# **The Guarded Suspension Pattern**



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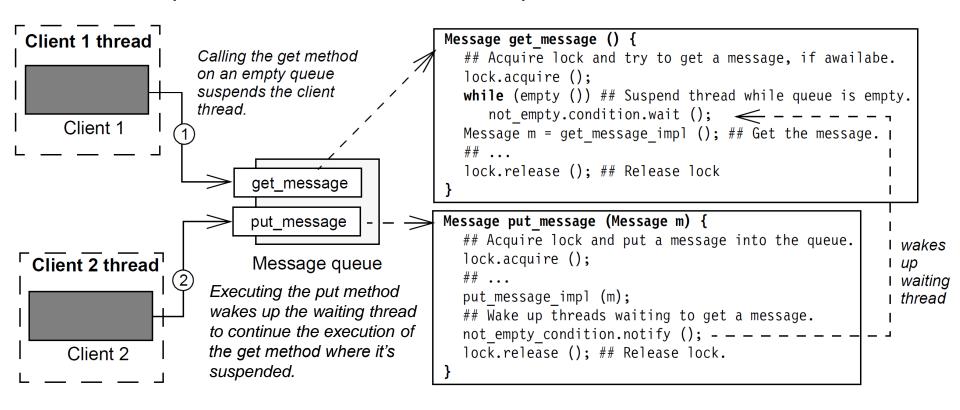
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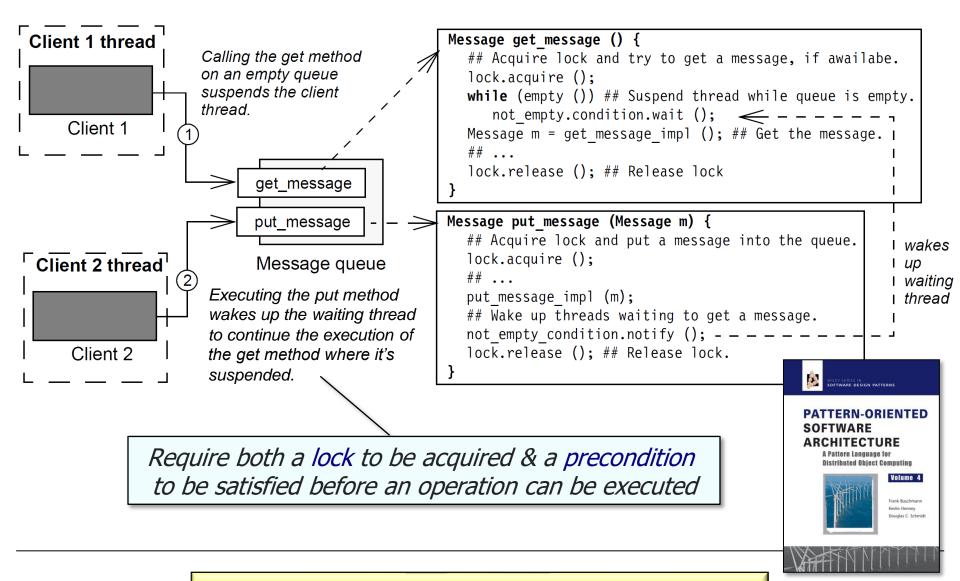


#### Learning Objectives in this Part of the Lesson

- Understand what condition variables are
- Note a human known use of condition variables
- Know what pattern condition variables implement

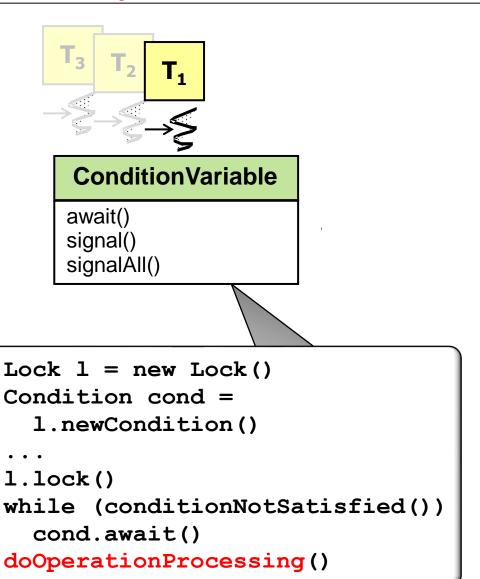


CVs are most often used to implement the Guarded Suspension pattern

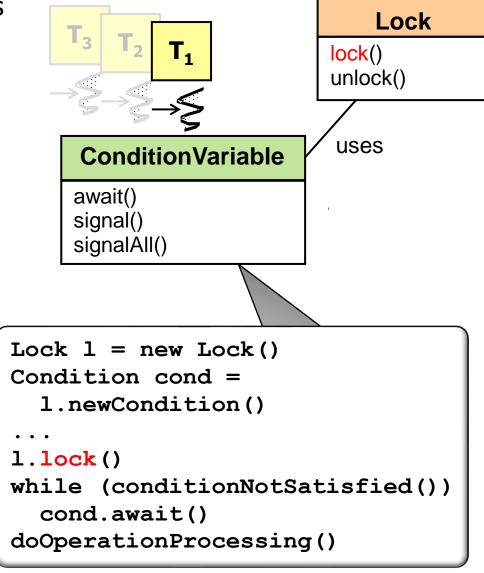


See en.wikipedia.org/wiki/Guarded\_suspension

 This pattern is applied to operations that can run only when a condition is satisfied

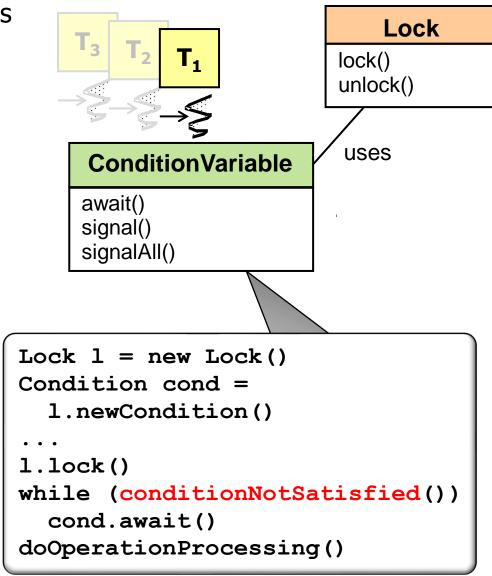


- This pattern is applied to operations that can run only when a condition is satisfied, e.g.,
  - a lock is acquired



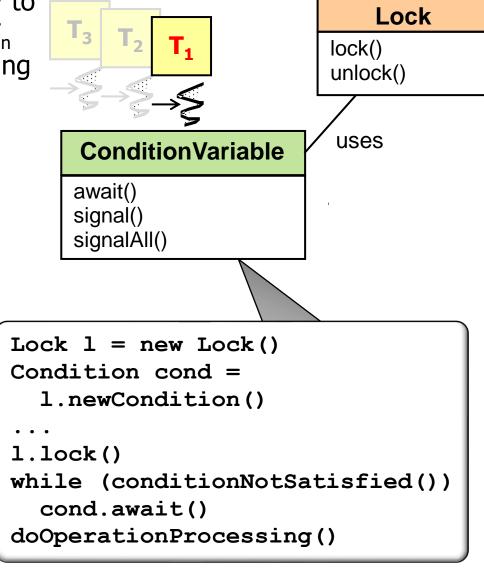
A condition variable is *always* associated with a lock

- This pattern is applied to operations that can run only when a condition is satisfied, e.g.,
  - a lock is acquired
  - a precondition holds

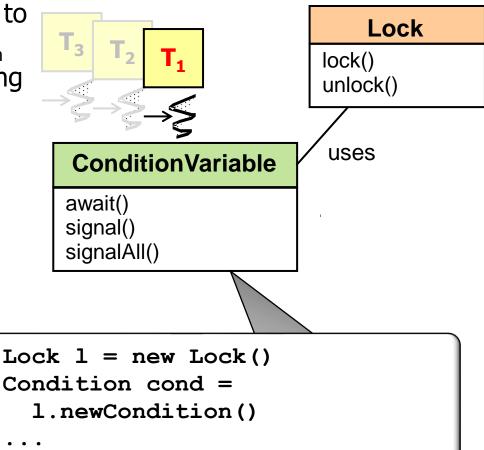


 In this example thread T<sub>1</sub> uses a CV to suspend its execution until thread T<sub>n</sub> notifies it that shared state it's waiting on *may* now be satisfied





• In this example thread  $T_1$  uses a CV to suspend its execution until thread  $T_n$  notifies it that shared state it's waiting on may now be satisfied



while (conditionNotSatisfied())



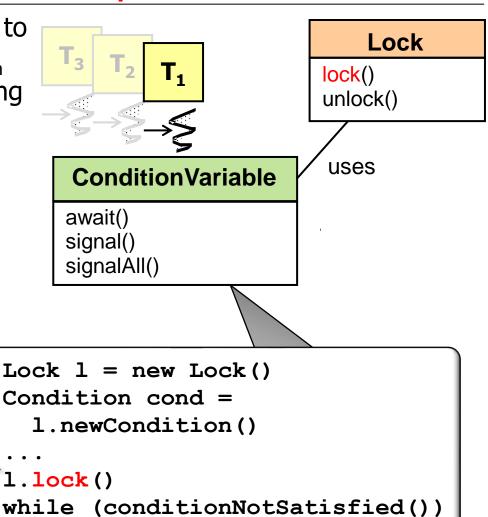
Note the tentative nature of "may"...

1.lock()

cond.await()

doOperationProcessing()

 In this example thread T<sub>1</sub> uses a CV to suspend its execution until thread T<sub>n</sub> notifies it that shared state it's waiting on *may* now be satisfied

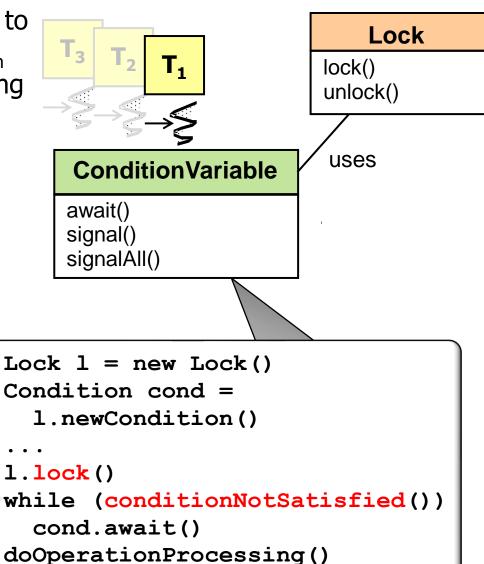


First, a lock must be acquired..

cond.await()

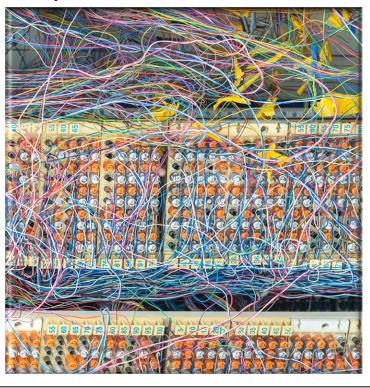
doOperationProcessing()

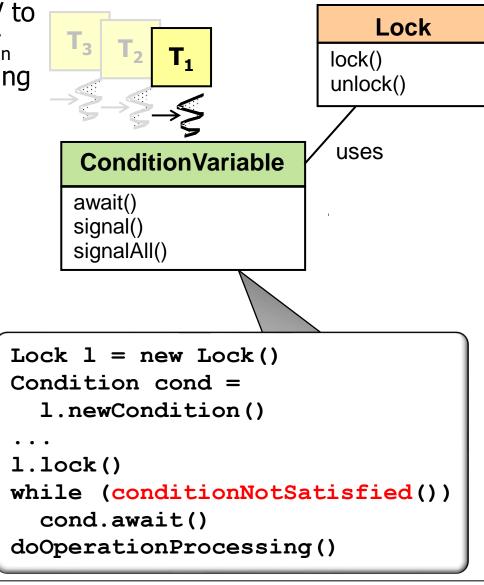
 In this example thread T<sub>1</sub> uses a CV to suspend its execution until thread T<sub>n</sub> notifies it that shared state it's waiting on *may* now be satisfied



Second, a condition is checked (in a loop) with the lock held..

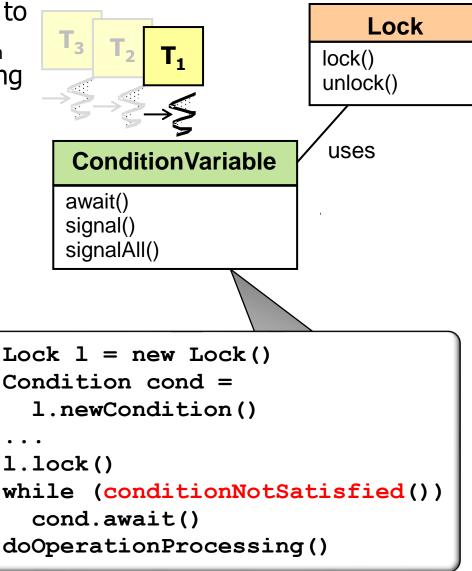
- In this example thread T<sub>1</sub> uses a CV to suspend its execution until thread T<sub>n</sub> notifies it that shared state it's waiting on *may* now be satisfied
  - A condition can be arbitrarily complex





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e.g., a method call, an expression that involves shared state, etc.



Any state shared between threads must be protected by a lock associated with the CV

- In this example thread T<sub>1</sub> uses a CV to suspend its execution until thread T<sub>n</sub> notifies it that shared state it's waiting on *may* now be satisfied
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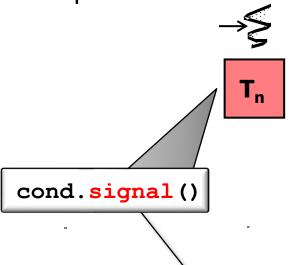
The calling thread will block (possibly repeatedly) while the condition is not satisfied (await() atomically releases the lock)

```
Lock
                           lock()
                           unlock()
                           uses
      ConditionVariable
     await()
     signal()
     signalAll()
Lock l = new Lock()
Condition cond =
  1.newCondition()
1.lock()
while (conditionNotSatisfied())
```

cond.await()

doOperationProcessing()

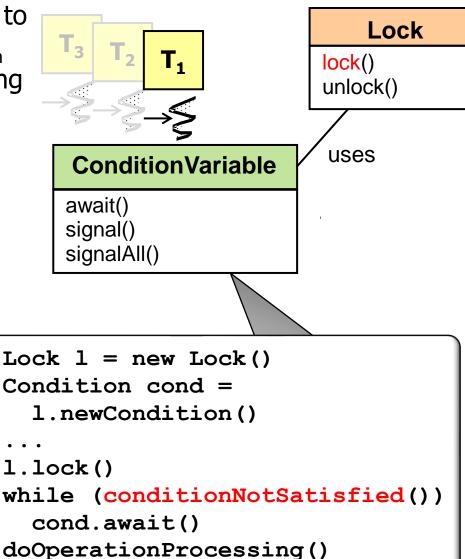
- In this example thread T<sub>1</sub> uses a CV to suspend its execution until thread T<sub>n</sub> notifies it that shared state it's waiting on *may* now be satisfied
  - A condition can be arbitrarily complex



Another thread can signal condition when shared state may now be true

```
Lock
                          lock()
                          unlock()
                          uses
     ConditionVariable
     await()
     signal()
     signalAll()
Lock l = new Lock()
Condition cond =
  l.newCondition()
1.lock()
while (conditionNotSatisfied())
  cond.await()
doOperationProcessing()
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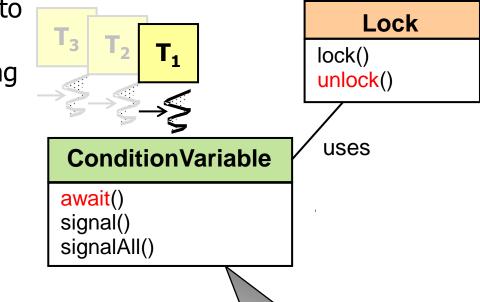
- In this example thread T<sub>1</sub> uses a CV to suspend its execution until thread T<sub>n</sub> notifies it that shared state it's waiting on may now be satisfied
  - A condition can be arbitrarily complex



await() reacquires the lock & condition is rechecked in loop

- In this example thread T<sub>1</sub> uses a CV to suspend its execution until thread T<sub>n</sub> notifies it that shared state it's waiting on *may* now be satisfied
  - A condition can be arbitrarily complex
  - Waiting on a CV releases the lock
     & suspends the thread atomically



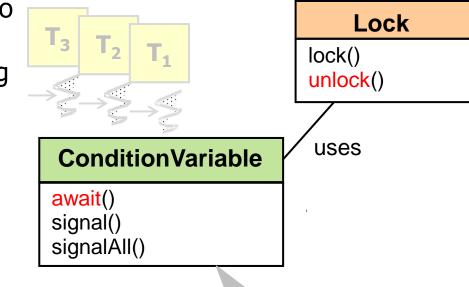


```
Lock 1 = new Lock()
Condition cond =
    l.newCondition()
...
l.lock()
while (conditionNotSatisfied())
    cond.await()
doOperationProcessing()
```

The lock is released when the thread is suspended on the CV

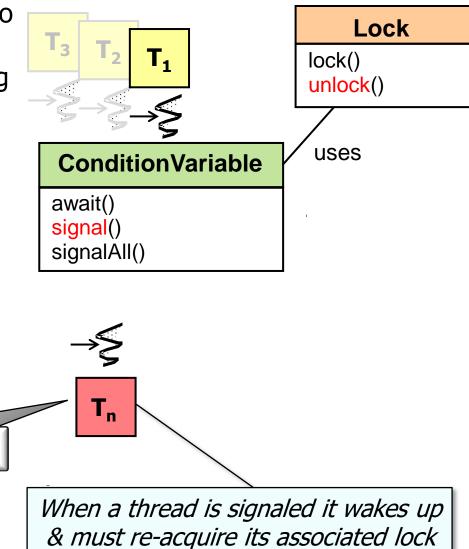
- In this example thread T<sub>1</sub> uses a CV to suspend its execution until thread T<sub>n</sub> notifies it that shared state it's waiting on *may* now be satisfied
  - A condition can be arbitrarily complex
  - Waiting on a CV releases the lock
     & suspends the thread atomically
    - Thread T<sub>1</sub> is suspended until thread T<sub>n</sub> signals the CV





```
Lock 1 = new Lock()
Condition cond =
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l.lock()
while (conditionNotSatisfied())
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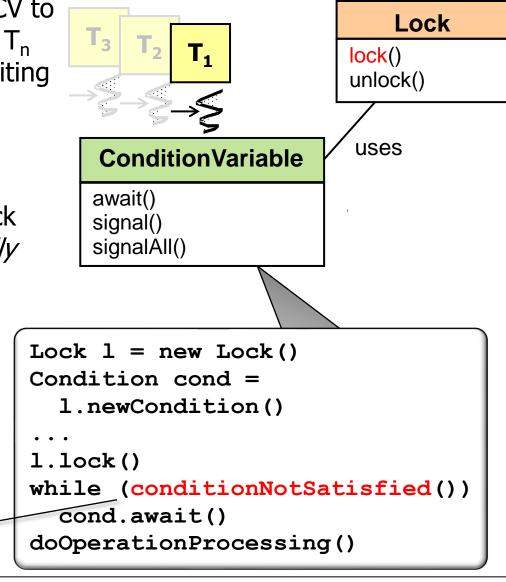
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    - Thread T<sub>1</sub> is suspended until thread T<sub>n</sub> signals the CV



cond.signal()

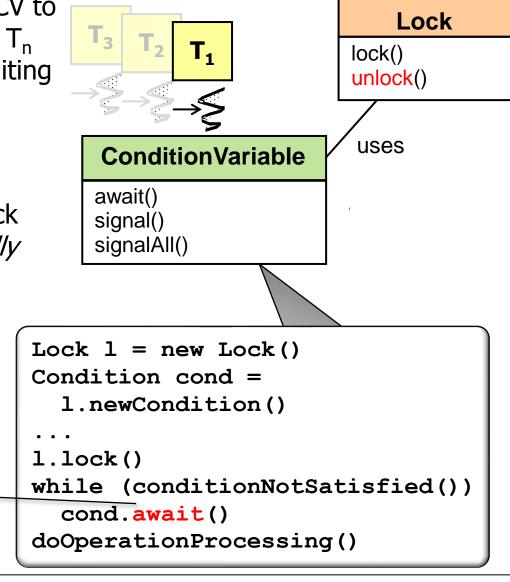
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  - A condition can be arbitrarily complex
  - Waiting on a CV releases the lock
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After lock is re-acquired the thread can reevaluate its condition to see if it's satisfied



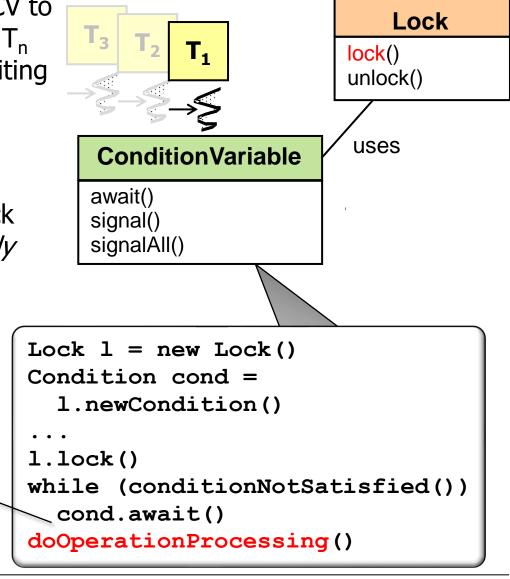
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  - Waiting on a CV releases the lock
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If condition is not satisfied the thread must wait (which releases the lock atomically)



- In this example thread T<sub>1</sub> uses a CV to suspend its execution until thread T<sub>n</sub> notifies it that shared state it's waiting on *may* now be satisfied
  - A condition can be arbitrarily complex
  - Waiting on a CV releases the lock
     & suspends the thread atomically
    - Thread T<sub>1</sub> is suspended until thread T<sub>n</sub> signals the CV

After the lock is re-acquired & the condition is satisfied the operation can proceed (with lock held)



# End of the Guarded Suspension Pattern