## Java Concurrent Collections: ConcurrentHashMap & BlockingQueue



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## Learning Objectives in this Lesson

- Understand the capabilities of Java's concurrent collections
- Recognize the capabilities of Java's ConcurrentHashMap & BlockingQueue

#### Interface BlockingQueue<E>

#### **Type Parameters:**

E - the type of elements held in this collection

#### **All Superinterfaces:**

Collection<E>, Iterable<E>, Queue<E>

#### **All Known Subinterfaces:**

BlockingDeque<E>, TransferQueue<E>

#### All Known Implementing Classes:

ArrayBlockingQueue, DelayQueue, LinkedBlockingDeque, LinkedBlockingQueue, LinkedTransferQueue, PriorityBlockingQueue, SynchronousQueue

public interface BlockingQueue<E>
extends Oueue<E>

A Queue that additionally supports operations that wait for the queue to become non-empty when retrieving an element, and wait for space to become available in the queue when storing an element.

#### Class ConcurrentHashMap<K,V>

java.lang.Object java.util.AbstractMap<K,V> java.util.concurrent.ConcurrentHashMap<K,V>

#### **Type Parameters:**

K - the type of keys maintained by this map

V - the type of mapped values

#### All Implemented Interfaces:

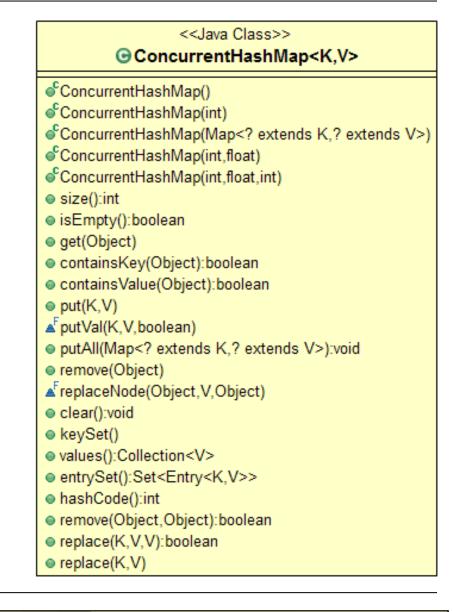
Serializable, ConcurrentMap<K,V>, Map<K,V>

public class ConcurrentHashMap<K,V>
extends AbstractMap<K,V>
implements ConcurrentMap<K,V>, Serializable

A hash table supporting full concurrency of retrievals and high expected concurrency for updates. This class obeys the same functional specification as Hashtable, and includes versions of methods corresponding to each method of Hashtable. However, even though all operations are thread-safe, retrieval operations do *not* entail locking, and there is *not* any support for locking the entire table in a way that prevents all access. This class is fully interoperable with Hashtable in programs that rely on its thread safety but not on its synchronization details.

Retrieval operations (including get) generally do not block, so may overlap with update operations (including put and remove). Retrievals reflect the results of the most recently *completed* update operations holding upon their onset. (More formally, an update operation for a given key bears a *happens-before* relation with any (non-null) retrieval for that key reporting the updated value.) For aggregate operations such as putAll and

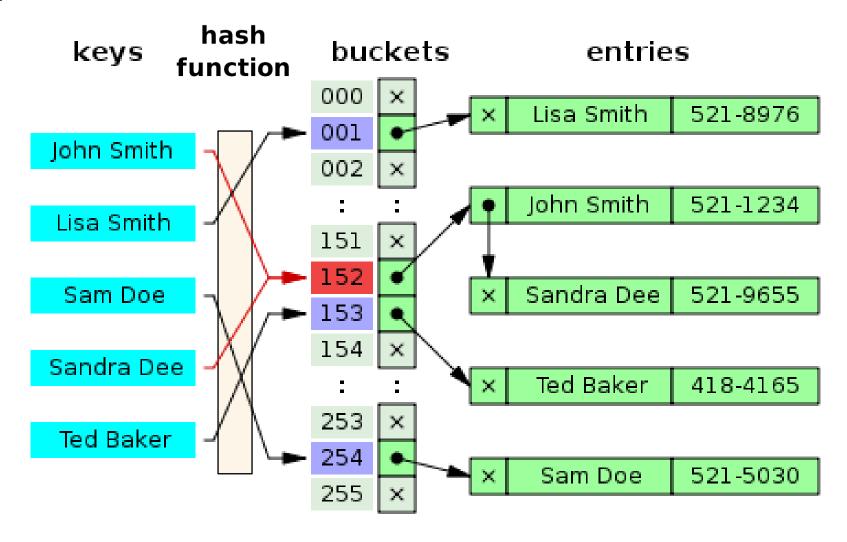
 Provides efficient concurrent operations on key/value pairs via OO & functional programming APIs



- Provides efficient concurrent operations on key/value pairs via OO & functional programming APIs
  - A highly-optimized "associative array"
    - Cannot contain duplicate keys
      - i.e., each key maps to at most one value

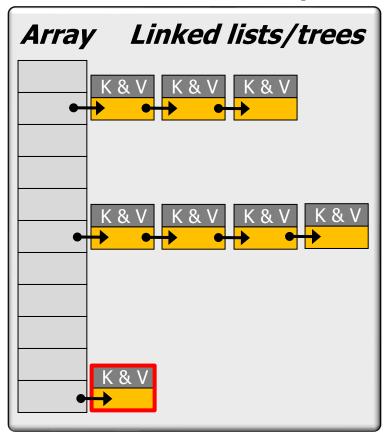
#### <<.lava Class>> • ConcurrentHashMap<K,V> ConcurrentHashMap() ConcurrentHashMap(int) ConcurrentHashMap(Map<? extends K,? extends V>) ConcurrentHashMap(int,float) ConcurrentHashMap(int,float,int) size():int isEmpty():boolean get(Object) containsKey(Object):boolean containsValue(Object):boolean put(K,V) FputVal(K,V,boolean) putAll(Map<? extends K,? extends V>):void remove(Object) replaceNode(Object, V, Object) clear():void keySet() values():Collection<V> entrySet():Set<Entry<K,V>> hashCode():int remove(Object,Object):boolean replace(K,V,V):boolean replace(K,V)

Implemented as a hash table



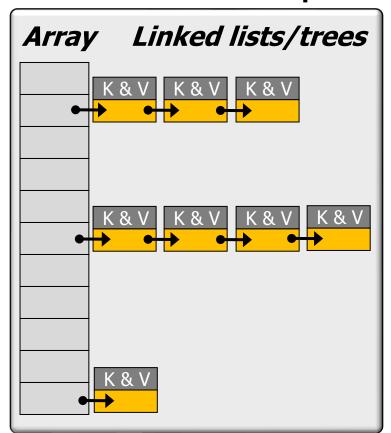
See <a href="mailto:en.wikipedia.org/wiki/Hash\_table">en.wikipedia.org/wiki/Hash\_table</a>

- Implemented as a hash table
  - Insert & retrieve data elements by key



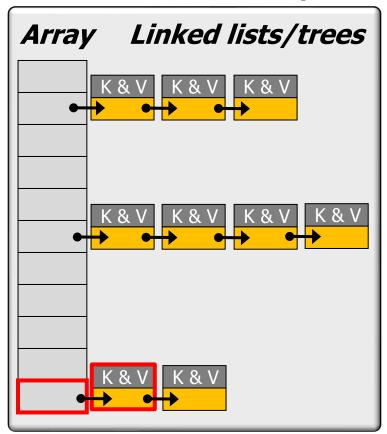
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```
Map<String, Integer> map
  = new ConcurrentHashMap<>();
initializeMap(map);
// Thread T1
map.put("key1", 42);
// Thread T2
Integer value = map.get("key1");
   put() in thread T1 must "happen-
      before" get() in thread T2
```

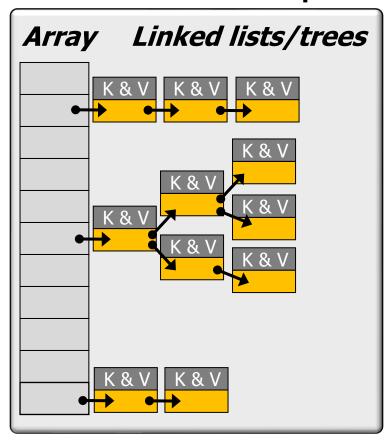


- Implemented as a hash table
  - Insert & retrieve data elements by key
  - Two items that hash to same location in the array are placed in linked list

```
map.put("key2", 1066);
```



- Implemented as a hash table
  - Insert & retrieve data elements by key
  - Two items that hash to same location in the array are placed in linked list
    - In Java 8+, a linked list is replaced by a binary tree when # of elements in a bucket reaches certain threshold



Optimized for multi-core CPUs



## Building a better HashMap

How ConcurrentHashMap offers higher concurrency without compromising thread safety



Brian Goetz Published on August 21, 2003



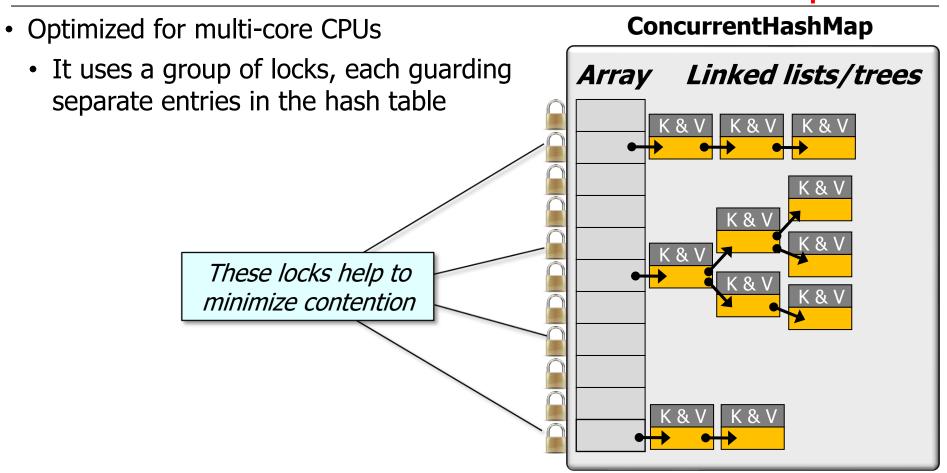
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#### Content series:

+ This content is part of the series: Java theory and practice

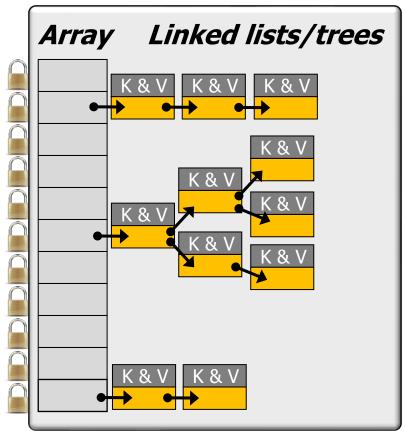
In July's installment of *Java theory and practice* ("Concurrent collections classes"), we reviewed scalability bottlenecks and discussed how to achieve higher concurrency and throughput in shared data structures. Sometimes, the best way to learn is to examine the work of the experts, so this month we're going to look at the implementation of ConcurrentHashMap from Doug Lea's util.concurrent package. A version of ConcurrentHashMap optimized for the new Java Memory Model (JMM), which is being specified by JSR 133, will be included in the java.util.concurrent package in JDK 1.5; the version in util.concurrent has been audited for thread-safety under both the old and new memory models.



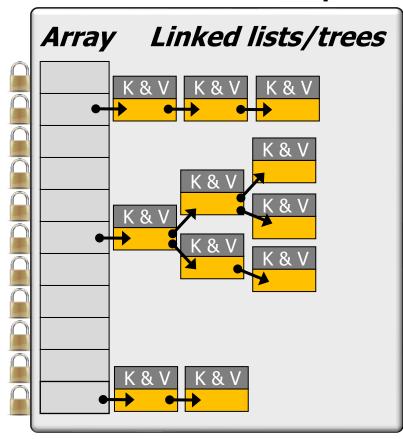
- Optimized for multi-core CPUs
  - It uses a group of locks, each guarding separate entries in the hash table



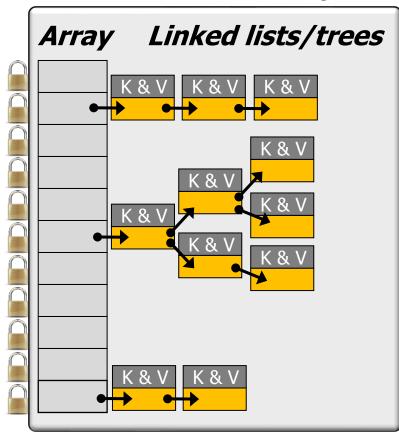
There are common human known uses of this approach!



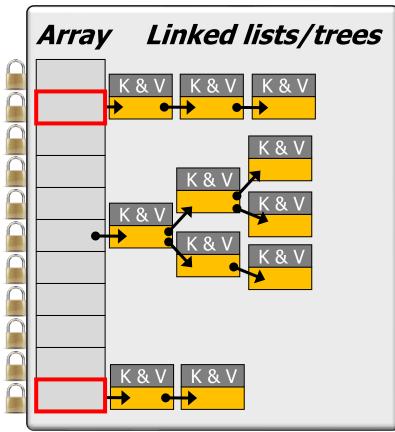
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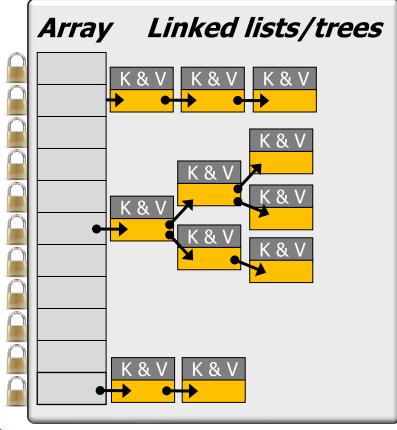


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    - Reads & writes are concurrent if they occur in different lists (or trees)

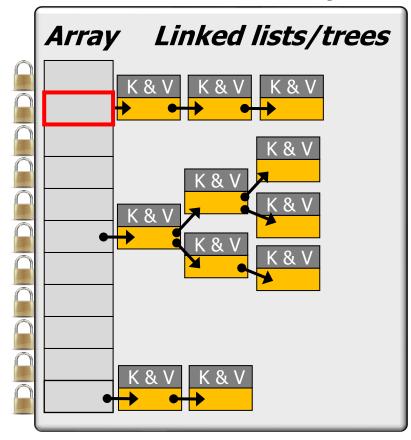


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    - Reads & writes to same list are optimized to avoid locking

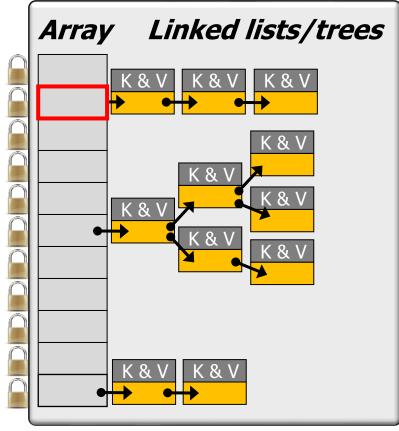




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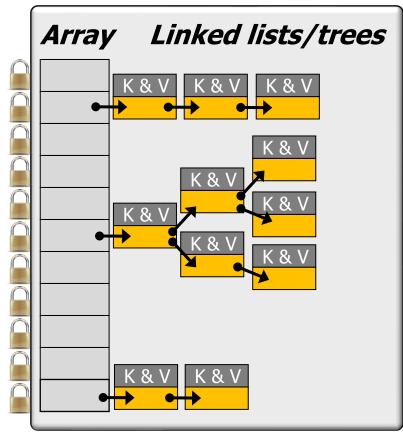


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      - Atomic add to head of list
      - Remove from list by setting data field to null, rebuild list to skip this cell
        - Unreachable cells are eventually garbage collected

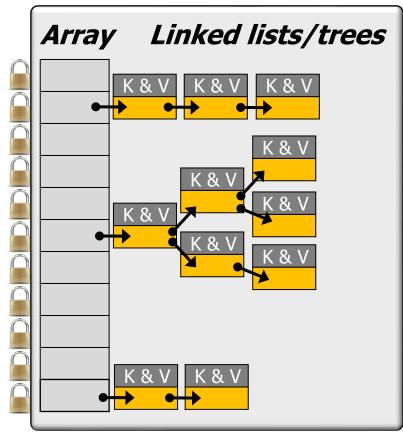


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    - Can be modified during iteration

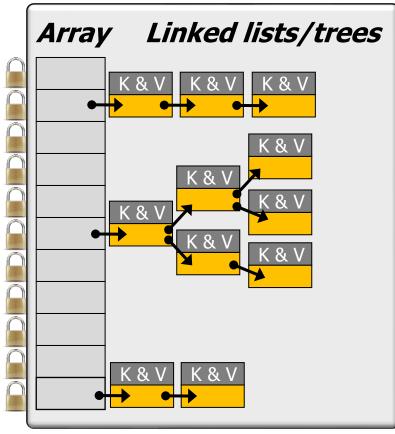




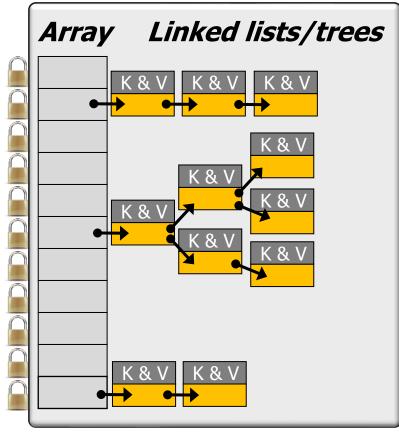
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      - Entire map isn't locked

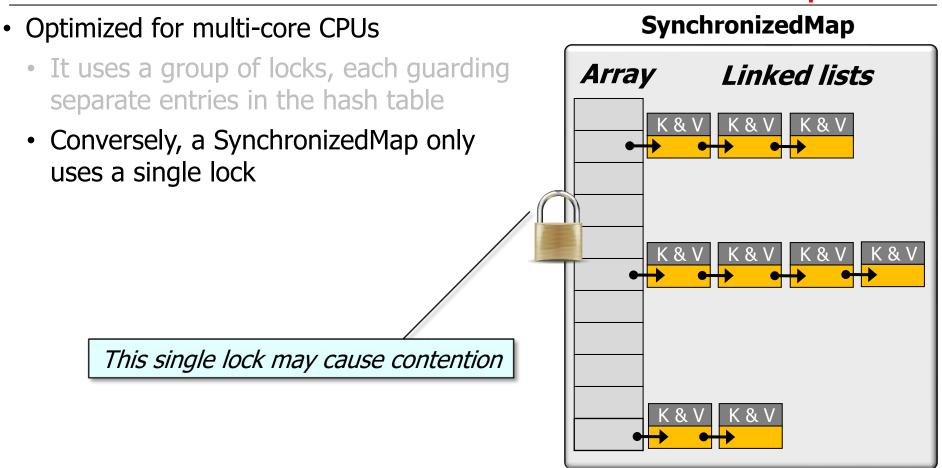


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      - ConcurrentModificationException isn't thrown



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    - Reads & writes are concurrent if they occur in different lists
    - Reads & writes to same list are optimized to avoid locking
    - Can be modified during iteration, e.g.
      - Entire map isn't locked
      - ConcurrentModificationException isn't thrown
      - However, changes may not be visible immediately

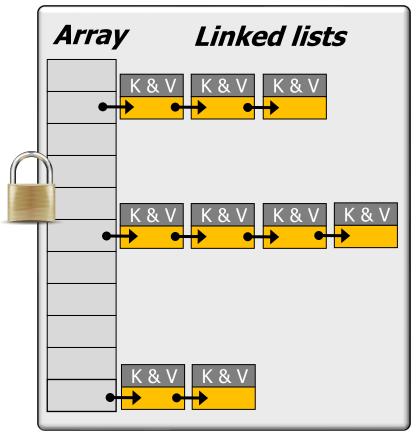




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  - Conversely, a SynchronizedMap only uses a single lock



#### **SynchronizedMap**

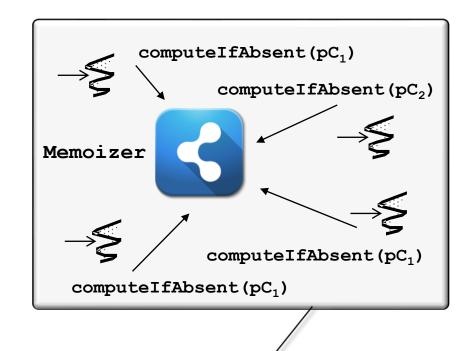


There are also common human known uses of this approach!

- Optimized for multi-core CPUs
  - It uses a group of locks, each guarding separate entries in the hash table
  - Conversely, a SynchronizedMap only uses a single lock
  - ConcurrentHashMaps are thus much more scalable than SynchronizedMaps

```
Starting ConcurrentHashMap
Leaving ConcurrentHashMap
Starting SynchronizedHashMap
Leaving SynchronizedHashMap
Printing 2 results from fastest to slowest
ConcurrentHashMap executed in 351 msecs
SynchronizedHashMap executed in 1409 msecs
```

 Provides "atomic get-and-maybe-set" methods



Only one computation per key is performed even if multiple threads call computeIfAbsent() using the same key

See <a href="mailto:dig.cs.illinois.edu/papers/checkThenAct.pdf">dig.cs.illinois.edu/papers/checkThenAct.pdf</a>

- Provides "atomic get-and-maybe-set" methods, e.g.
  - If key isn't already associated w/a value, compute its value using the given function & enter it into map

```
Use
return map.computeIfAbsent
  (key,
   k -> mappingFunc(k)));
instead of
V value = map.get(key);
if (value == null) {
  value =
    mappingFunc.apply(key);
  if (value != null)
    map.put(key, value);
return value;
```

- Provides "atomic get-and-maybe-set" methods, e.g.
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  - Replaces entry for a key only if currently mapped to some value

```
Use

return map.replace(key, value);
instead of

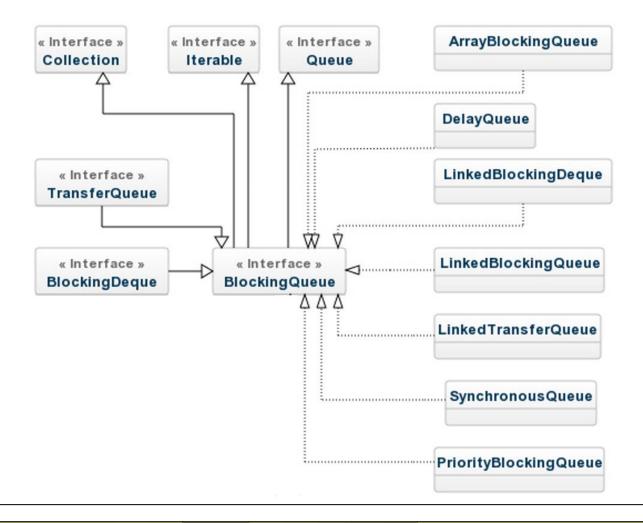
if (map.containsKey(key))
  return map.put(key, value);
else
  return null;
```

- Provides "atomic get-and-maybe-set" methods, e.g.
  - If key isn't already associated w/a value, compute its value using the given function & enter it into map
  - If a key isn't already associated w/a value, associate it with the value
  - Replaces entry for a key only if currently mapped to some value
  - Replaces entry for a key only if currently mapped to given value

```
Use
return map.replace(key,
                    oldValue,
                    newValue);
instead of
if (map.containsKey(key) &&
    Objects.equals(map.get(key),
                  oldValue)) {
   map.put(key, newValue);
   return true;
} else
  return false:
```

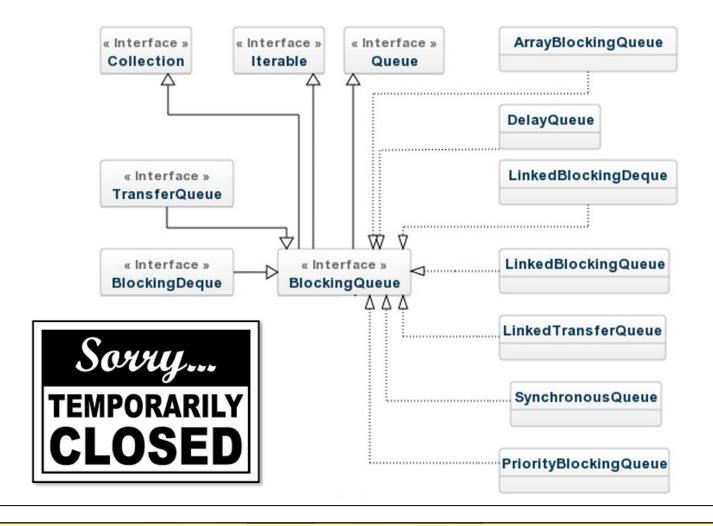
See <u>docs.oracle.com/javase/8/docs/api/java/util/</u> concurrent/ConcurrentHashMap.html#replace

A Queue supporting operations with certain properties



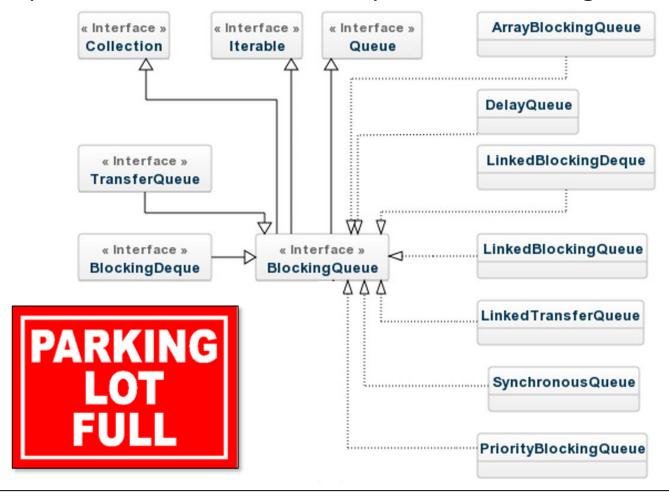
See <a href="https://docs.oracle.com/javase/8/docs/api/java/util/concurrent/BlockingQueue.html">docs.oracle.com/javase/8/docs/api/java/util/concurrent/BlockingQueue.html</a>

- A Queue supporting operations with certain properties
  - wait for the queue to become non-empty when retrieving an element &



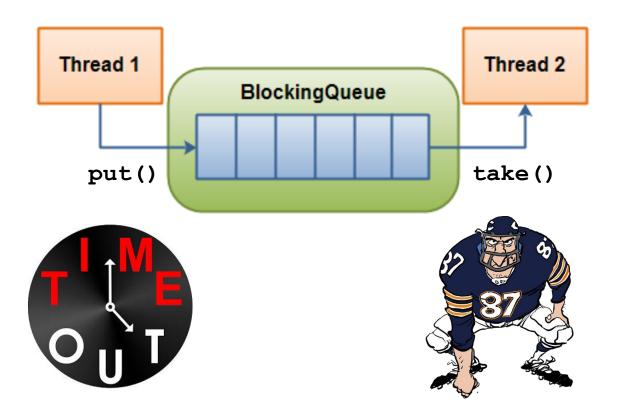
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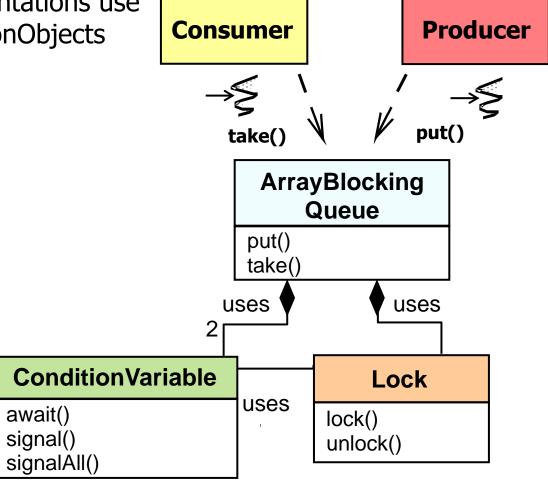


See docs.oracle.com/javase/8/docs/api/java/util/concurrent/BlockingQueue.html

 When adding to a full queue or retrieving from an empty queue clients can either block indefinitely, timeout after waiting for a designated time, or poll



 Many BlockingQueue implementations use Java ReentrantLock & ConditionObjects



## End of Java Concurrent Collections: ConcurrentHash Map & BlockingQueue