

Java ExecutorCompletionService: Designing a Memoizer

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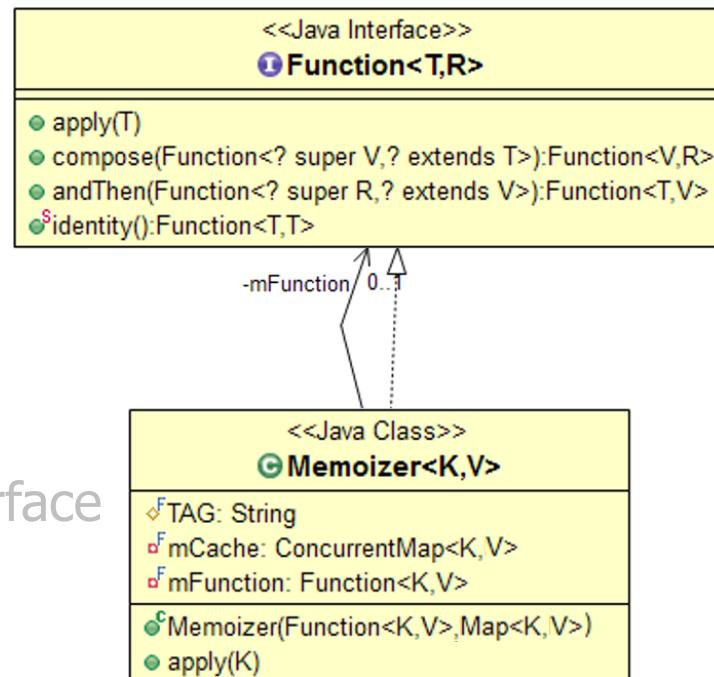
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Nashville, Tennessee, USA**



Learning Objectives in this Part of the Lesson

- Understand how the Java CompletionService interface defines a framework for handling the completion of asynchronous tasks
- Know how to instantiate the Java ExecutorCompletionService
- Recognize the key methods in the Java CompletionService interface
- Visualize the ExecutorCompletionService in action
- Be aware of how the Java ExecutorCompletionService implements the CompletionService interface
- Know how to apply the Java ConcurrentHashMap class to design a “memoizer”



Memoizer caches function call results & returns cached results for same inputs

Overview of Memoizer

Overview of Memoization

- Memoization is optimization technique used to speed up programs



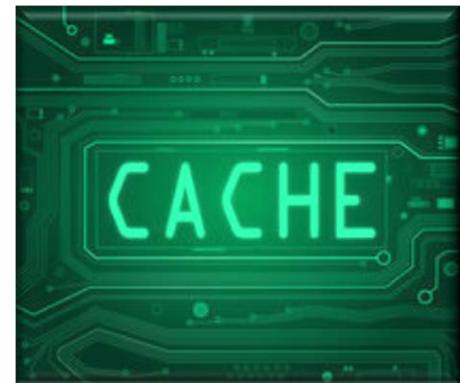
Overview of Memoization

- Memoization is optimization technique used to speed up programs
 - It caches the results of expensive function calls

```
v computeIfAbsent(K key, Function func) {  
    1. If key doesn't exist in cache perform a  
       long-running function associated w/key  
       & store the resulting value via the key  
    2. Return value associated with key  
}
```



Memoizer



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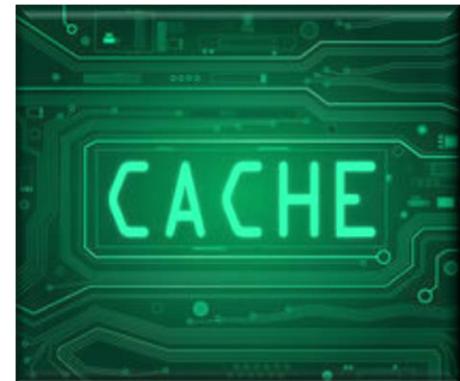
```
v computeIfAbsent(K key, Function func) {
```

- 1. If key doesn't exist in cache perform a long-running function associated w/key & store the resulting value via the key*
- 2. Return value associated with key*

```
}
```



Memoizer



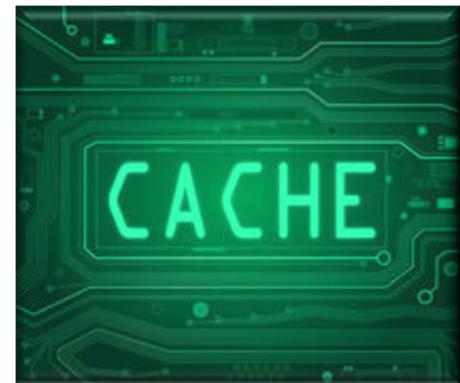
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Memoizer



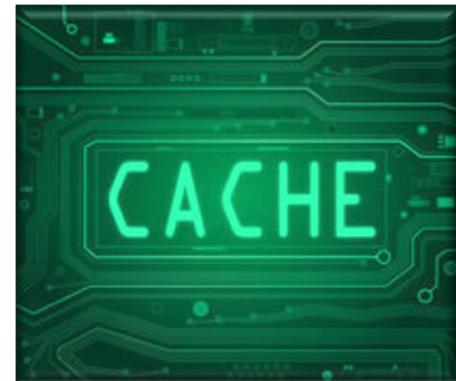
Overview of Memoization

- Memoization is optimization technique used to speed up programs
 - It caches the results of expensive function calls
 - When the same inputs occur again the cached results are simply returned

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v computeIfAbsent(K key, Function func) {  
    1. If key already exists in cache  
        return cached value associated w/key  
    }  
}
```



Memoizer



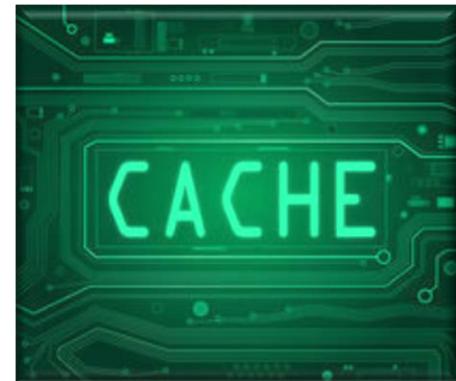
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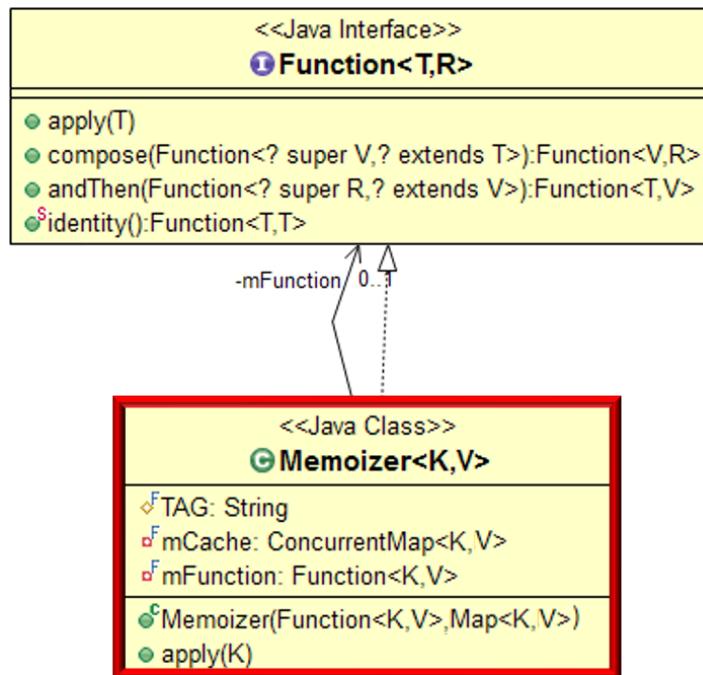
Memoizer



Designing a Memoizer with ConcurrentHashMap

Designing a Memoizer with ConcurrentHashMap

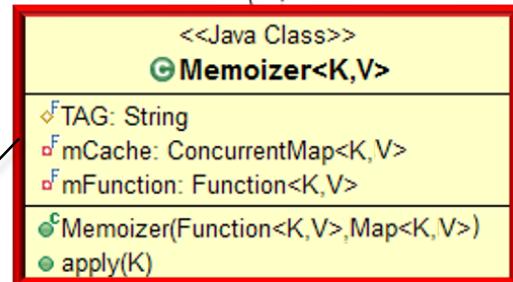
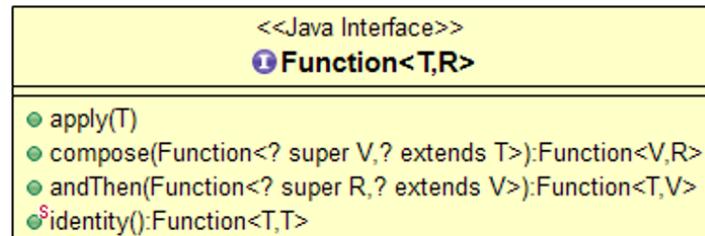
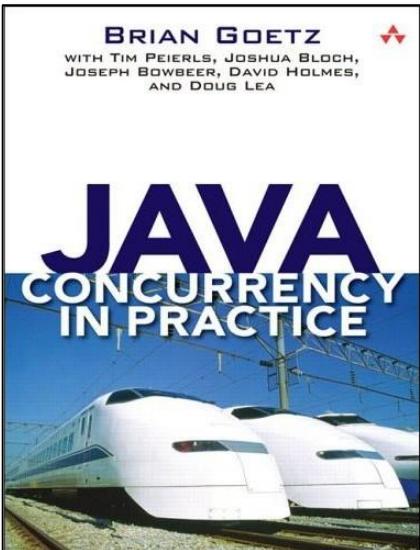
- The Memoizer cache returns a value produced by applying a function to a key



See [PrimeExecutorService/app/src/main/java/vandy/mooc/prime/utils/Memoizer.java](#)

Designing a Memoizer with ConcurrentHashMap

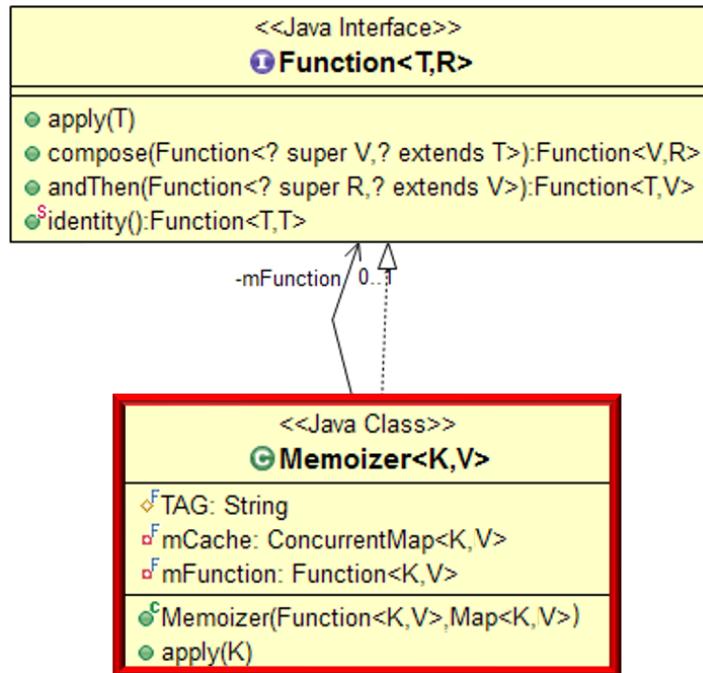
- The Memoizer cache returns a value produced by applying a function to a key



This class is based on "Java Concurrency in Practice" by Brian Goetz et al.

Designing a Memoizer with ConcurrentHashMap

- The Memoizer cache returns a value produced by applying a function to a key
 - A value computed for a key is returned, rather than reapplying the function



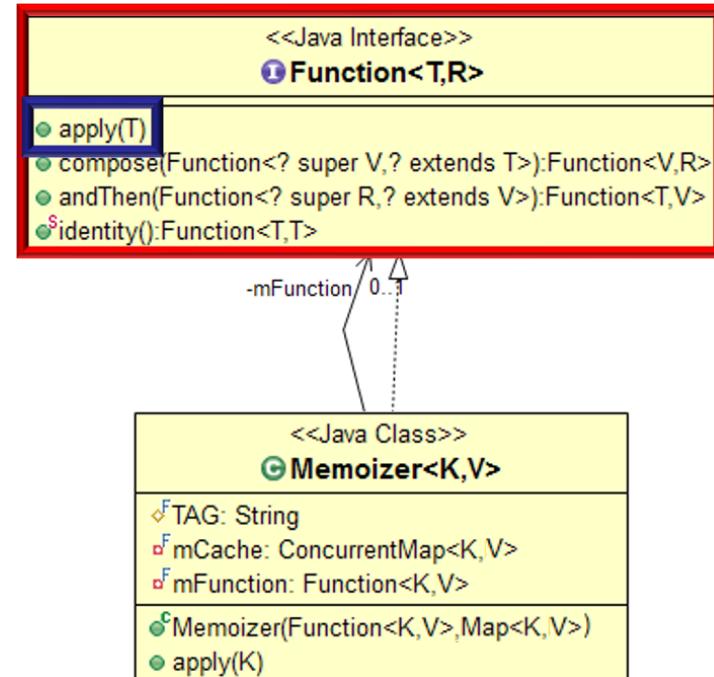
Designing a Memoizer with ConcurrentHashMap

- The Memoizer cache returns a value produced by applying a function to a key

- A value computed for a key is returned, rather than reapplying the function

- Can be used when a Function is expected

```
Function<Long, Long> func =  
    doMemoization  
        ? new Memoizer<>  
            (PrimeCheckers::isPrime,  
             new ConcurrentHashMap());  
        : PrimeCheckers::isPrime;  
  
    ...  
    new PrimeCallable(randomNumber, func); ...
```



Designing a Memoizer with ConcurrentHashMap

- The Memoizer cache returns a value produced by applying a function to a key

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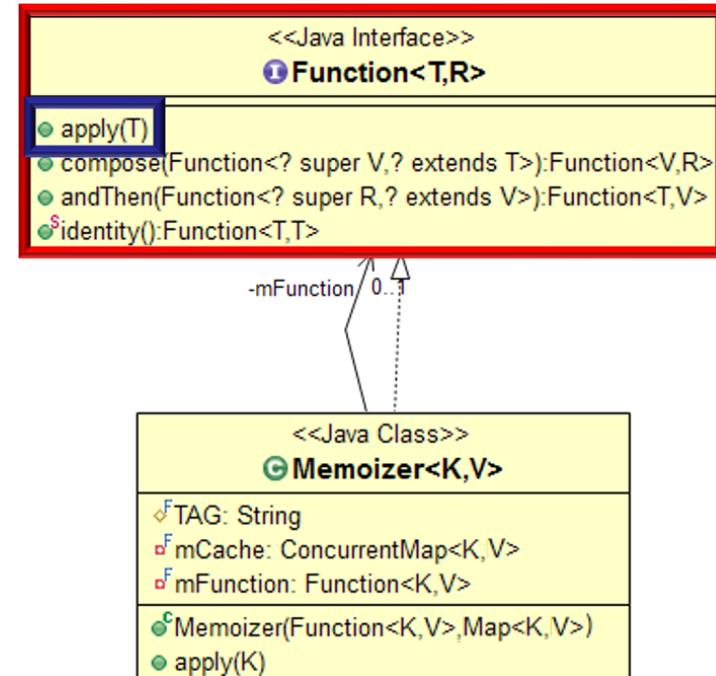
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Function<Long, Long> func =  
    doMemoization  
    ? new Memoizer<>  
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    : PrimeCheckers::isPrime;
```

Use memoizer

...

```
new PrimeCallable(randomNumber, func); ...
```



Designing a Memoizer with ConcurrentHashMap

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- A value computed for a key is returned, rather than reapplying the function

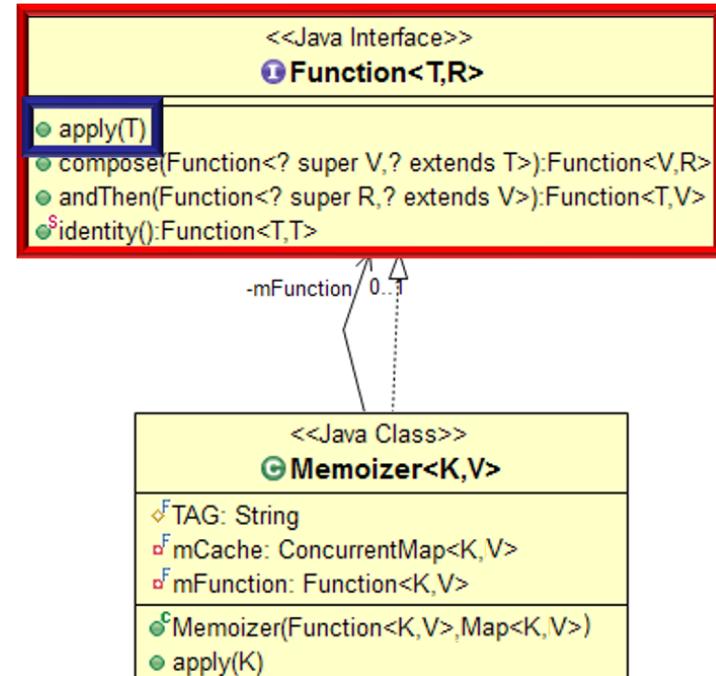
- Can be used when a Function is expected

```
Function<Long, Long> func =  
doMemoization  
? new Memoizer<>  
    (PrimeCheckers::isPrime,  
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: PrimeCheckers::isPrime;
```

Don't use memoizer

...

```
new PrimeCallable(randomNumber, func); ...
```



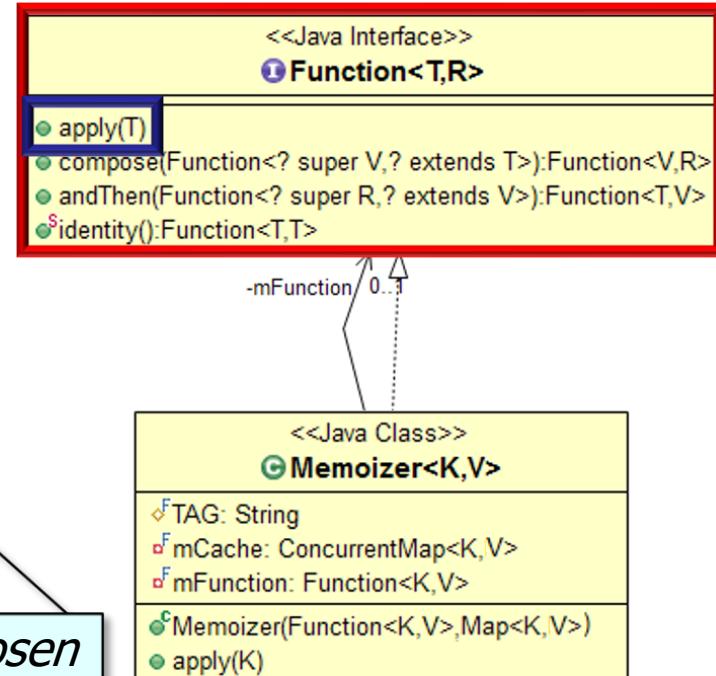
Designing a Memoizer with ConcurrentHashMap

- The Memoizer cache returns a value produced by applying a function to a key
 - A value computed for a key is returned, rather than reapplying the function
 - Can be used when a Function is expected

```
Function<Long, Long> func =  
    doMemoization  
    ? new Memoizer<>  
        (PrimeCheckers::isPrime,  
         new ConcurrentHashMap());  
    : PrimeCheckers::isPrime;
```

func is identical, regardless of which branch is chosen

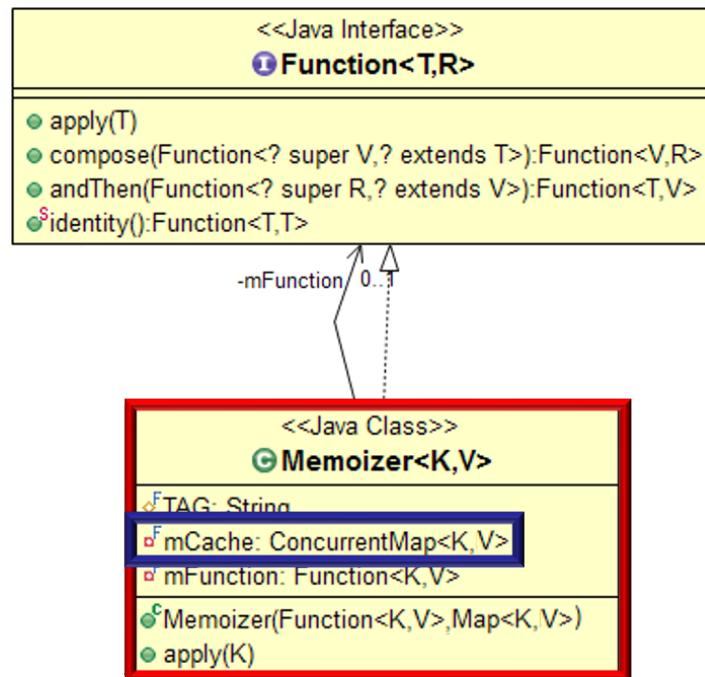
```
...  
new PrimeCallable(randomNumber, func)); ...
```



See upcoming part of this lesson on "Application to PrimeChecker App"

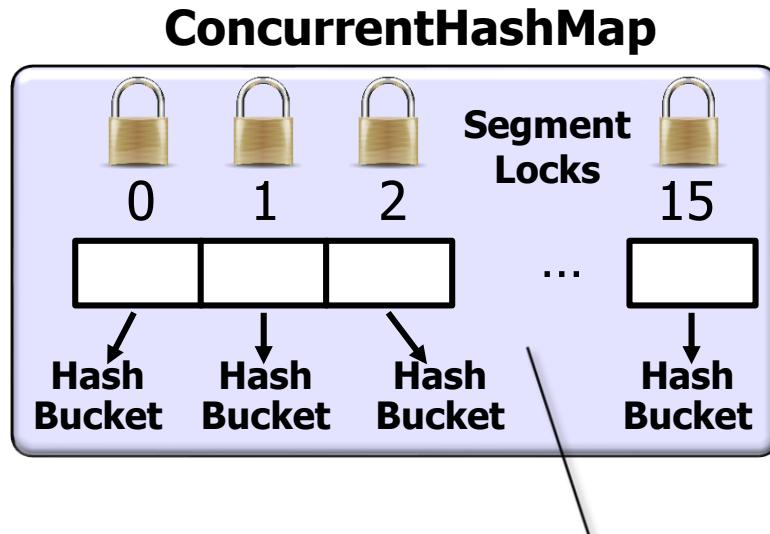
Designing a Memoizer with ConcurrentHashMap

- Memoizer uses a ConcurrentHashMap to minimize synchronization overhead

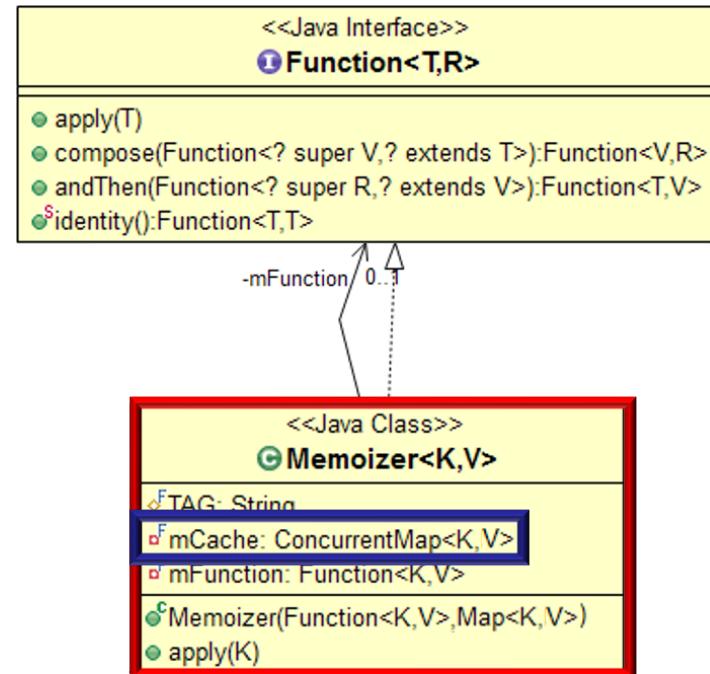


Designing a Memoizer with ConcurrentHashMap

- Memoizer uses a ConcurrentHashMap to minimize synchronization overhead
 - A group of locks guard different subsets of the hash buckets

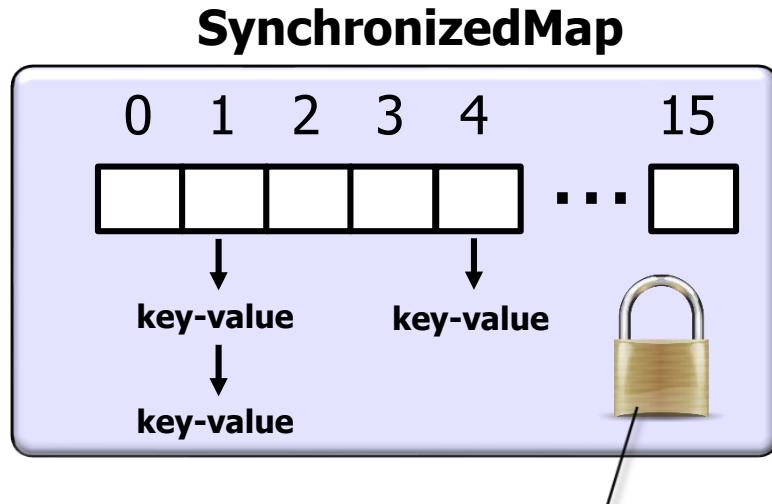


Contention is low due to use of multiple locks

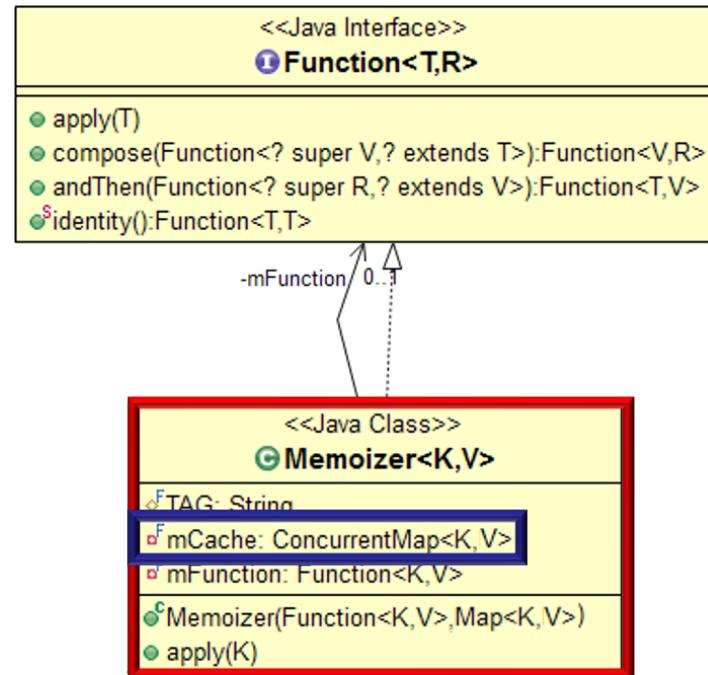


Designing a Memoizer with ConcurrentHashMap

- Memoizer uses a ConcurrentHashMap to minimize synchronization overhead
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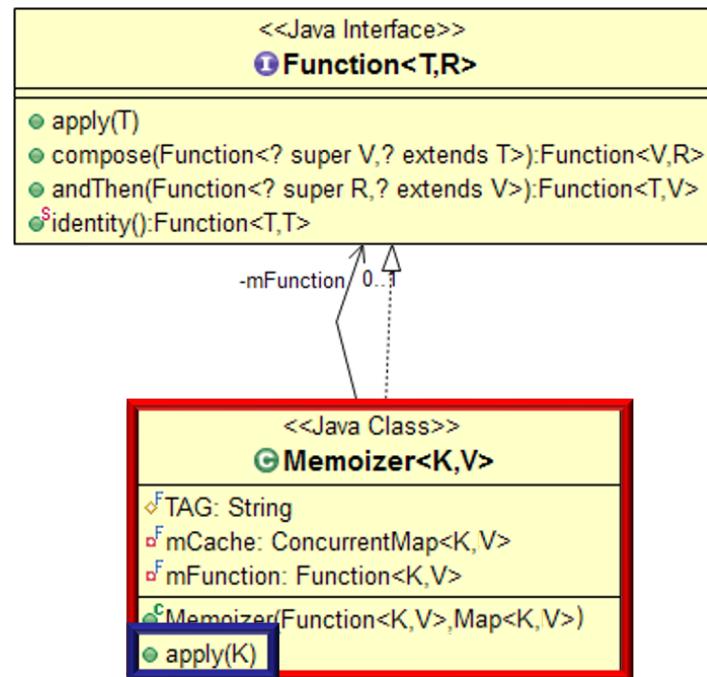


In contrast, a SynchronizedMap uses a single lock



Designing a Memoizer with ConcurrentHashMap

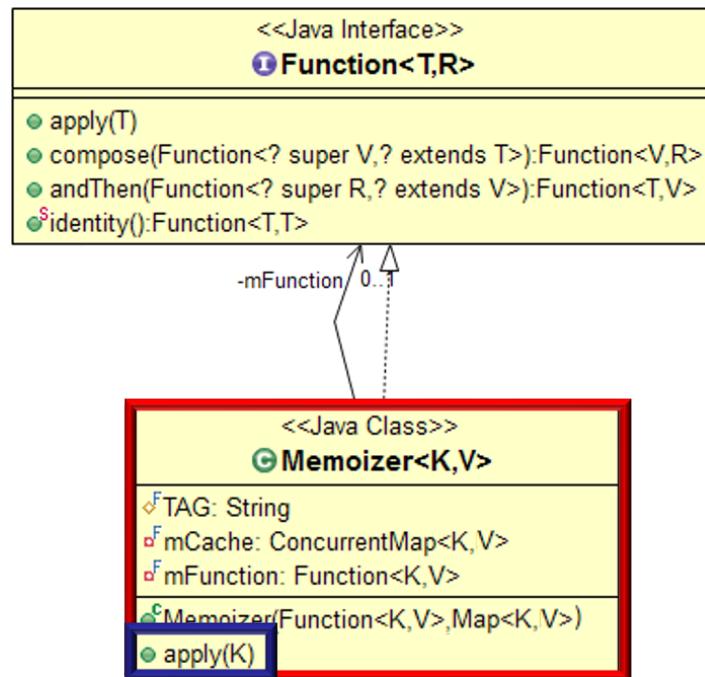
- Memoizer's apply() hook method uses computeIfAbsent() to ensure a function only runs when a key is added to cache



Designing a Memoizer with ConcurrentHashMap

- Memoizer's apply() hook method uses computeIfAbsent() to ensure a function only runs when a key is added to cache, e.g.
- This method implements "atomic check-then-act" semantics

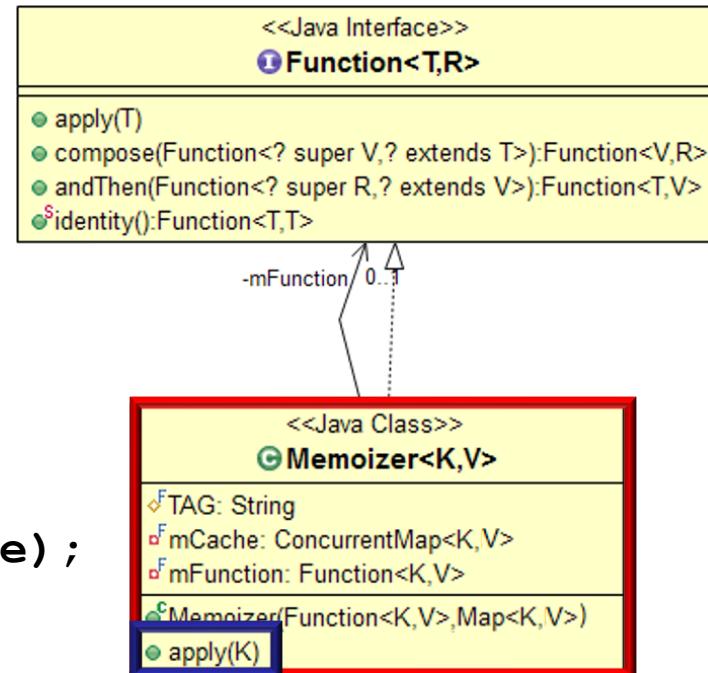
```
return map.computeIfAbsent  
(key,  
    k -> mappingFunc(k));
```



Designing a Memoizer with ConcurrentHashMap

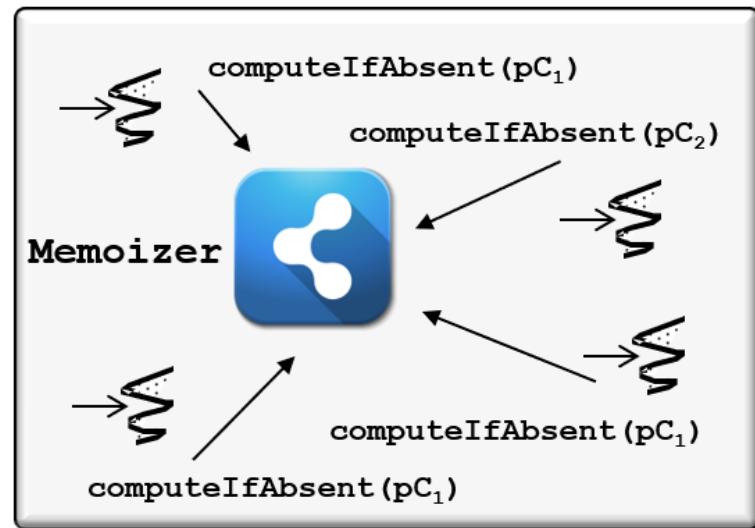
- Memoizer's apply() hook method uses computeIfAbsent() to ensure a function only runs when a key is added to cache, e.g.
 - This method implements “atomic check-then-act” semantics
 - Here's the equivalent sequence of Java (non-atomic/-optimized) code

```
v value = map.get(key);  
if (value == null) {  
    value = mappingFunc.apply(key);  
    if (value != null) map.put(key, value);  
}  
return value;
```



Designing a Memoizer with ConcurrentHashMap

- Memoizer's apply() hook method uses computeIfAbsent() to ensure a function only runs when a key is added to cache, e.g.
 - This method implements "atomic check-then-act" semantics
 - Here's the equivalent sequence of Java (non-atomic/-optimized) code
 - Only one computation per key is performed even if multiple threads simultaneously call computeIfAbsent() using the same key



End of Java ExecutorCompletion Service: Designing a Memoizer