Managing the Java Thread Lifecycle: Java Thread Interrupts vs. Hardware/OS Interrupts



Douglas C. Schmidt

<u>d.schmidt@vanderbilt.edu</u>

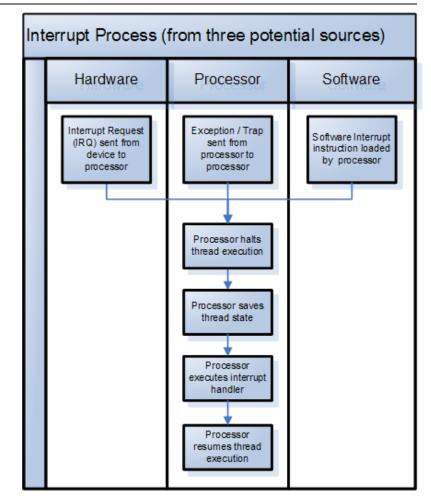
www.dre.vanderbilt.edu/~schmidt

Institute for Software Integrated Systems Vanderbilt University Nashville, Tennessee, USA

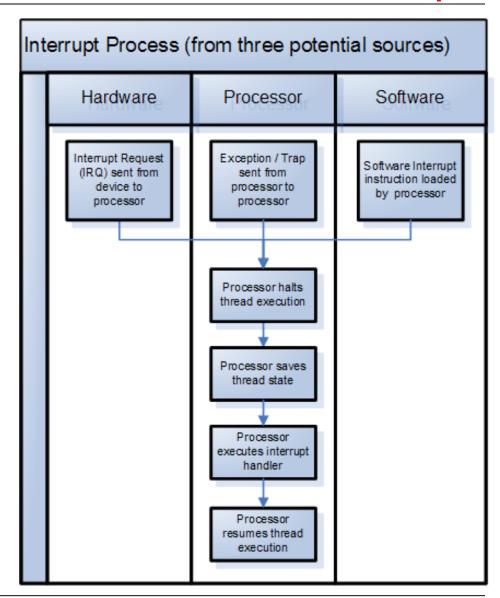


Learning Objectives in this Part of the Lesson

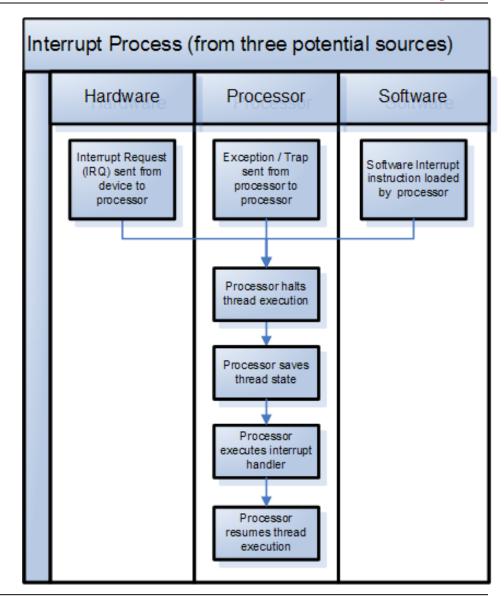
- Know various ways to stop Java threads
 - Stopping a thread with a volatile flag
 - Stopping a thread with an interrupt request
 - Learn the patterns of interrupting Java threads
 - Understand differences between a Java thread interrupt & a hardware/OS interrupt



Interrupts at the hardware or OS layers have several properties



- Interrupts at the hardware or OS layers have several properties
 - Asynchronous
 - Can occur essentially anytime
 & are independent of the instruction currently running

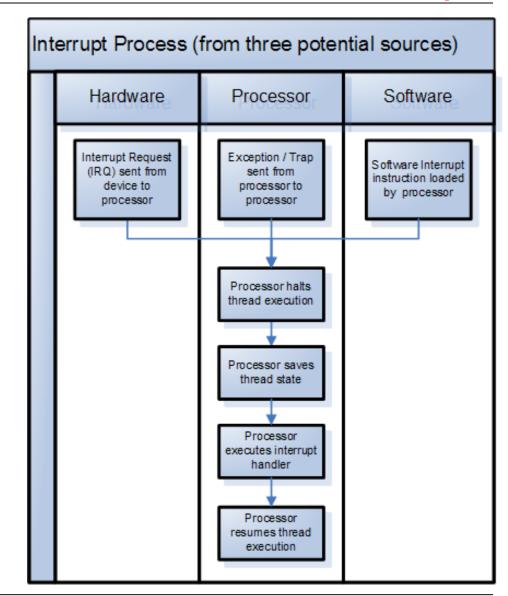


See vujungle.blogspot.com/2010/12/differentiate-synchronous-and.html

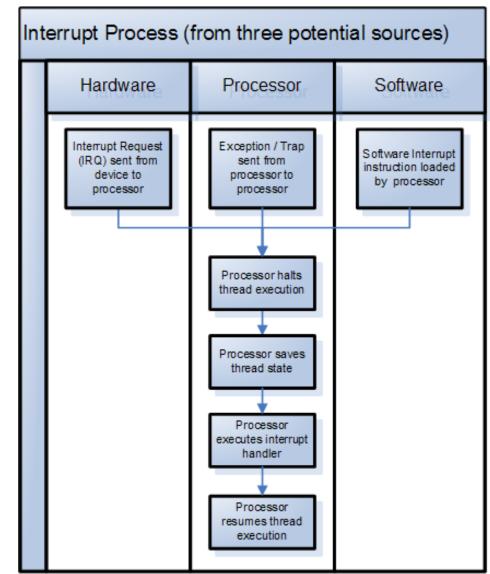
Interrupts at the hardware or OS layers have several properties

Asynchronous

- Can occur essentially anytime
 & are independent of the instruction currently running
- A program needn't test for them explicitly since they occur "out-of-band"



- Interrupts at the hardware or OS layers have several properties
 - Asynchronous
 - Preemptive
 - Pause (& then later resume) the execution of currently running code without its cooperation



See en-wikipedia.org/wiki/Preemption_(computing)

}

 This example shows how to catch the UNIX SIGINT signal

```
void sig_handler(int signo) {
   if (signo == SIGINT)
      printf("received SIGINT\n");
int main(void) {
  if (signal(SIGINT, sig_handler)
      == SIG ERR)
    printf("can't catch SIGINT\n");
  for (;;)
    sleep(10);
  return 0;
```

- This example shows how to catch the UNIX SIGINT signal
 - It occurs asynchronously

The SIGINT interrupt is typically generated by typing ^C in a UNIX shell

```
void sig_handler(int signo) {
   if (signo == SIGINT)
      printf("received SIGINT\n");
int main(void) {
  if (signal(SIGINT, sig_handler)
      == SIG ERR)
    printf("can't catch SIGINT\n");
  for (;;)
    sleep(10);
  return 0;
```

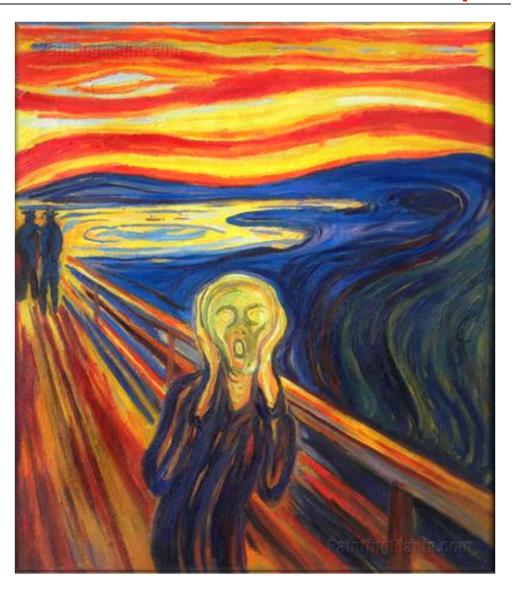
- This example shows how to catch the UNIX SIGINT signal
 - It occurs asynchronously
 - It preempts the current instruction

```
void sig handler(int signo) {
   if (signo == SIGINT)
      printf("received SIGINT\n");
int main(void) {
  if (signal(SIGINT, sig_handler)
      == SIG ERR)
    printf("can't catch SIGINT\n");
  for (;;)
    sleep(10);
  return 0;
```

- This example shows how to catch the UNIX SIGINT signal
 - It occurs asynchronously
 - It preempts the current instruction
 - It needn't be tested for explicitly

```
void sig_handler(int signo) {
   if (signo == SIGINT)
      printf("received SIGINT\n");
int main(void) {
  if (signal(SIGINT, sig_handler)
      == SIG ERR)
    printf("can't catch SIGINT\n");
  for (;;)
    sleep(10);
  return 0;
```

 Asynchronous & preemptive interrupt handling make it hard to reason about programs



See en.wikipedia.org/wiki/Unix_signal#Risks

 Asynchronous & preemptive interrupt handling make it hard to reason about programs, e.g.

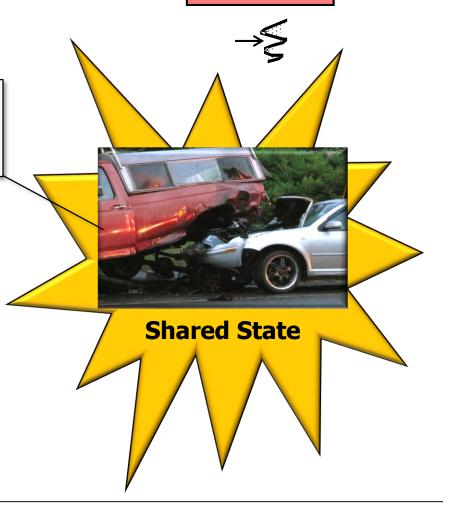
Race conditions

Race conditions occur when a program depends on the sequence or timing of threads for it to operate properly

Thread₁

->\$

Thread₂



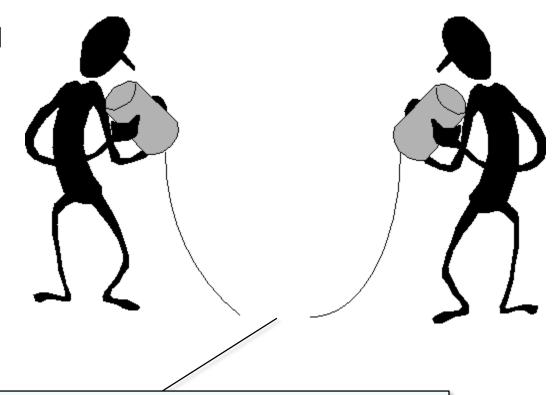
See en.wikipedia.org/wiki/Race_condition#Software

- Asynchronous & preemptive interrupt handling make it hard to reason about programs, e.g.
 - Race conditions
 - Re-entrancy problems

A non-reentrant function cannot be interrupted in the middle of its execution & then safely called again before its previous invocations complete execution



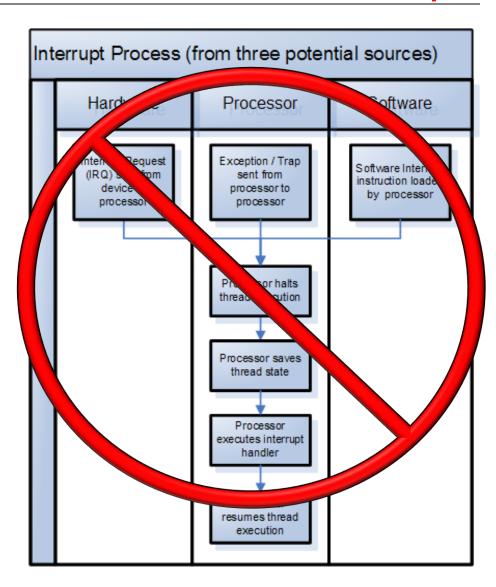
- Asynchronous & preemptive interrupt handling make it hard to reason about programs, e.g.
 - Race conditions
 - Re-entrancy problems
 - Non-transparent restarts



e.g., an I/O operation returns the # of bytes transferred & it is up to the application to check this & manage its own resumption of the operation until all the bytes have been transferred

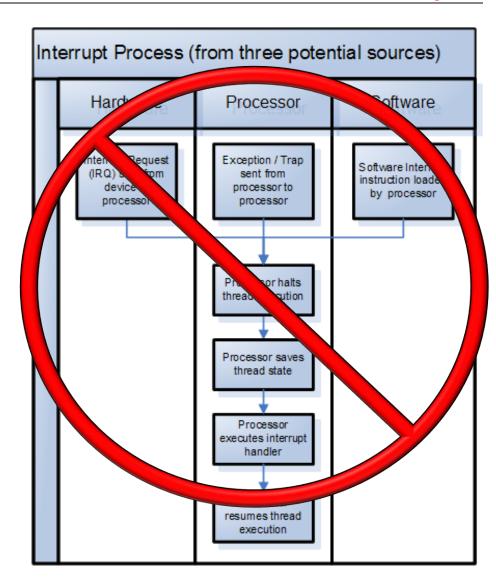
See en.wikipedia.org/wiki/PCLSRing#Unix-solution: restart on request

 Java thread interrupts differ from hardware or operating system interrupts



See docs.oracle.com/javase/tutorial/essential/concurrency/interrupt.html

- Java thread interrupts differ from hardware or operating system interrupts, e.g.
 - Delivery is synchronous & non-preemptive rather than asynchronous & preemptive
 - i.e., they don't occur at an arbitrary point & don't pause (& later resume) running code



- Java thread interrupts differ from hardware or operating system interrupts, e.g.
 - Delivery is synchronous & non-preemptive rather than asynchronous & preemptive
 - A program must test for them explicitly

- Java thread interrupts differ from hardware or operating system interrupts, e.g.
 - Delivery is synchronous & non-preemptive rather than asynchronous & preemptive
 - A program must test for them explicitly
 - i.e., InterruptedException is (usually) thrown synchronously
 & is handled synchronously

- Java thread interrupts differ from hardware or operating system interrupts, e.g.
 - Delivery is synchronous & non-preemptive rather than asynchronous & preemptive
 - A program must test for them explicitly
 - Certain operations cannot be interrupted
 - e.g., blocking I/O calls that aren't "interruptable channels"

```
static class SleeperThread
       extends Thread {
 public void run() {
    int c;
    try {
       c = System.in.read();
                Please
```

End of Managing the Java Thread Lifecycle: Java Thread Interrupts vs. Hardware/OS Interrupts