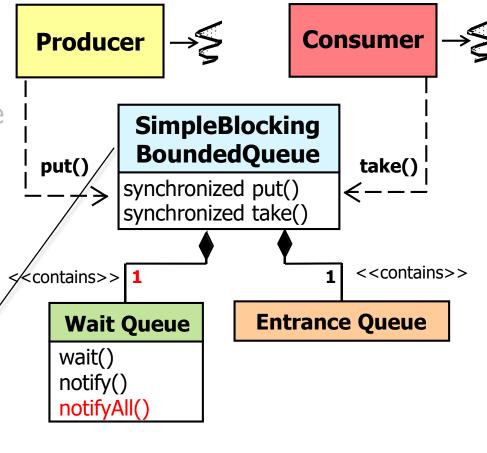
Conditions under which notify() can be used

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 - Choosing between notify()
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 - Use notify() when possible since it's more efficient & avoids the "Thundering Herd" problem..

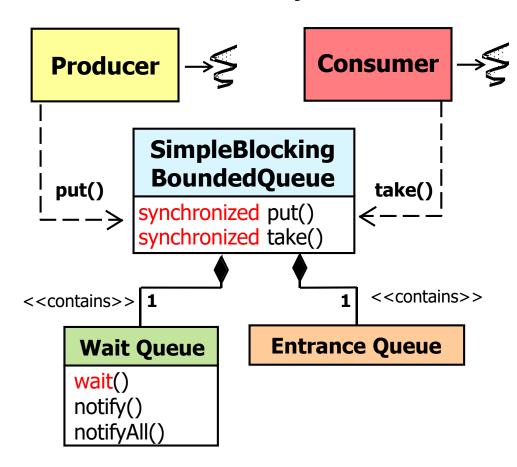


- Programmers must be aware of issues with Java monitor objects
 - Monitor objects are limited
 - Choosing between notify()
 & notifyAll() is tricky
 - Use notify() when possible since it's more efficient & avoids the "Thundering Herd" problem..
 - However, notifyAll() is often needed since there's just one wait queue..

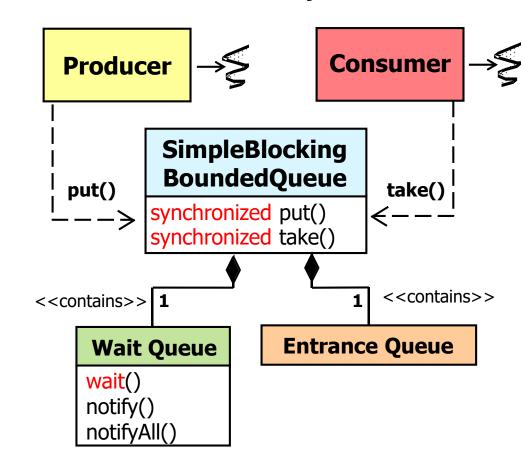
A monitor object may need to wait for different condition expression



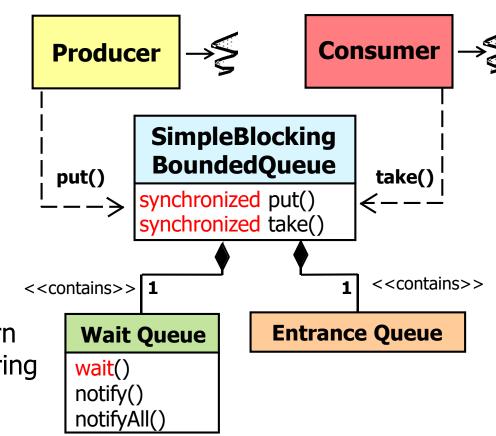
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 - Fairness issues arise due to the order in which waiting threads are notified



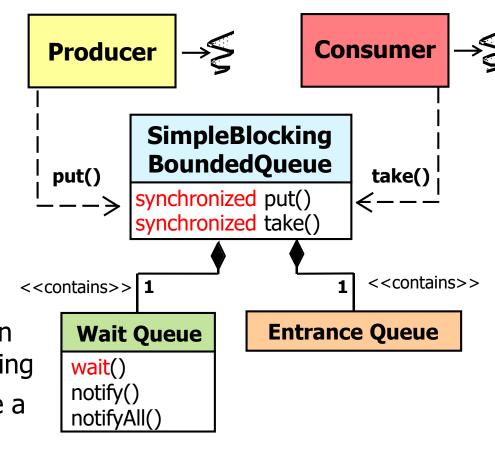
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 - Monitor object's implement "haphazard notification" to optimize performance



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 - The Specific Notification pattern can be applied to control ordering



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 - Monitor object's implement "haphazard notification" to optimize performance
 - The Specific Notification pattern can be applied to control ordering
 - i.e., programmatically choose a particular thread to run from a family of waiting threads



- In practice, you often need more than Java's Java monitor mechanisms
 - java.util.concurrent & java.util.concurrent.locks

package Added in API level 1

java.util.concurrent.locks

Interfaces and classes providing a framework for locking and waiting for conditions that is distinct from built-in synchronization and monitors. The framework permits much greater flexibility in the use of locks and conditions, at the expense of more awkward syntax.

The Lock interface supports locking disciplines that differ in semantics (reentrant, fair, etc), and that can be used in non-block-structured contexts including hand-over-hand and lock reordering algorithms. The main implementation is ReentrantLock.

package Added in API level 1

java.util.concurrent

Utility classes commonly useful in concurrent programming. This package includes a few small standardized extensible frameworks, as well as some classes that provide useful functionality and are otherwise tedious or difficult to implement. Here are brief descriptions of the main components. See also the java.util.concurrent.locks and java.util.concurrent.atomic packages.

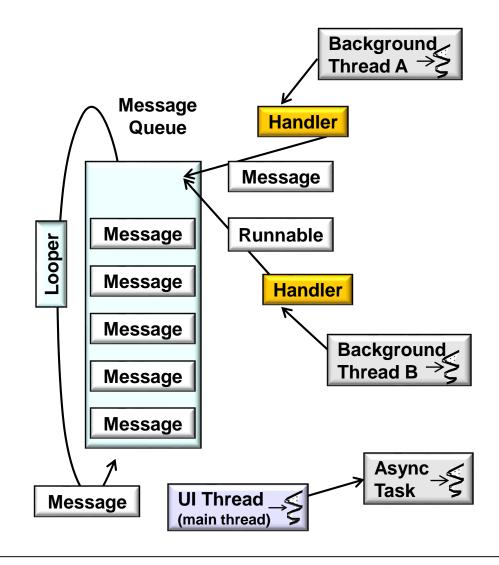
See developer.android.com/reference/java/util/concurrent/package-summary.html

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 - java.util.concurrent & java.util.concurrent.locks
 - e.g., ReentrantLock & ConditionObject

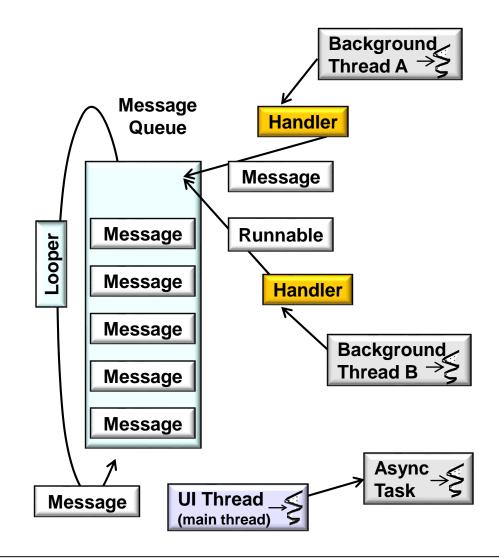
Used to protect the object state from race conditions

```
public class ArrayBlockingQueue<E>
           extends AbstractQueue<E>
        implements BlockingQueue<E>,
              java.io.Serializable {
  /** Main lock guarding access */
  final ReentrantLock lock;
  /** Condition for waiting takes */
  private final Condition notEmpty;
  /** Condition for waiting puts */
  private final Condition notFull;
```

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 - Android concurrency frameworks



- In practice, you often need more than Java's Java monitor mechanisms
 - java.util.concurrent & java.util.concurrent.locks
 - Android concurrency frameworks
 - Message passing may avoid need for monitor objects & synchronization altogether



End of Java Monitor Objects: Usage Considerations