Java ConditionObject: The Guarded Suspension Pattern



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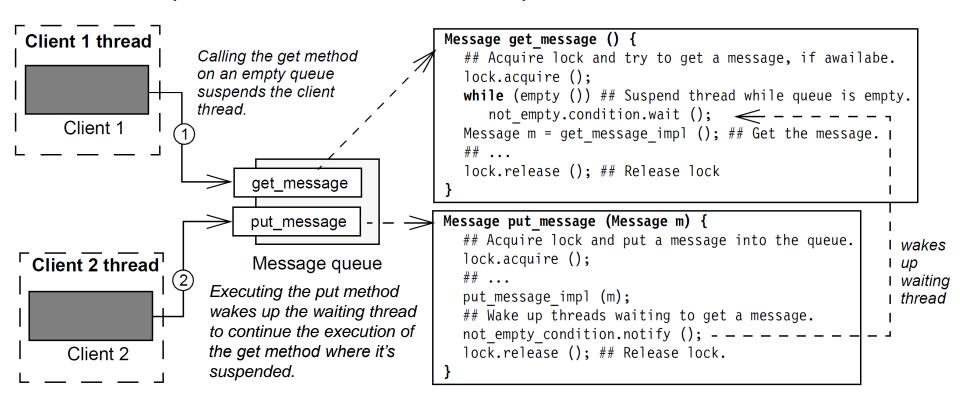
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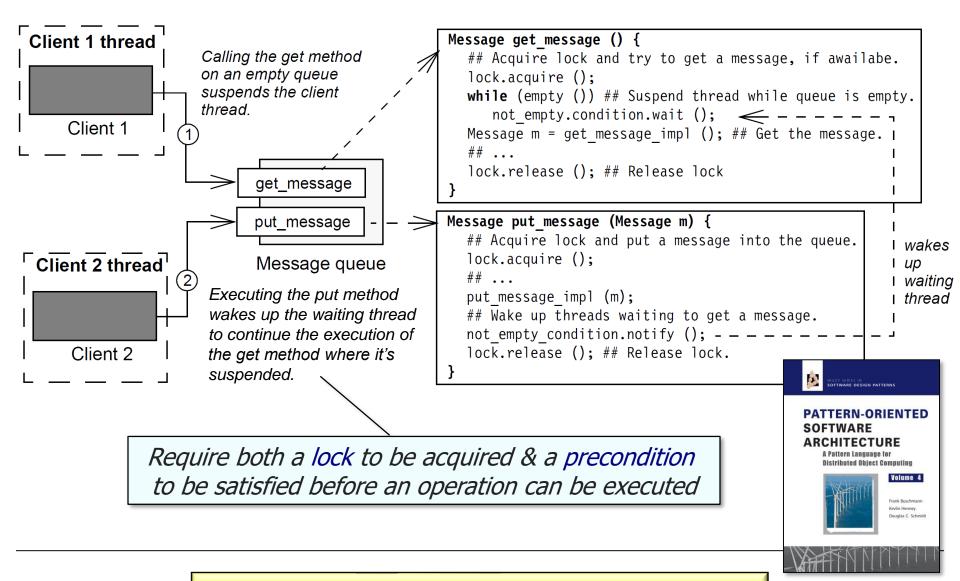


Learning Objectives in this Part of the Lesson

- Understand what condition variables are
- Note a human known use of condition variables
- Know what pattern condition variables implement

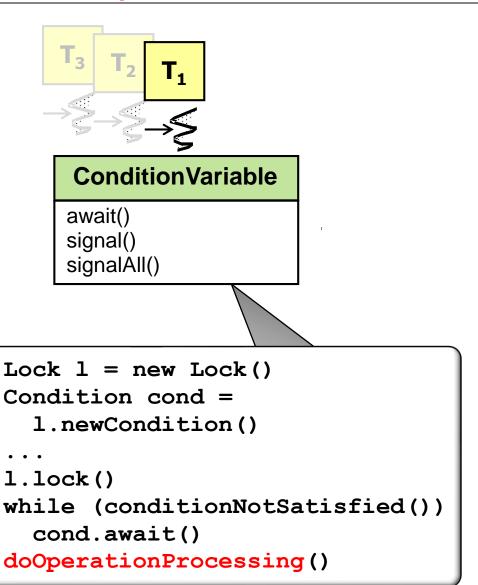


CVs are most often used to implement the Guarded Suspension pattern

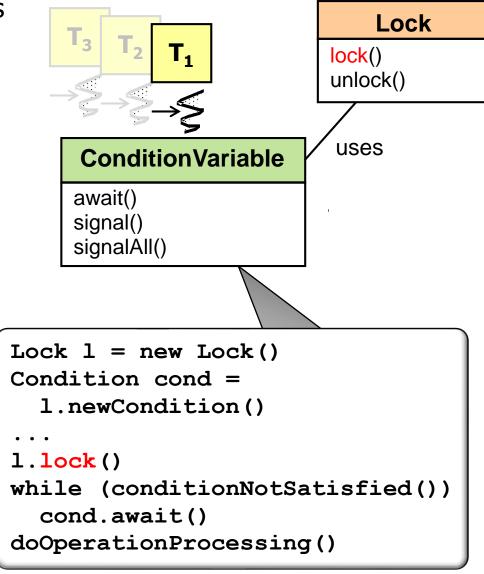


See en.wikipedia.org/wiki/Guarded_suspension

 This pattern is applied to operations that can run only when a condition is satisfied

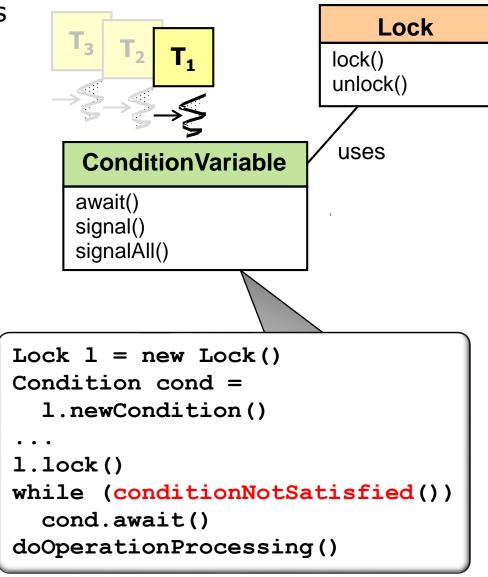


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 - a lock is acquired



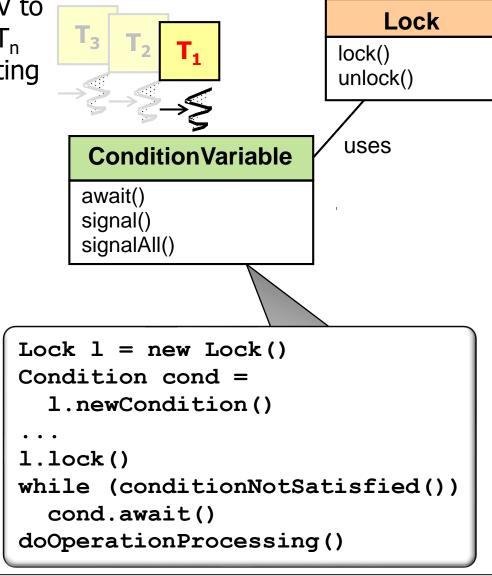
A condition variable is *always* associated with a lock

- This pattern is applied to operations that can run only when a condition is satisfied, e.g.,
 - a lock is acquired
 - a precondition holds

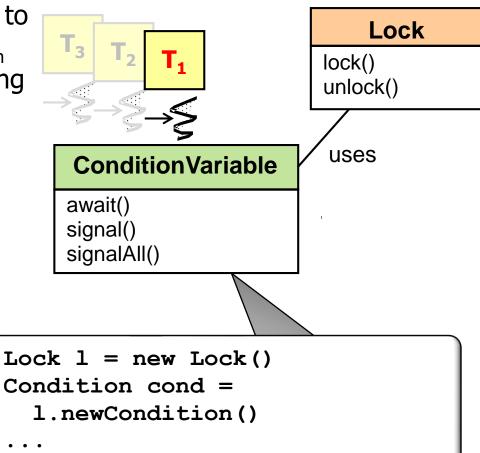


 In this example thread T₁ uses a CV to suspend its execution until thread T_n notifies it that shared state it's waiting on *may* now be satisfied





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while (conditionNotSatisfied())



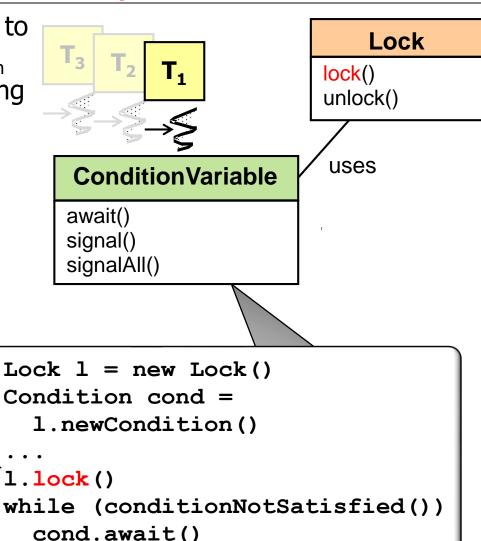
Note the tentative nature of "may"...

1.lock()

cond.await()

doOperationProcessing()

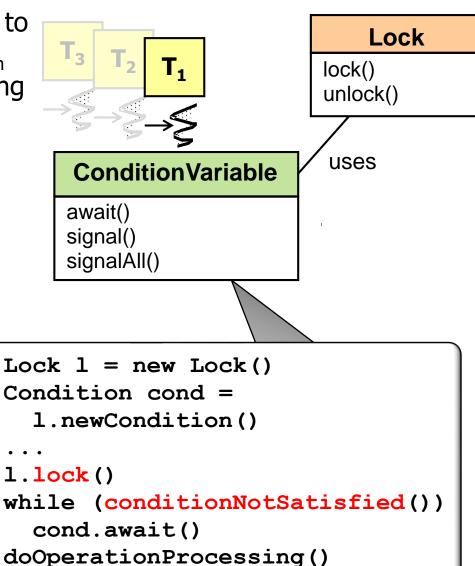
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doOperationProcessing()

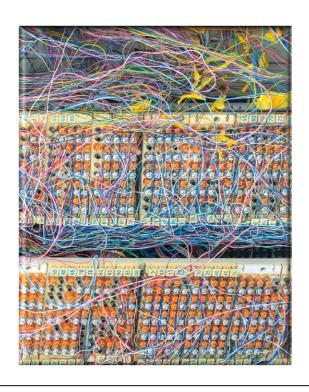
First, a lock must be acquired..

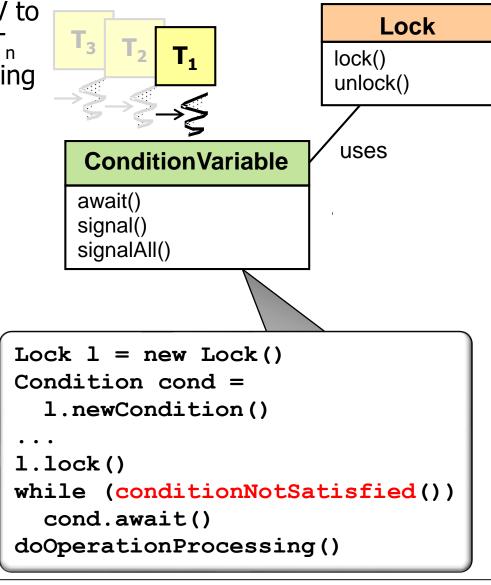
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Second, a condition is checked (in a loop) with the lock held..

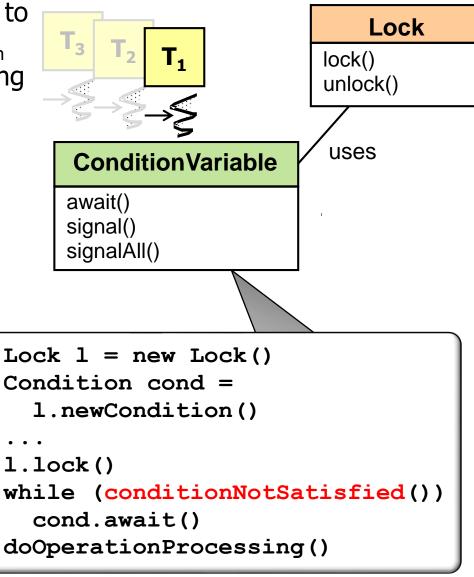
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e.g., a method call, an expression that involves shared state, etc.



Any state shared between threads must be protected by a lock associated with the CV

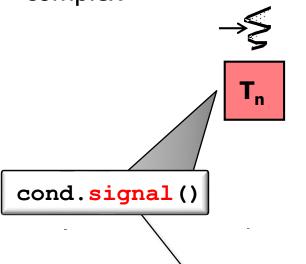
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The calling thread will block (possibly repeatedly) while the condition is not satisfied (await() atomically releases the lock)

```
Lock
                          lock()
                          unlock()
                           uses
     ConditionVariable
     await()
     signal()
     signalAll()
Lock l = new Lock()
Condition cond =
  1.newCondition()
1.lock()
while (conditionNotSatisfied())
  cond.await()
```

doOperationProcessing()

- In this example thread T₁ uses a CV to suspend its execution until thread T_n notifies it that shared state it's waiting on *may* now be satisfied
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Another thread can signal condition when shared state may now be true

```
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                          unlock()
                          uses
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Lock l = new Lock()
Condition cond =
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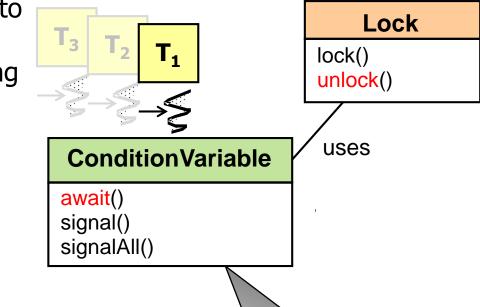
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1.lock()
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```

await() reacquires the lock & condition is rechecked in loop

- In this example thread T₁ uses a CV to suspend its execution until thread T_n notifies it that shared state it's waiting on *may* now be satisfied
 - A condition can be arbitrarily complex
 - Waiting on a CV releases the lock
 & suspends the thread atomically



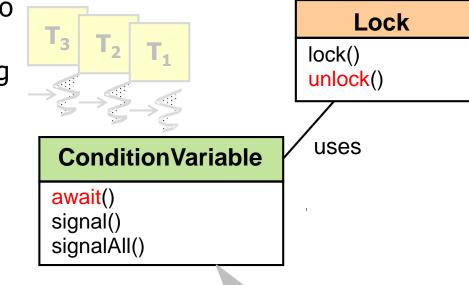


```
Lock 1 = new Lock()
Condition cond =
    l.newCondition()
...
l.lock()
while (conditionNotSatisfied())
    cond.await()
doOperationProcessing()
```

The lock is released when the thread is suspended on the CV

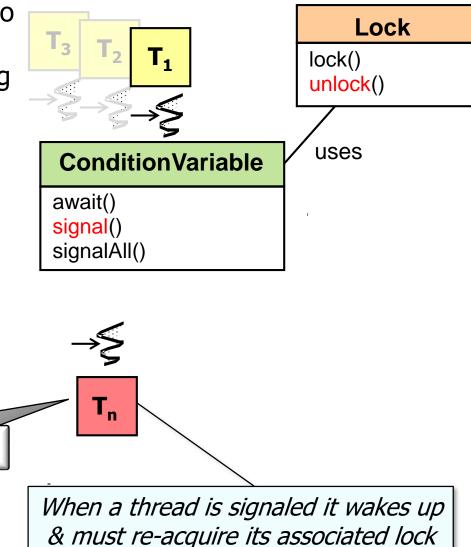
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 - Thread T₁ is suspended until thread T_n signals the CV





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Condition cond =
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l.lock()
while (conditionNotSatisfied())
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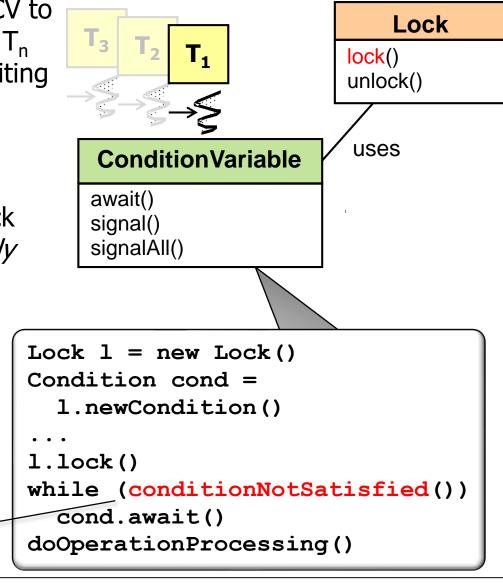
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cond.signal()

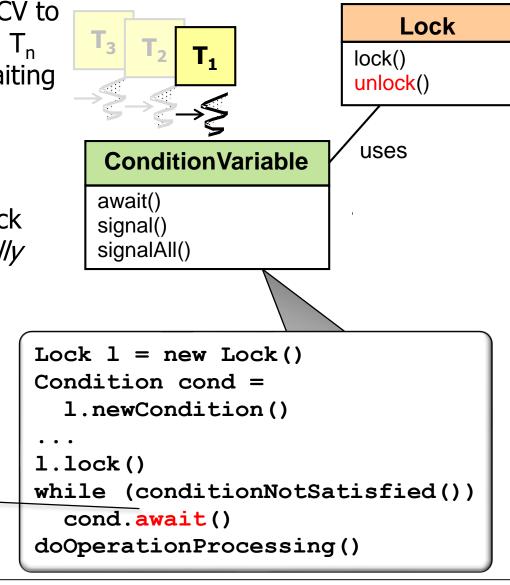
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After lock is re-acquired the thread can reevaluate its condition to see if it's satisfied



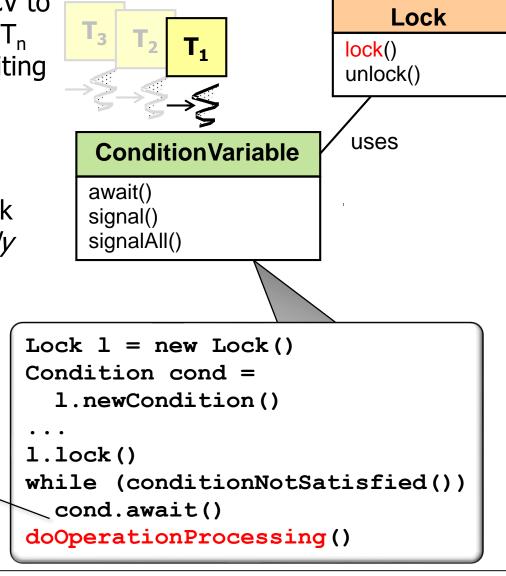
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 - Thread T₁ is suspended until thread T_n signals the CV

If condition is not satisfied the thread must wait (which releases the lock atomically)



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 - A condition can be arbitrarily complex
 - Waiting on a CV releases the lock
 & suspends the thread atomically
 - Thread T₁ is suspended until thread T_n signals the CV

After the lock is re-acquired & the condition is satisfied the operation can proceed (with lock held)



End of Java ConditionObject: The Guarded Suspension Pattern